CS 333
Introduction to Operating Systems

Class 1 - Introduction to OS-related Hardware and Software

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About the instructor

- **Instructor - Jonathan Walpole**
  - Professor at PSU since 2004, OGI 1989 - 2004
  - Research Interests: *Operating System Design, Parallel and Distributed Computing Systems*
  - [http://www.cs.pdx.edu/~walpole](http://www.cs.pdx.edu/~walpole)
About CS 333

- **Goals of the class**
  - understand the basic concepts of operating systems
    - *designing & building* operating systems, not *using* them!
  - gain some practical experience so it's not just words!

- **Expectations**
  - reading assignments should be read *before* class
  - active participation in class discussions
  - no cheating!
Grading

- Exams
  - Mid-term - 25%
  - Final - 25%

- Coursework
  - Project - 50%
Text books


“The BLITZ System” by Harry Porter
The project

- You will read, understand and write real operating system code!
- We will be using the BLITZ system, written by Harry Porter
- About BLITZ
  - CPU emulator, assembler, high-level language, operating system, and debugging environment
  - Simple enough to understand in detail how everything works!
  - Realistic enough to understand in detail how everything works!
  - Runs on the departmental Sun machines (cs.pdx.edu), plus Macs and x86/Linux
Administrative

- **Class web site**
  - [www.cs.pdx.edu/~walpole/class/cs333/winter2013/home.html](http://www.cs.pdx.edu/~walpole/class/cs333/winter2013/home.html)
  - Find my website from the faculty listing on the department website. Follow *teaching* link to current classes

- **Class mailing list**
  - [https://mailhost.cecs.pdx.edu/cgi-bin/mailman/listinfo/cs333](https://mailhost.cecs.pdx.edu/cgi-bin/mailman/listinfo/cs333)

- **Project 0**
  - read the class web site
  - join the class mailing list

- **Project 1**
  - due next week!
  - see class web site for project assignments
Class 1 - Introduction to OS-related Hardware and Software
Overview

- What is an Operating System?
- A review of OS-related hardware
What is an operating system?

- Operating system -- “a program that controls the execution of application programs and implements an interface between the user of a computer and the computer hardware”
  - Narrow view of a computer and OS
    - Traditional computer with applications running on it (e.g. PCs, Workstations, Servers)
  - Broad view of a computer and OS
    - Anything that needs to manage resources (e.g. router OS, embedded system, cell phone OS …)
Two key OS functions

- **Abstract Machine**
  - Hides complex details of the underlying hardware
  - Provides common API to applications and services
  - Simplifies application writing

- **Resource Manager**
  - Controls accesses to *shared* resources
    - CPU, memory, disks, network, ...
  - Allows for global policies to be implemented
Why is abstraction important?

- Without OSs and abstract interfaces, application writers must program all device access directly:
  - load device command codes into device registers
  - handle initialization, recalibration, sensing, timing etc for physical devices
  - understand physical characteristics and layout
  - control motors
  - interpret return codes ... etc

- Applications suffer severe code bloat!
  - very complicated maintenance and upgrading
  - no portability
  - writing this code once, and sharing it, is how OS began!
Providing abstraction via system calls
Providing abstraction via system calls

System Calls: `read()`, `open()`, `write()`, `mkdir()`, `kill()` ...
OS as a resource manager

- Allocating resources to applications across space and time
  - time sharing a resource (scheduling)
  - space sharing a resource (allocation)

- Making efficient use of limited resources
  - improving utilization
  - minimizing overhead
  - improving throughput/good put

- Protecting applications from each other
  - enforcement of boundaries
Problems an OS must solve

- Time sharing the CPU among applications
- Space sharing the memory among applications
- Space sharing the disk among users
- Time sharing access to the disk
- Time sharing access to the network
More problems an OS must solve

- **Protection**
  - of applications from each other
  - of user data from other users
  - of hardware/devices
  - of the OS itself!

- The OS is just a program! It needs help from the hardware to accomplish these tasks!
  - When an application is running, the OS is not running!
  - When the OS is not running, it can’t do anything!
Overview

- **What is an Operating System?**
- **A review of OS-related hardware**
Instruction sets

- A CPU's instruction set defines what it can do
  - different for different CPU architectures
  - all have load and store instructions for moving items between memory and registers
    - Load a word located at an address in memory into a register
    - Store the contents of a register to a word located at an address in memory
  - many instructions for comparing and combining values in registers and putting result into a register

- Look at the Blitz instruction set which is similar to a SUN SPARC instruction set
Basic anatomy on a CPU

- Program Counter (PC)
Basic anatomy on a CPU

- **Program Counter (PC)**
  - Holds the memory address of the next instruction
Basic anatomy on a CPU

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Basic anatomy on a CPU

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- **Arithmetic and Logic Unit (ALU)**
Basic anatomy on a CPU

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  - Holds the memory address of the next instruction

- **Instruction Register (IR)**
  - Holds the instruction currently being executed

- **General Registers (Reg. 1..n)**
  - Hold variables and temporary results

- **Arithmetic and Logic Unit (ALU)**
  - Performs arithmetic functions and logic operations
Basic anatomy on a CPU

- Stack Pointer (SP)
Basic anatomy on a CPU

- Stack Pointer (SP)
  - holds memory address of a stack with a frame for each active procedure’s parameters & local variables
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- Processor Status Word (PSW)
Basic anatomy on a CPU

- **Stack Pointer (SP)**
  - holds memory address of a stack with a frame for each active procedure’s parameters & local variables

- **Processor Status Word (PSW)**
  - contains various control bits including the *mode bit* which determines whether privileged instructions can be executed at this time
Basic anatomy on a CPU

- **Stack Pointer (SP)**
  - holds memory address of a stack with a frame for each active procedure’s parameters & local variables

- **Processor Status Word (PSW)**
  - contains various control bits including the mode bit which determines whether privileged instructions can be executed

- **Memory Address Register (MAR)**
  - contains address of memory to be loaded from/stored to

- **Memory Data Register (MDR)**
  - contains memory data loaded or to be stored
Program execution

- The Fetch/Decode/Execute cycle
  - fetch next instruction pointed to by PC
  - decode it to find its type and operands
  - execute it
  - repeat

- At a fundamental level, fetch/decode/execute is all a CPU does, regardless of which program it is executing
Fetch/decode/execute cycle

CPU

PC
IR
Reg. 1
...
Reg. n
ALU

MAR
MDR

Memory
Fetch/decode/execute cycle

While (1) {
    Fetch instruction from memory
    Execute instruction
        (Get other operands if necessary)
    Store result
}
Fetch/decode/execute cycle

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  PC  IR
  Reg. 1
  ...
  Reg. n

ALU

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MDR
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CPU

PC  IR

MAR  MDR

ALU

Reg. 1

... 

Reg. n

Memory
Fetch/decode/execute cycle

While (1) {
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Fetch/decode/execute cycle

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  Fetch instruction from memory
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}
The OS is just a program!

- The OS is a sequence of instructions that the CPU will fetch/decode/execute
  - How can the OS cause application programs to run?
  - How can the OS switch the CPU to run a different application and later resume the first one?
  - How can the OS maintain control?
  - In what ways can application code try to seize control indefinitely (ie. cheat)?
  - And how can the OS prevent such cheating?
  - How can applications programs cause the OS to run?
How can the OS invoke an application?
How can the OS invoke an application?

- Somehow, the OS must load the address of the application’s starting instruction into the PC

  - The computer boots and begins running the OS
    - OS code must be loaded into memory somehow
    - fetch/decode/execute OS instructions
    - OS requests user input to identify application “file”
    - OS loads application file (executable) into memory
    - OS loads the memory address of the application’s starting instruction into the PC
    - CPU fetches/decodes/executes the application's instructions
How can OS guarantee to regain control?

- What if a running application doesn’t make a system call and hence hogs the CPU?
  - OS needs interrupts from a timer device!
  - OS must register a future timer interrupt before it hands control of the CPU over to an application
  - When the timer interrupt goes off the interrupt hardware jumps control back into the OS at a pre-specified location called an interrupt handler
  - The interrupt handler is just a program (part of the OS)
  - The address of the interrupt handler’s first instruction is placed in the PC by the interrupt h/w
What if the application tries to cheat?

- What stops the running application from disabling the future timer interrupt so that the OS can not take control back from it?
  - Disabling interrupts must be a privileged instruction which is not executable by applications
  - The CPU knows whether or not to execute privileged instructions based on the value of the mode bit in the PSW!

- Privileged instructions can only be executed when the mode bit is set
  - disabling interrupts
  - setting the mode bit!
  - Attempted execution in non-privileged mode generally causes an interrupt (trap) to occur
What other ways are there to cheat?

- What stops the running application from modifying the OS?
  - Specifically, what stops it from modifying the timer interrupt handler to jump control back to the application?
What other ways are there to cheat?

- What stops the running application from modifying the OS?
  - Memory protection!
  - Memory protection instructions must be privileged
  - They can only be executed with the mode bit set ...

- Why must the OS clear the mode bit before it hands control to an application?
How can applications invoke the OS?

- Why not just set PC to an OS instruction address and transfer control that way?
How can applications invoke the OS?

- Special instruction causes a trap / interrupt
- Trap instruction changes PC to point to a predetermined OS entry point instruction and simultaneously sets the mode bit
  - application calls a library procedure that includes the appropriate trap instruction
  - fetch/decode/execute cycle begins at a pre-specified OS entry point called a system call
  - CPU is now running in privileged mode
- Traps, like interrupts, are hardware events, but they are caused by the executing program rather than a device external to the CPU
How can the OS switch to a new application?

- To suspend execution of an application simply capture its memory state and processor state
  - preserve all the memory values of this application
  - copy values of all CPU registers into a data structure which is saved in memory
  - restarting the application from the same point just requires reloading the register values
Recap

- Why do we need a timer device?
- Why do we need an interrupt mechanism?
- Why do we need privileged instructions?
- Why are system calls different to procedure calls?
- How are system calls different to interrupts?
- Why is memory protection necessary?
- How can the OS switch from one application to another?
What to do before next class

- Reading for today’s class
- Reading for class 2
- Assignment 0 - read class web page and join class email list
- Start project 1 - Introduction to BLITZ