

# CS321 Languages and Compiler Design I

## Winter 2012

### Lecture 1

## COURSE GOALS

- Improve understanding of **languages** and **machines**.
- Learn practicalities of **translation**.
- Learn “**anatomy**” of programming languages.
- Apply computer science **theory** to practical problems (using **tools**).
- Do large programming **project**.

# COMPILERS

A **compiler** is a **translator** from “high-level” language to assembly code/object language.

Language L  $\longrightarrow$  **TRANSLATOR**  $\longrightarrow$  Language L'

Examples of translators:

Pascal, C, etc.  $\longrightarrow$  **Compiler**  $\longrightarrow$  Machine Code

Java  $\longrightarrow$  **Compiler**  $\longrightarrow$  Byte Code

Ratfor  $\longrightarrow$  **Preprocessor**  $\longrightarrow$  Fortran

Tex  $\longrightarrow$  **Text Formatter**  $\longrightarrow$  Postscript

SQL  $\longrightarrow$  **DB Optimizer**  $\longrightarrow$  Query plan

We study common features of translators, by building one.

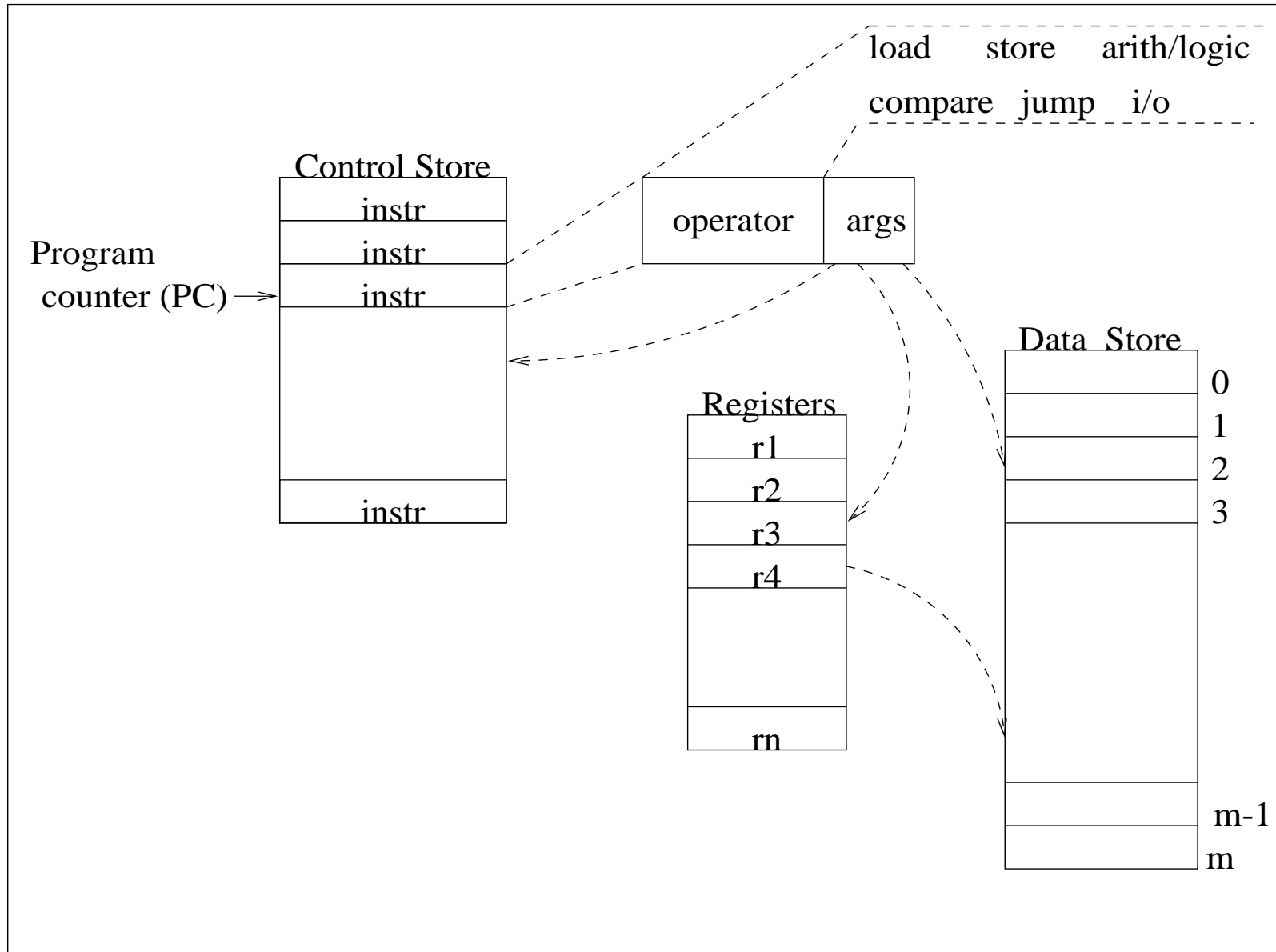
## LANGUAGE DESIGN

We study languages mainly from an **implementor's** viewpoint.

- How do **compilation feasibility** and **runtime efficiency** affect language design?

(There are more “theoretical” approaches to studying programming languages, and there are interesting and useful languages that don't compile easily...)

# "VON NEUMANN" MACHINE



## FEATURES OF LOW-LEVEL CODE

- Sequential control flow + labels + jumps
- Small set of built-in data types and operators (e.g., byte, integer, floating point)
- Flat linear address space.
- Memory hierarchy (registers faster than memory faster than disk).

## “HIGH-LEVEL” LANGUAGES

E.g., Fortran, Pascal, C, Cobol, Java, JavaScript, Python...

Example

```
func rev (a: @real, n:int) {  
  var i := 0;  
  var j := n - 1;  
  while i < j do {  
    var x := a[i];  
    a[i] := a[j];  
    a[j] := x;  
    i := i + 1;  
    j := j - 1  
  }  
}
```

## FEATURES OF HIGH-LEVEL CODE

- Expressions (arithmetic, logical)
- Control structures (loops, conditionals, etc.)
- Type declarations and type checking
- Composite types (arrays, records, etc.)
- Procedures/Functions, with private scope
- **Abstraction** facilities!



## MEETING IN THE MIDDLE

How can we make high-level language run on a Von Neumann machine?

Answer:

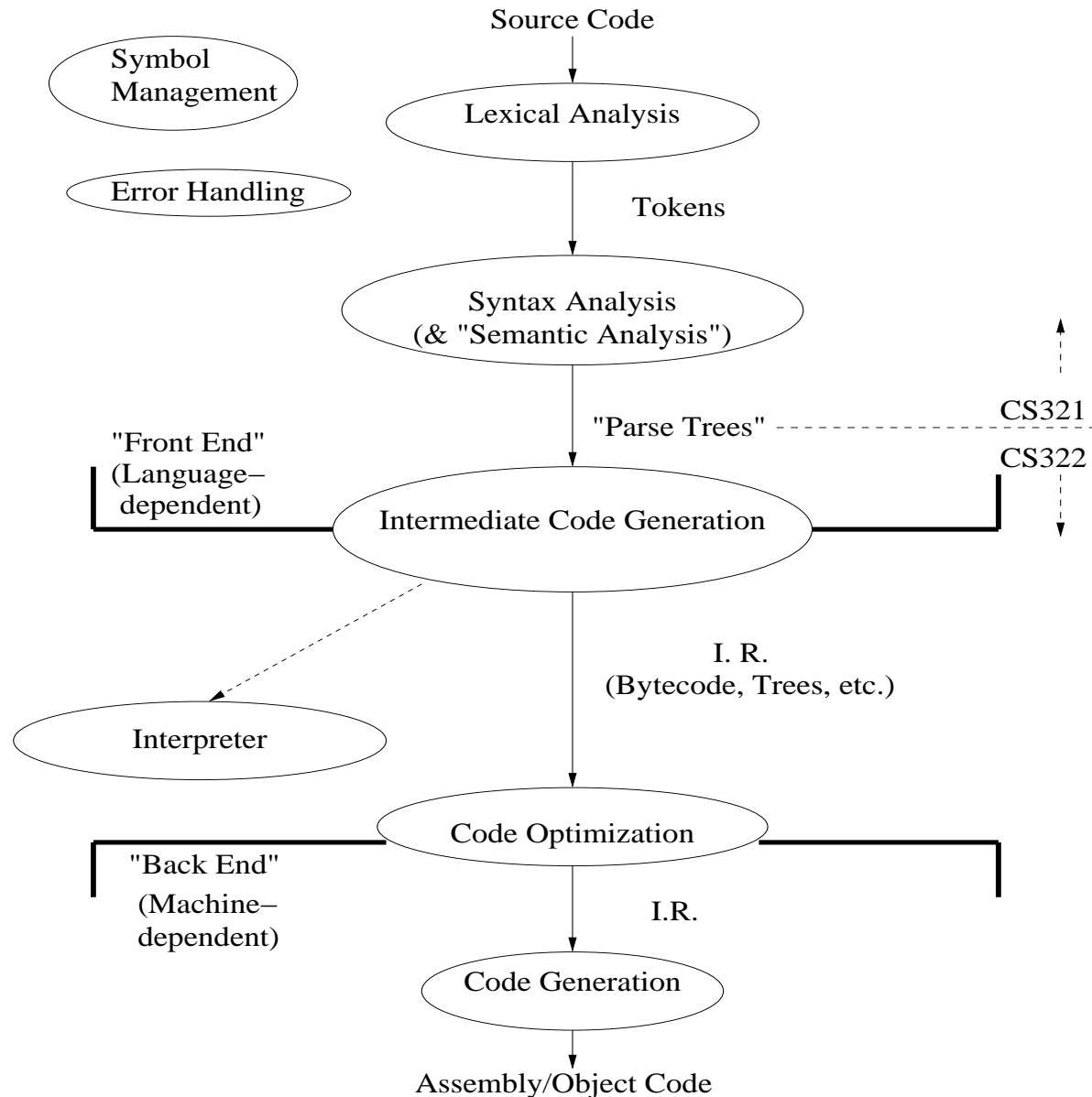
- Translate HLL into lower-level code  
(in traditional compiler, to machine code.)

and/or

- Build a “higher level” virtual machine  
(in traditional interpreter, perhaps a stack machine.)

In practice, we do some of both, even in a compiler, since generated machine code makes use of a runtime library and operating system.

# COMPILER STRUCTURE: WANT SIMPLICITY AND FLEXIBILITY



# FRONT-END EXAMPLE

Source characters: `if (a <= b[i]) a := 4.5 ;`

Lexical Analysis  
"linear"

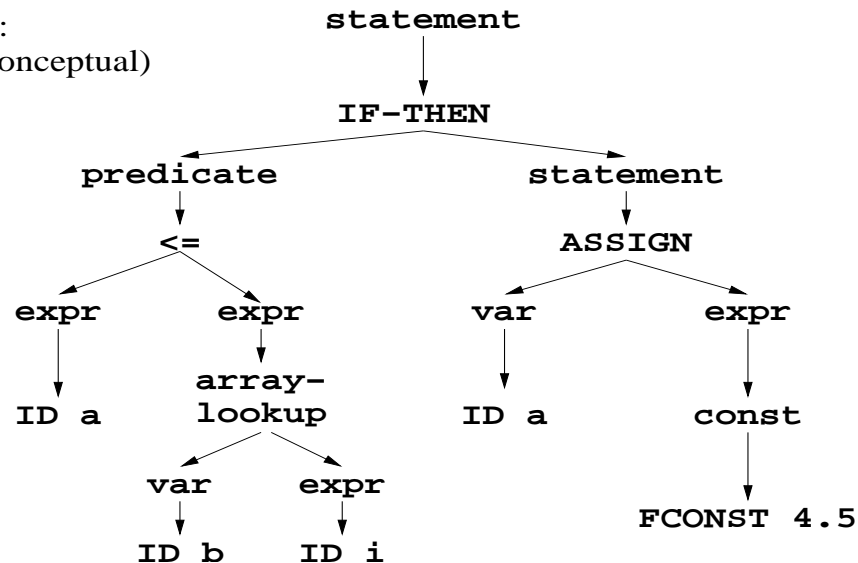
Theory: regular languages, FAs  
Tools: lex, Jflex, etc.

Token stream: `IF '(' (ID a) LE (ID b) '[' (ID i) ']' ')' (ID a) ASSGN (FCONST 4.5) ';'`

Syntax Analysis  
"hierarchical"

Theory: context-free languages,  
PDA's  
Tools: yacc, javaCup, jacc, etc.

Parse tree:  
(real or conceptual)



**Syntax** is easy.

- Well-understood.
- Good theory: regular and context-free languages and automata.
- Good tools, even for complex cases.

**Semantics** are hard.

- Inherently complex.
- Variety of choices:

Informal — Reference Manual

Operational — Definitional interpreter  
(↑ *we will focus here*)

Axiomatic — Logic

Denotational — Mathematical functions  
etc.

- Few tools.