

Converting an NFA to a DFA

Given:

A *non-deterministic* finite state machine (NFA)

Goal:

Convert to an equivalent *deterministic* finite state machine (DFA)

Why?

Faster recognizer!

Approach:

Consider simulating a NFA.

Work with sets of states.

IDEA: Each *state* in the DFA will correspond to a *set of* NFA states.

Worst-case:

There can be an exponential number $O(2^N)$ of sets of states.

The DFA can have exponentially many more states than the NFA

... but this is rare.

NFA to DFA

Input: A NFA

$S = \text{States} = \{ s_0, s_1, \dots, s_N \} = S_{\text{NFA}}$

$\delta = \text{Move function} = \text{Move}_{\text{NFA}}$

$\text{Move}(S, a) \rightarrow \text{Set of states}$

Output: A DFA

$S = \text{States} = \{ ?, ?, \dots, ? \} = S_{\text{DFA}}$

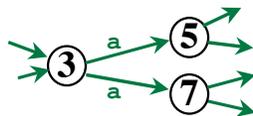
$\delta = \text{Move function} = \text{Move}_{\text{DFA}}$

$\text{Move}(s, a) \rightarrow \text{Single state from } S_{\text{DFA}}$

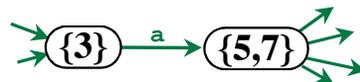
Main Idea:

Each state in S_{DFA} will be a set of states from the NFA

$S_{\text{DFA}} = \{ \{ \dots \}, \{ \dots \}, \dots, \{ \dots \} \}$

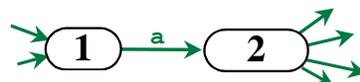


NFA



DFA

(The names of the states is arbitrary and can be changed later, if desired.)



Algorithm: Convert NFA to DFA

We'll use...

- $\text{Move}_{\text{NFA}}(\mathbf{S}, a)$ the transition function from NFA
- $\epsilon\text{-Closure}(s)$ where s is a single state from NFA
- $\epsilon\text{-Closure}(\mathbf{S})$ where \mathbf{S} is a set of states from NFA

We'll construct...

- \mathbf{S}_{DFA} the set of states in the DFA
Initially, we'll set \mathbf{S}_{DFA} to $\{\}$
- Add \mathbf{X} to \mathbf{S}_{DFA} where \mathbf{X} is some *set of* NFA states
Example: "Add $\{\mathbf{3,5,7}\}$ to \mathbf{S}_{DFA} "
- We'll "mark" some of the states in the DFA.
Marked = "We've done this one" (✓)
Unmarked = "Still need to do this one"

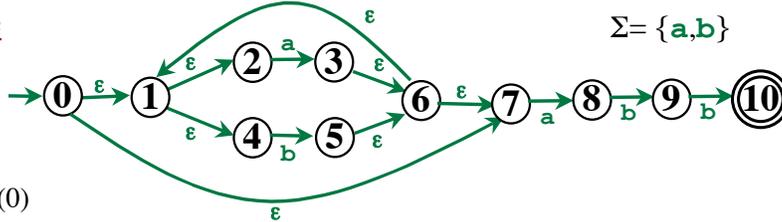
- $\text{Move}_{\text{DFA}}(\mathbf{T}, b)$ The transition function from DFA
To add an edge to the growing DFA...

Set $\text{Move}_{\text{DFA}}(\mathbf{T}, b)$ to \mathbf{S}



...where \mathbf{S} and \mathbf{T} are sets of NFA states

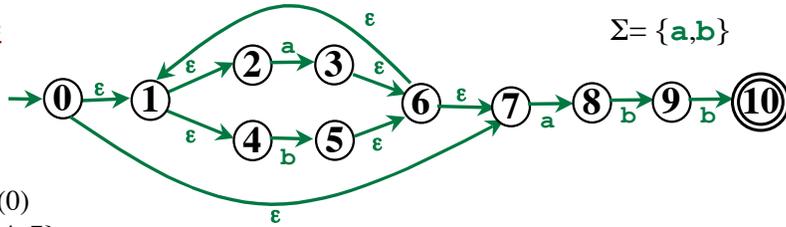
Example



Start state:
 $\epsilon\text{-Closure}(0)$
 =

Lexical Analysis - Part 3

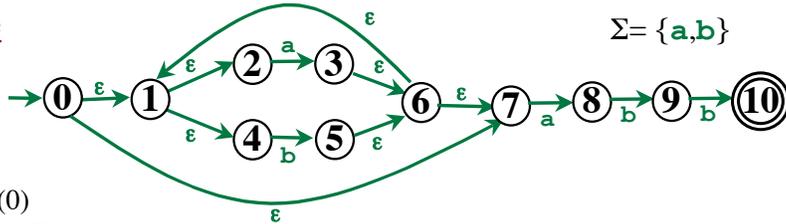
Example



Start state:
 ϵ -Closure (0)
= {0, 1, 2, 4, 7}

Lexical Analysis - Part 3

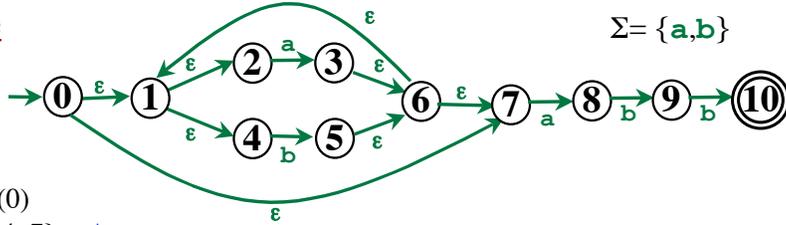
Example



Start state:
 ϵ -Closure (0)
= {0, 1, 2, 4, 7} = A

Lexical Analysis - Part 3

Example



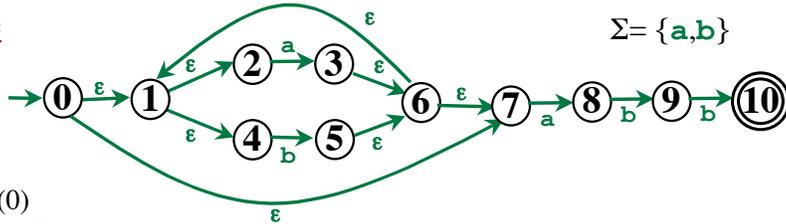
Start state:
 ϵ -Closure (0)
 = {0, 1, 2, 4, 7} = A

Move_{DFA}(A,a)
 =

Move_{DFA}(A,b)
 =

Lexical Analysis - Part 3

Example



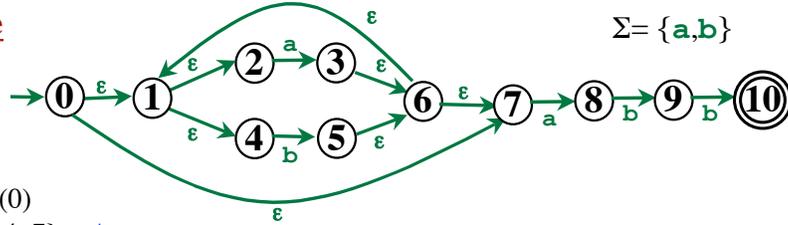
Start state:
 ϵ -Closure (0)
 = {0, 1, 2, 4, 7} = A

Move_{DFA}(A,a)
 = ϵ -Closure (Move_{NFA}(A,a))
 =

Move_{DFA}(A,b)
 =

Lexical Analysis - Part 3

Example



Start state:

$$\begin{aligned} &\epsilon\text{-Closure}(0) \\ &= \{0, 1, 2, 4, 7\} = A \end{aligned}$$

$\text{Move}_{\text{DFA}}(A, a)$

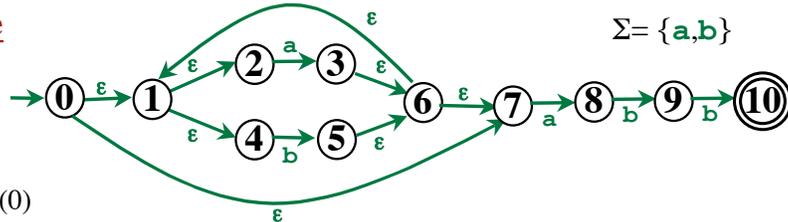
$$\begin{aligned} &= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(A, a)) \\ &= \epsilon\text{-Closure}(\{3, 8\}) \\ &= \end{aligned}$$

$\text{Move}_{\text{DFA}}(A, b)$

=

Lexical Analysis - Part 3

Example



Start state:

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$\text{Move}_{\text{DFA}}(A, a)$

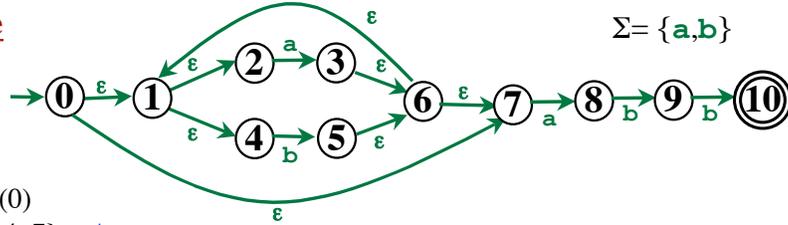
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$\text{Move}_{\text{DFA}}(A, b)$

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Lexical Analysis - Part 3

Example



Start state:

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$\text{Move}_{\text{DFA}}(A, a)$

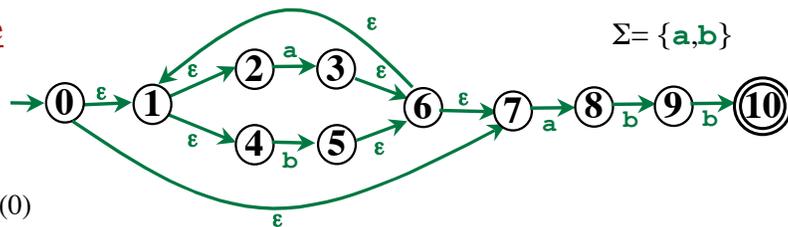
$$\begin{aligned} &= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(A, a)) \\ &= \epsilon\text{-Closure}(\{3, 8\}) \\ &= \{1, 2, 3, 4, 6, 7, 8\} = B \end{aligned}$$

$\text{Move}_{\text{DFA}}(A, b)$

=

Lexical Analysis - Part 3

Example



Start state:

$$\begin{aligned} &\epsilon\text{-Closure}(0) \\ &= \{0, 1, 2, 4, 7\} = A \end{aligned}$$

$\text{Move}_{\text{DFA}}(A, a)$

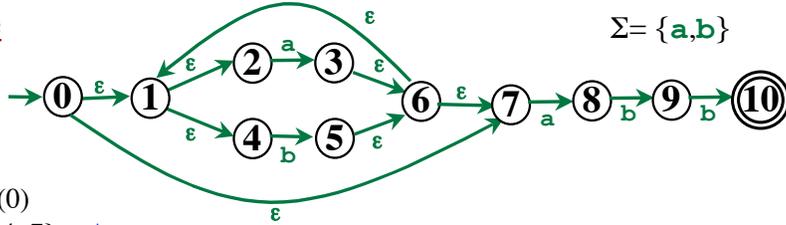
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$\text{Move}_{\text{DFA}}(A, b)$

$$\begin{aligned} &= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(A, b)) \\ &= \end{aligned}$$

Lexical Analysis - Part 3

Example



Start state:

$$\begin{aligned} &\epsilon\text{-Closure}(0) \\ &= \{0, 1, 2, 4, 7\} = A \end{aligned}$$

$\text{Move}_{\text{DFA}}(A, a)$

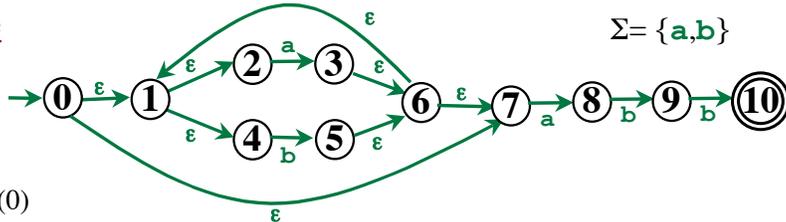
$$\begin{aligned} &= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(A, a)) \\ &= \epsilon\text{-Closure}(\{3, 8\}) \\ &= \{1, 2, 3, 4, 6, 7, 8\} = B \end{aligned}$$

$\text{Move}_{\text{DFA}}(A, b)$

$$\begin{aligned} &= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(A, b)) \\ &= \epsilon\text{-Closure}(\{5\}) \\ &= \end{aligned}$$

Lexical Analysis - Part 3

Example



Start state:

$$\begin{aligned} &\epsilon\text{-Closure}(0) \\ &= \{0, 1, 2, 4, 7\} = A \end{aligned}$$

$\text{Move}_{\text{DFA}}(A, a)$

$$\begin{aligned} &= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(A, a)) \\ &= \epsilon\text{-Closure}(\{3, 8\}) \\ &= \{1, 2, 3, 4, 6, 7, 8\} = B \end{aligned}$$

$\text{Move}_{\text{DFA}}(A, b)$

$$\begin{aligned} &= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(A, b)) \\ &= \epsilon\text{-Closure}(\{5\}) \\ &= \{1, 2, 4, 5, 6, 7\} = C \end{aligned}$$

Example $\Sigma = \{a,b\}$

Start state:
 ϵ -Closure (0)
 = {0, 1, 2, 4, 7} = A

Move_{DFA}(A,a)
 = ϵ -Closure (Move_{NFA}(A,a))
 = ϵ -Closure ({3,8})
 = {1,2,3,4,6,7,8} = B

Move_{DFA}(A,b)
 = ϵ -Closure (Move_{NFA}(A,b))
 = ϵ -Closure ({5})
 = {1,2,4,5,6,7} = C

So far:

Example $\Sigma = \{a,b\}$

Start state:
 ϵ -Closure (0)
 = {0, 1, 2, 4, 7} = A

Move_{DFA}(A,a)
 = ϵ -Closure (Move_{NFA}(A,a))
 = ϵ -Closure ({3,8})
 = {1,2,3,4,6,7,8} = B

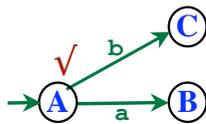
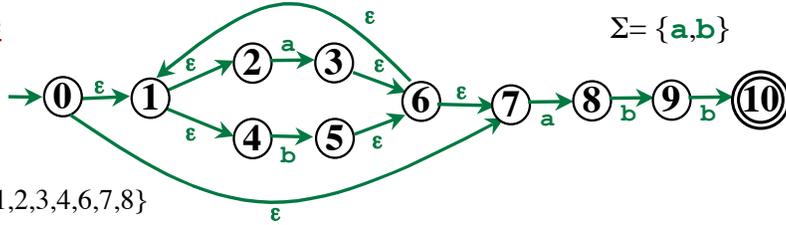
Move_{DFA}(A,b)
 = ϵ -Closure (Move_{NFA}(A,b))
 = ϵ -Closure ({5})
 = {1,2,4,5,6,7} = C

So far:

A is now done; mark it!
 B and C are unmarked.
 Let's do B next...

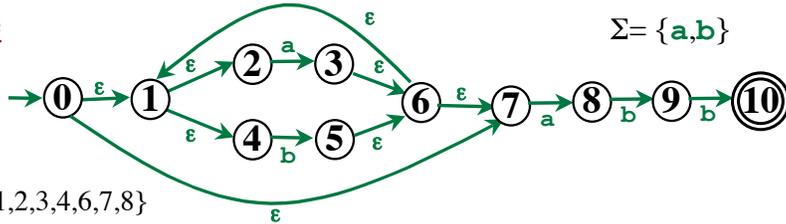
Lexical Analysis - Part 3

Example

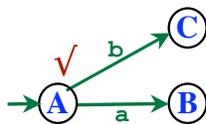


Lexical Analysis - Part 3

Example

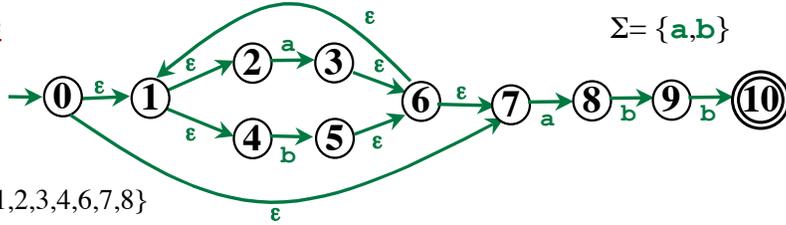


$\text{Move}_{\text{DFA}}(B, a)$
 $=$



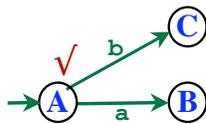
Lexical Analysis - Part 3

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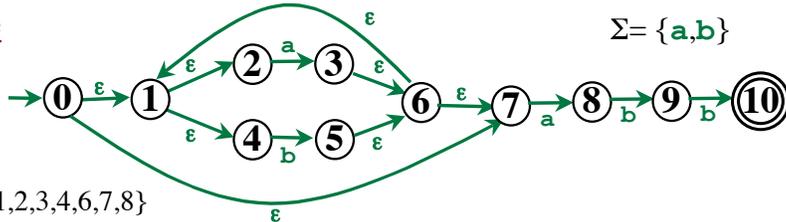
Process $B = \{1,2,3,4,6,7,8\}$

$\text{Move}_{\text{DFA}}(B,a)$
 $= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(B,a))$
 $=$



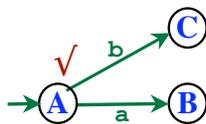
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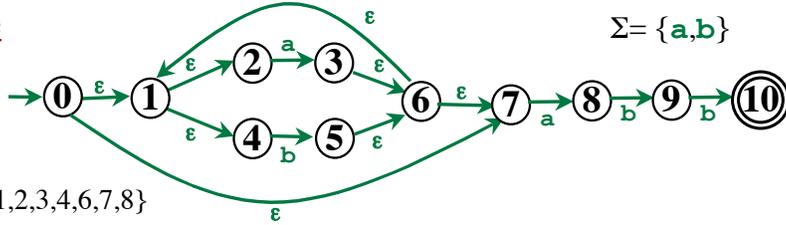
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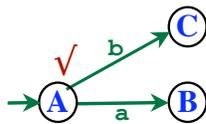
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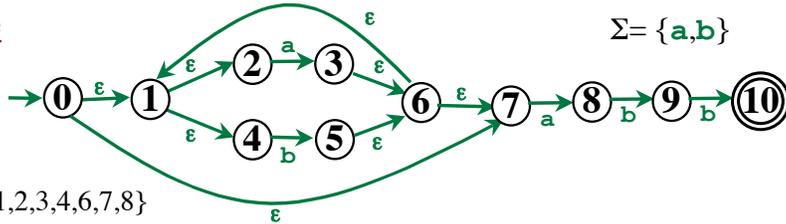
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 $= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(B,a))$
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 $= \{1,2,3,4,6,7,8\} = B$



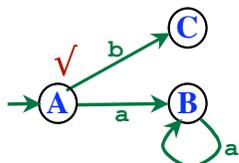
Lexical Analysis - Part 3

Example



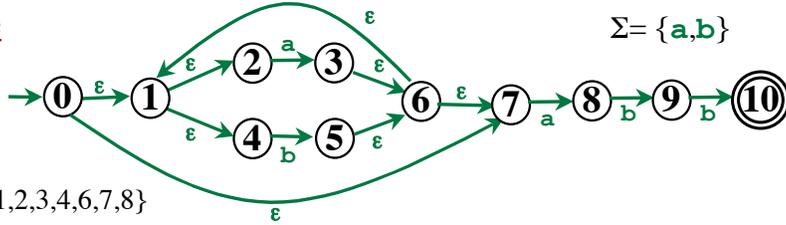
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Lexical Analysis - Part 3

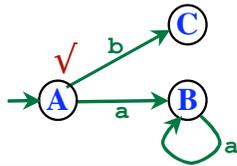
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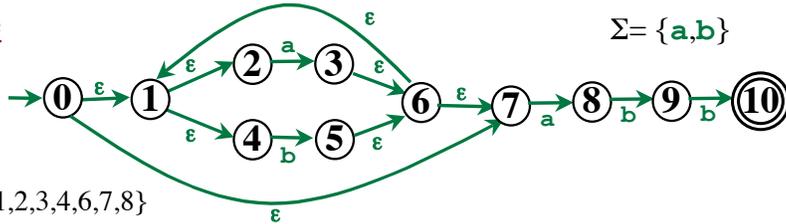
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$\text{Move}_{\text{DFA}}(B, b)$
 $=$



Lexical Analysis - Part 3

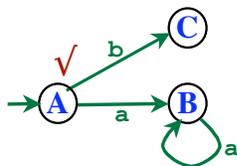
Example



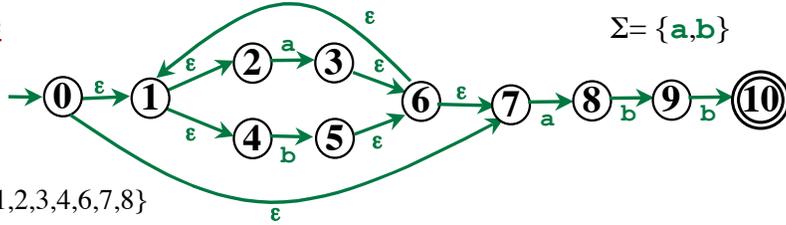
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 $= \epsilon\text{-Closure}(\{3, 8\})$
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$\text{Move}_{\text{DFA}}(B, b)$
 $= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(B, b))$
 $=$



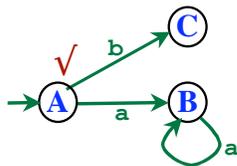
Example



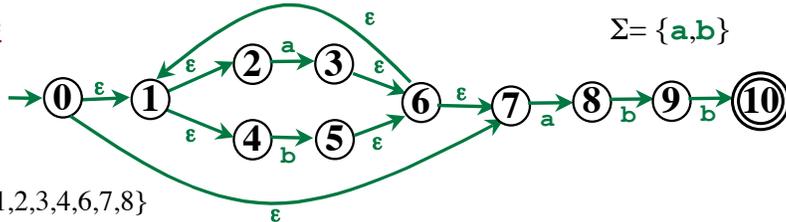
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$\text{Move}_{\text{DFA}}(B,b)$
 $= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(B,b))$
 $= \epsilon\text{-Closure}(\{5,9\})$
 $=$



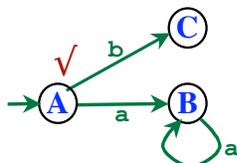
Example



Process $B = \{1,2,3,4,6,7,8\}$

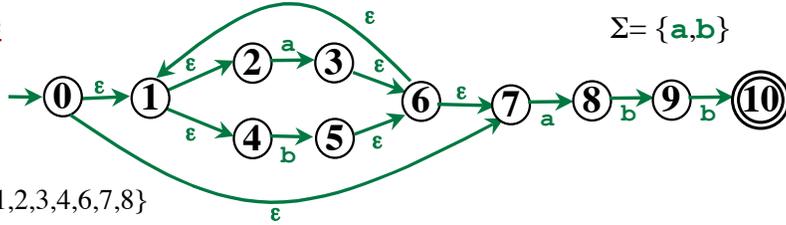
$\text{Move}_{\text{DFA}}(B,a)$
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 $= \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(B,b)$
 $= \epsilon\text{-Closure}(\text{Move}_{\text{NFA}}(B,b))$
 $= \epsilon\text{-Closure}(\{5,9\})$
 $= \{1,2,4,5,6,7,9\} = D$



Lexical Analysis - Part 3

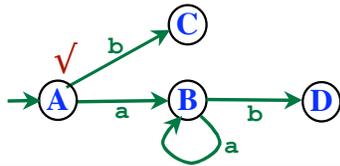
Example



Process $B = \{1,2,3,4,6,7,8\}$

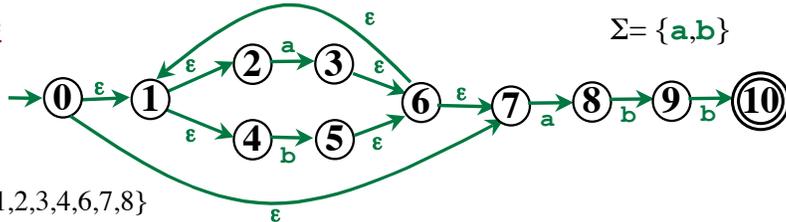
$\text{Move}_{\text{DFA}}(B,a)$
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$\text{Move}_{\text{DFA}}(B,b)$
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 $= \epsilon\text{-Closure}(\{5,9\})$
 $= \{1,2,4,5,6,7,9\} = D$



Lexical Analysis - Part 3

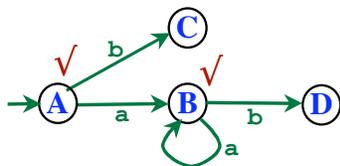
Example



Process $B = \{1,2,3,4,6,7,8\}$

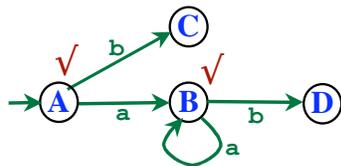
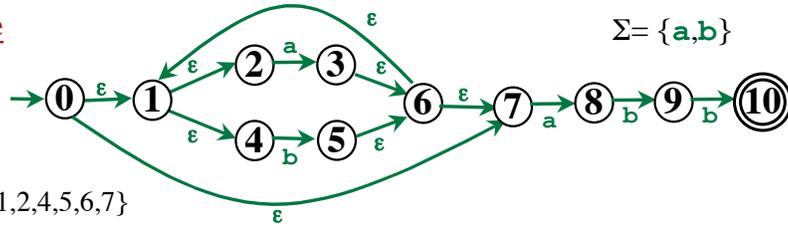
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$\text{Move}_{\text{DFA}}(B,b)$
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 $= \{1,2,4,5,6,7,9\} = D$



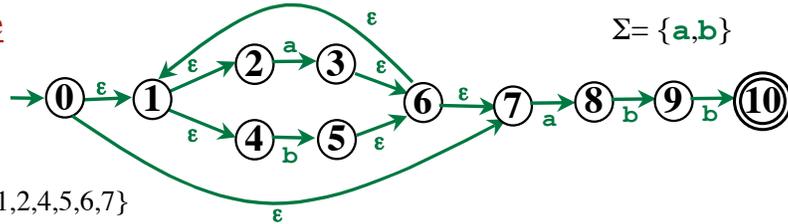
Lexical Analysis - Part 3

Example



Lexical Analysis - Part 3

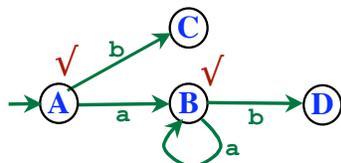
Example



Process $C = \{1,2,4,5,6,7\}$

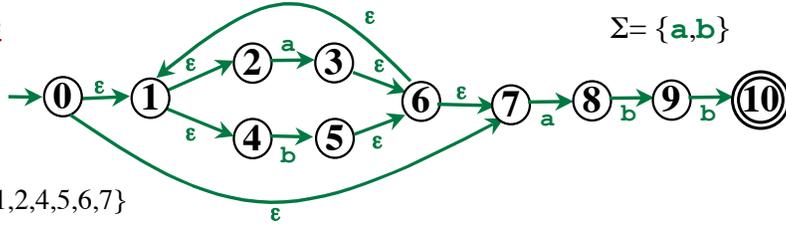
$\text{Move}_{\text{DFA}}(C,a) =$

$\text{Move}_{\text{DFA}}(C,b) =$



Lexical Analysis - Part 3

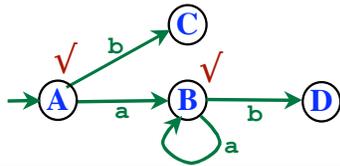
Example



Process $C = \{1,2,4,5,6,7\}$

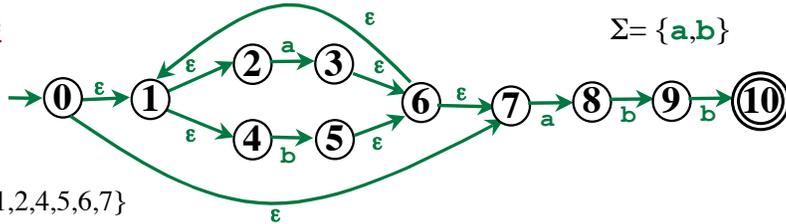
$\text{Move}_{\text{DFA}}(C,a) = \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(C,b) =$



Lexical Analysis - Part 3

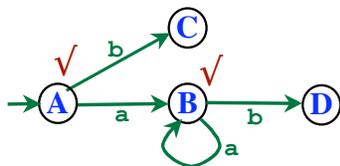
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Process $C = \{1,2,4,5,6,7\}$

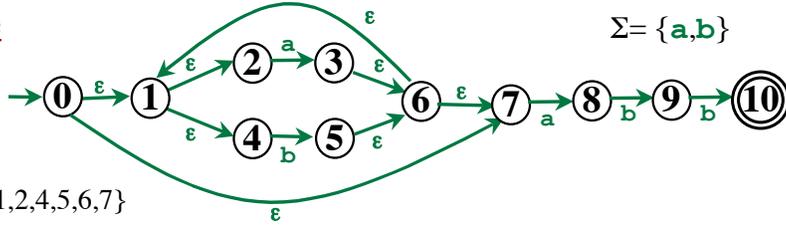
$\text{Move}_{\text{DFA}}(C,a) = \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(C,b) = \{1,2,4,5,6,7\} = C$



Lexical Analysis - Part 3

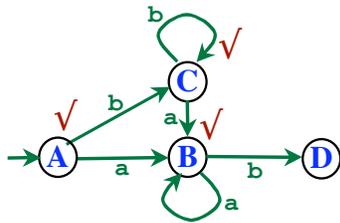
Example



Process $C = \{1,2,4,5,6,7\}$

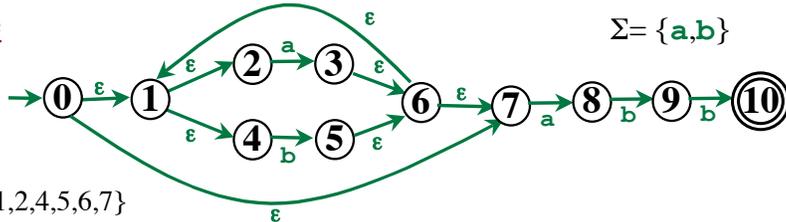
$\text{Move}_{\text{DFA}}(C,a) = \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(C,b) = \{1,2,4,5,6,7\} = C$



Lexical Analysis - Part 3

Example



Process $C = \{1,2,4,5,6,7\}$

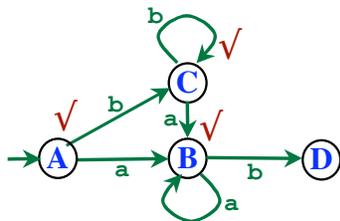
$\text{Move}_{\text{DFA}}(C,a) = \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(C,b) = \{1,2,4,5,6,7\} = C$

Process $D = \{1,2,4,5,6,7,9\}$

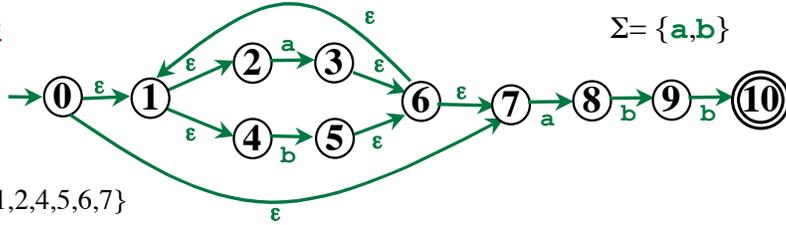
$\text{Move}_{\text{DFA}}(D,a) =$

$\text{Move}_{\text{DFA}}(D,b) =$



Lexical Analysis - Part 3

Example



Process $C = \{1,2,4,5,6,7\}$

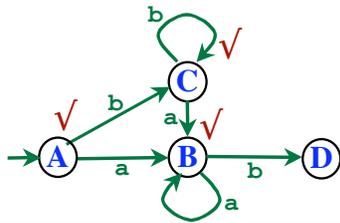
$\text{Move}_{\text{DFA}}(C,a) = \{1,2,3,4,6,7,8\} = B$

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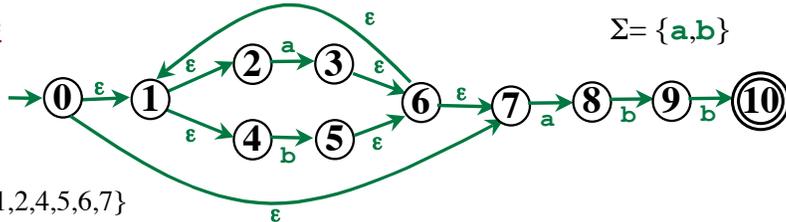
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Lexical Analysis - Part 3

Example



Process $C = \{1,2,4,5,6,7\}$

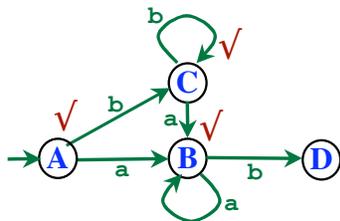
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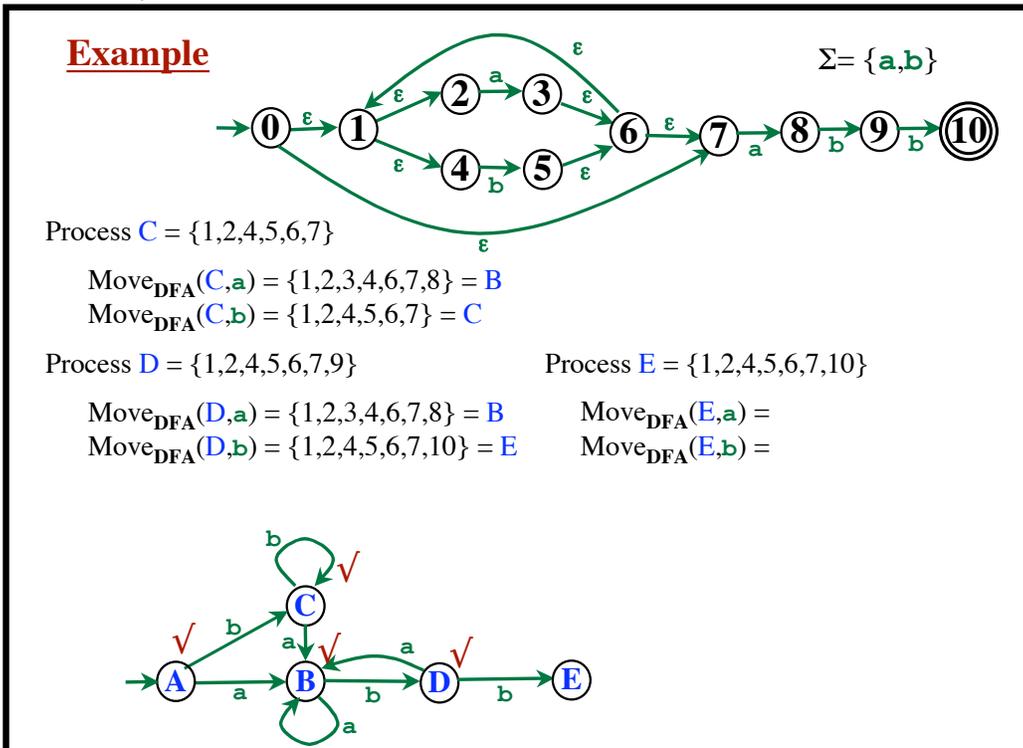
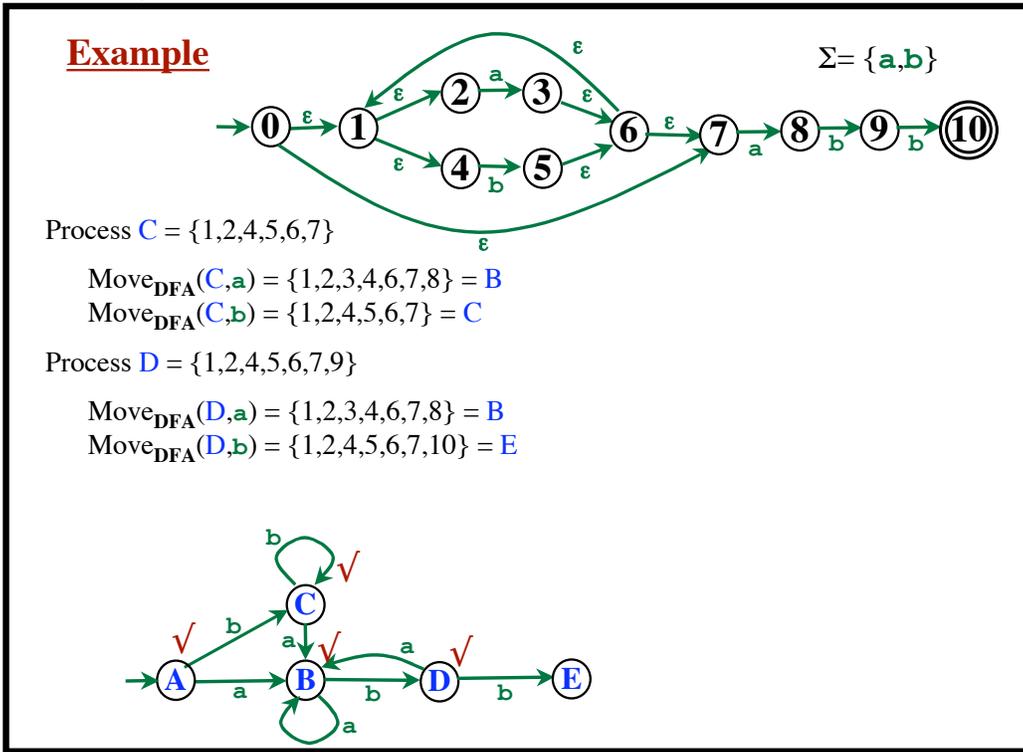
$\text{Move}_{\text{DFA}}(C,b) = \{1,2,4,5,6,7\} = C$

Process $D = \{1,2,4,5,6,7,9\}$

$\text{Move}_{\text{DFA}}(D,a) = \{1,2,3,4,6,7,8\} = B$

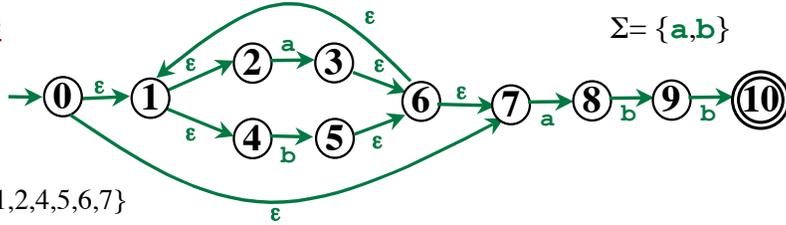
$\text{Move}_{\text{DFA}}(D,b) = \{1,2,4,5,6,7,10\} = E$





Lexical Analysis - Part 3

Example



Process $C = \{1,2,4,5,6,7\}$

$\text{Move}_{\text{DFA}}(C,a) = \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(C,b) = \{1,2,4,5,6,7\} = C$

Process $D = \{1,2,4,5,6,7,9\}$

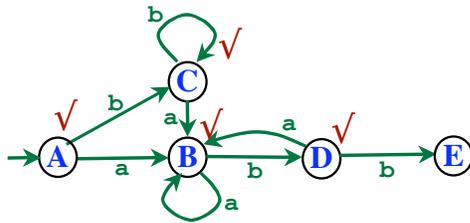
Process $E = \{1,2,4,5,6,7,10\}$

$\text{Move}_{\text{DFA}}(D,a) = \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(E,a) = \{1,2,3,4,6,7,8\} = B$

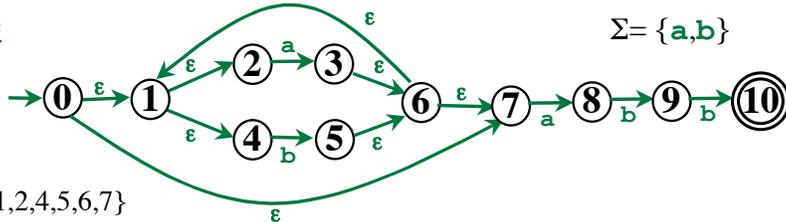
$\text{Move}_{\text{DFA}}(D,b) = \{1,2,4,5,6,7,10\} = E$

$\text{Move}_{\text{DFA}}(E,b) =$



Lexical Analysis - Part 3

Example



Process $C = \{1,2,4,5,6,7\}$

$\text{Move}_{\text{DFA}}(C,a) = \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(C,b) = \{1,2,4,5,6,7\} = C$

Process $D = \{1,2,4,5,6,7,9\}$

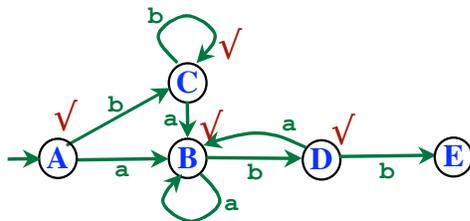
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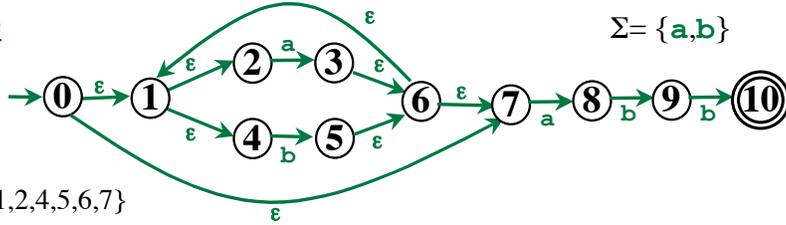
$\text{Move}_{\text{DFA}}(D,b) = \{1,2,4,5,6,7,10\} = E$

$\text{Move}_{\text{DFA}}(E,b) = \{1,2,4,5,6,7\} = C$



Lexical Analysis - Part 3

Example



Process $C = \{1, 2, 4, 5, 6, 7\}$

$\text{Move}_{\text{DFA}}(C, a) = \{1, 2, 3, 4, 6, 7, 8\} = B$

$\text{Move}_{\text{DFA}}(C, b) = \{1, 2, 4, 5, 6, 7\} = C$

Process $D = \{1, 2, 4, 5, 6, 7, 9\}$

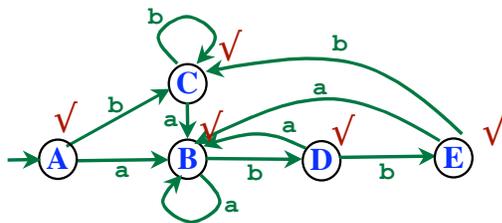
$\text{Move}_{\text{DFA}}(D, a) = \{1, 2, 3, 4, 6, 7, 8\} = B$

$\text{Move}_{\text{DFA}}(D, b) = \{1, 2, 4, 5, 6, 7, 10\} = E$

Process $E = \{1, 2, 4, 5, 6, 7, 10\}$

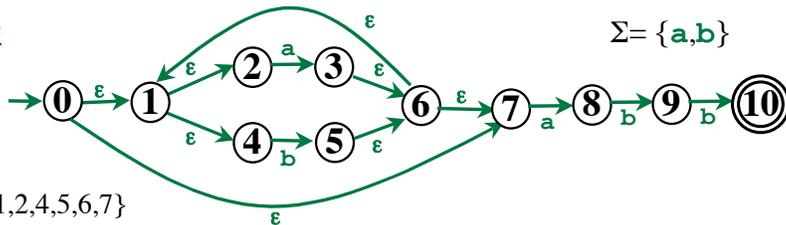
$\text{Move}_{\text{DFA}}(E, a) = \{1, 2, 3, 4, 6, 7, 8\} = B$

$\text{Move}_{\text{DFA}}(E, b) = \{1, 2, 4, 5, 6, 7\} = C$



Lexical Analysis - Part 3

Example



Process $C = \{1, 2, 4, 5, 6, 7\}$

$\text{Move}_{\text{DFA}}(C, a) = \{1, 2, 3, 4, 6, 7, 8\} = B$

$\text{Move}_{\text{DFA}}(C, b) = \{1, 2, 4, 5, 6, 7\} = C$

Process $D = \{1, 2, 4, 5, 6, 7, 9\}$

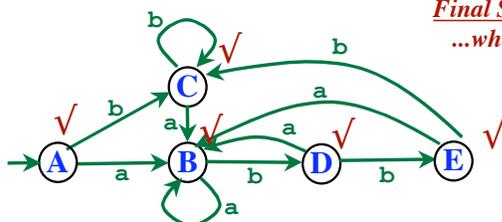
$\text{Move}_{\text{DFA}}(D, a) = \{1, 2, 3, 4, 6, 7, 8\} = B$

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$\text{Move}_{\text{DFA}}(E, b) = \{1, 2, 4, 5, 6, 7\} = C$



Final States in DFA?
...which state(s) contain 10?

Example $\Sigma = \{a,b\}$

Process $C = \{1,2,4,5,6,7\}$

$\text{Move}_{\text{DFA}}(C,a) = \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(C,b) = \{1,2,4,5,6,7\} = C$

Process $D = \{1,2,4,5,6,7,9\}$

$\text{Move}_{\text{DFA}}(D,a) = \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(D,b) = \{1,2,4,5,6,7,10\} = E$

Process $E = \{1,2,4,5,6,7,10\}$

$\text{Move}_{\text{DFA}}(E,a) = \{1,2,3,4,6,7,8\} = B$

$\text{Move}_{\text{DFA}}(E,b) = \{1,2,4,5,6,7\} = C$

Final Result:

Final States in DFA?
...which state(s) contain 10?

Algorithm: Convert NFA to DFA

```

SDFA = {}
Add ε-Closure(s0) to SDFA as the start state
Set the only state in SDFA to "unmarked"
while SDFA contains an unmarked state do
  Let T be that unmarked state
  Mark T
  for each a in Σ do
    S = ε-Closure(MoveNFA(T,a))
    if S is not in SDFA already then
      Add S to SDFA (as an "unmarked" state)
    endIf
    Set MoveDFA(T,a) to S
  endFor
endWhile
for each S in SDFA do
  if any s∈S is a final state in the NFA then
    Mark S as a final state in the DFA
  endIf
endFor
    
```

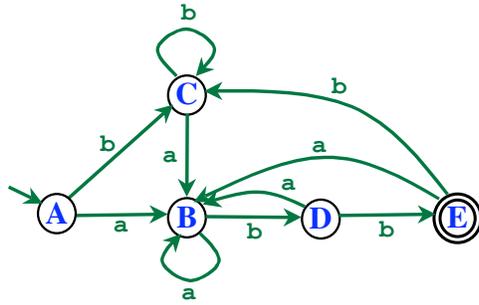
A set of NFA states

Everywhere you could possibly get to on an a

i.e., add an edge to the DFA...

Lexical Analysis - Part 3

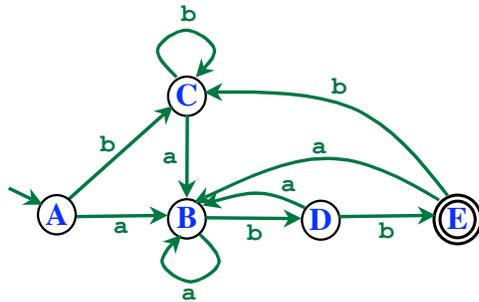
Resulting DFA for $(a|b)^*abb$



Is it minimal?

Lexical Analysis - Part 3

Resulting DFA for $(a|b)^*abb$

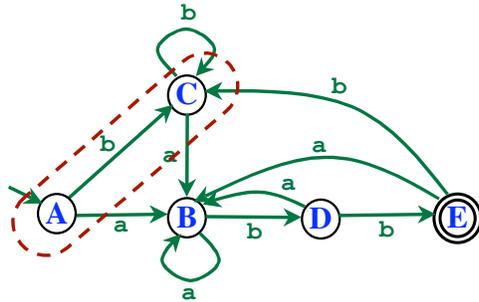


Is it minimal?

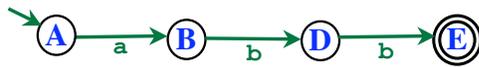


Lexical Analysis - Part 3

Resulting DFA for $(a|b)^*abb$

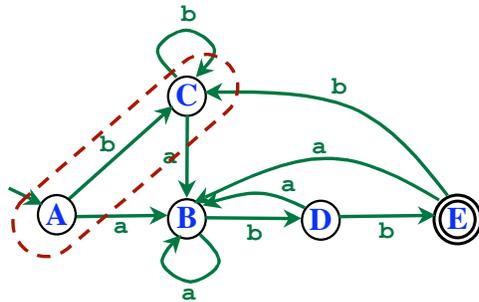


Is it minimal?

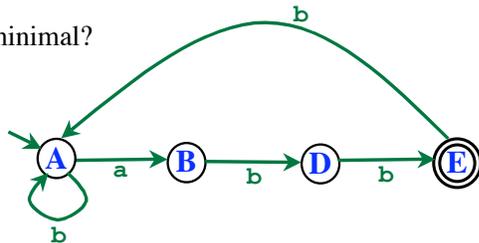


Lexical Analysis - Part 3

Resulting DFA for $(a|b)^*abb$

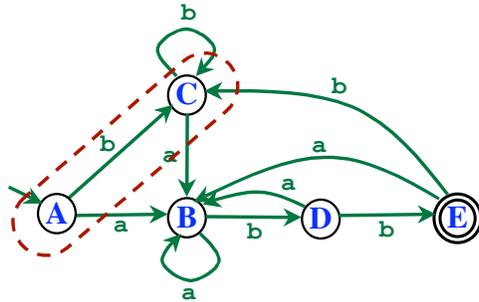


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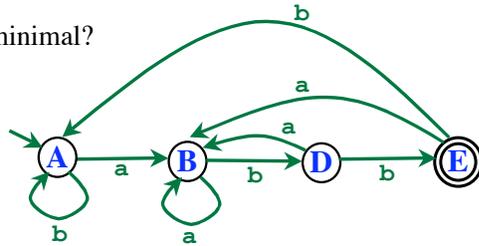


Lexical Analysis - Part 3

Resulting DFA for $(a|b)^*abb$

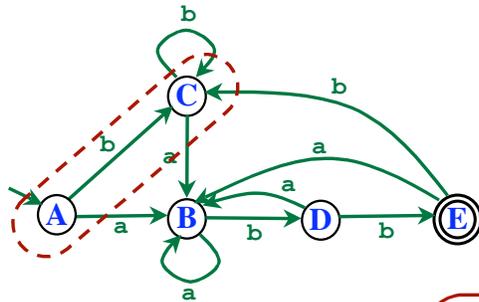


Is it minimal?

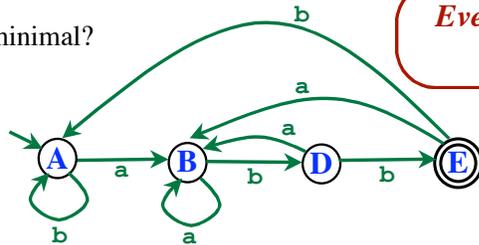


Lexical Analysis - Part 3

Resulting DFA for $(a|b)^*abb$



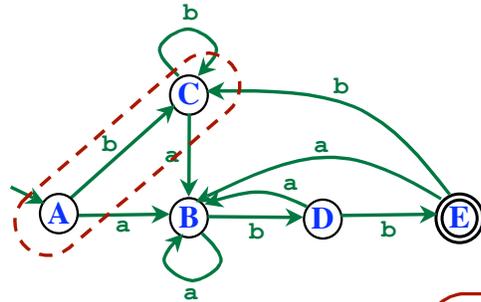
Is it minimal?



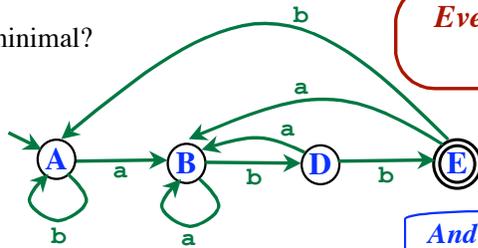
Every Regular Set is recognized by a minimal DFA!

Lexical Analysis - Part 3

Resulting DFA for $(a|b)^*abb$



Is it minimal?



Every Regular Set is recognized by a minimal DFA!

And it is unique, up to renaming of states