CS 591: Introduction to Computer Security

Lecture 2: Voting Machine Study Access Control

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Objectives:

- Review/Discuss Analysis of Diebold machine
- Introduce Access Control

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Discussion

- Thompson, Can You Coun on Voting Machines?, NY Times, January, 2008
 - Reaction?
 - Conflicts of interest in design and deployment?
 - Conflicts of interest in Testing? How independent?
 - Role of the vendor in operations?
 - How to prove the presence of a transient bug? The absence?

Discussion

- If we can make a good ATM why is it hard to make a good voting machine?
- "You have to be able to convince the loser they lost" ... "Not only must the losing candidate believe in the loss; the public has to believe in it, too."

Discussion

 Feldman, Halderman, and Felten, Security Analysis of the Diebold AccuVote-TS Voting Machine, September 2006

- Reaction to the paper?

Discussion Questions

- What was the basic architecture of the voting machine?
- How did FHF steal votes?
- What other attacks did FHF consider?
- How did the viral propagation mechanism work?

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Discussion Questions

- Is the analysis credible?
- Is the threat model credible?
- Is this representative of commercial systems today?
- Did Diebold follow best practices?
- Are the FHF results reproducible?
- Did Felton's lab follow a sound methodology in analyzing the machine?

Discussion Questions

- Having read the analysis of the Diebold machine, are you surprised that Sequoia used a threat of law suit to prevent Felten's lab from analyzing their machine?
- Having seen this analysis of a fielded commercial system, are you more or less concerned about the discrepencies observed in Union County elections?

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Discussion Questions

- Do you like Oregon's vote by mail system?
- Are Appel's comments on Minnesota relevant to voting in Oregon?

Case Study

- We will use the FHF paper as a case study
- As we encounter concepts we will attempt to instantiate them in the context of the voting machine domain

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Voting Machine Architecture



Boot Process

- Boot device specified by hardware jumpers (inside box)
 - EPROM
 - on-board flash (default)
 - ext flash
- On Boot:
 - Copy bootloader into RAM; init hardware
 - Scan Removable flash for special files
 - "fboot.nb0" => replace bootloader in on-board flash
 - "nk.bin" => replace OS in on-board flash
 - "EraseFFX.bsq" => erase file system on on-board flash
 - If no special files uncompress OS image
 - Jump to entry point of OS

Boot (continued)

- On OS start up:
 - run Filesys.exe
 - unpacks registry
 - runs programs in HKEY_LOCAL_MACHINE\Init
 - shell.exe (debug shell)
 - device.exe (Device manager)
 - gwes.exe (graphics and event)
 - taskman.exe (Task Manager)
 - Device.exe mounts file systems
 - \ (root): RAM only
 - \FFX: mount point for on-board flash
 - \Storage Card: mount point for removable flash

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Boot (continued)

- Customized taskman.exe
 - Check removable flash
 - explorer.glb => launch windows explorer
 - *.ins => run proprietary scripts
 - (script language has buffer overflow vulnerabilities)
 - used to configure election data
 - default => launch "BallotStation"
 - \FFX\Bin\BallotStation.exe

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BallotStation

- Four modes: pre-download, preelection testing, election, post-election
- Mode recorded in election results file
 - \Storage Card\CurrentElection\election.brs

Stealing Votes

- Malicious processes runs in parallel with BallotStation
- Polls election results file every 15 seconds
 - If election mode and new results
 - temporarily suspend Ballot Station
 - steal votes
 - resume Ballot Station

Viral propagation

- Malicious bootloader
 - Infects host by replacing existing bootloader in on-board flash
 - subsequent bootloader updates print appropriate messages but do nothing
- fboot.nb0
 - package contains malicious boot loader
 - and vote stealing software

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Objectives

- Introduce the concept of Access Control
- Relate mechanism to Confidentiality, Integrity and Availability

Access Control Model

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Articulating Policy

- How do we articulate a security policy?
- How do we provide mechanisms to enforce policy?
- Voting
 - Different individuals in different roles
 Voter, Poll worker, ...
 - Different actions
 - Vote, define ballot, start and stop election, ...
 - Logical and physical entities
 - Ballot, stored tally, final tally, voting machine, removable flash, on-board flash, ...

Ad hoc policies

- Discus
 - Only voters should vote
 - Only poll workers should start and start elections

Access Control Matrix Model

– Lampson '71, refined by Graham and Denning ('71, '72)

- Concepts

- Objects, the protected entities, O
- Subjects, the active entities acting on the objects, S
- Rights, the controlled operations subjects can perform on objects, R
- Access Control Matrix, A, maps Objects and Subjects to sets of Rights
- State: (S, O, A)

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Voting: Subjects, Objects, Rights

- Subjects: (Roles)
 - Voter, Poll worker, ...
- Rights: (Actions)
 - Vote, define ballot, start and stop election, ...
- Objects: (Logical and physical entities)
 - Ballot, stored tally, final tally, voting machine, removable flash, on-board flash, ...
- Question: Is every voter a subject? Or is the role of voter a subject? One-person-one-vote?

Exercise

 Sketch Access Control Matrix (ACM) for Voting

	Ballot	Stored Tally	Final Tally	
Voter	read	increment		
Poll			print	
worker				

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Questions

- What about modes?
 - Once the election starts the ballot should not change
 - Voters should only vote when the election is happening

Questions

- Levels of abstraction
 - Some objects are physical, some are logical
 - When considering the programming model you now have processes and files (and possibly modes of operation)
- Exercise:
 - Sketch ACMs with processes as subjects and files as objects for voting and postelection modes

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Exercise

- Compare the ACMs for files and processes with the original ACM
- Is every operation specified in the original feasible in the refined ACMs?
- Is every feasible operation in the refined ACMs allowed in the original?

Mechanisms

- Policy specifies abstract goals
- Mechanisms are concrete devices, algorithms, or processes that assist in implementing a policy
- For example, passwords are a mechanism that can support an authentication policy
 - Mechanisms are not always perfect!

Access Control Mechanisms

- Most operating systems provide some mechanisms for supporting access control
- Typically:
 - Processes are associated with users (or user identification numbers), which are the subjects
 - Files are objects
 - Rights are: read, write, append, execute, search, ...

Applying the Mechanism

- Can a generic Access Control mechanism help make the Voting machine more trustworthy?
- What about modes?
 - Mode is not part of typical AC mechanisms
 - However rights can be changed
 - A typical right is "own" which in discretionary access control generally allows the subject to change rights
 - Analysis of systems that actively change rights is potentially difficult

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Limitations on Mechanisms

- Simple mechanisms are preferred
- All computational mechanisms must be decidable
- In general, useful mechanisms must be computationally cheap

Access Control

- Is Access Control biased to
 - Confidentiality
 - Integrity
 - Availability
- Exercise
 - Develop scenarios in which a confidentiality (integrity, availability) property is expressed using an access control matrix

Model vs. Mechanism

- Earlier I presented the model of the AC Matrix
- Does UNIX implement the full AC Matrix?
 - What key simplifications does UNIX adopt?Why?
- Is the full ACM mechanism a good idea?
 - Is it a good model?

A Good Model

 ACM is a good model because any mechanism of compatible granularity can be described in terms of how it approximates the ACM model

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Next Lecture

- Discussion
 - Tibet

- NY Times article <u>http://www.nytimes.com/2009/03/29/</u> technology/29spy.html?emc=eta1
- Nagaraja and Anderson tech report <u>http://www.cl.cam.ac.uk/techreports/UCAM-CL-</u> <u>TR-746.html</u>