

# How FIFO is Your Concurrent FIFO Queue?

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# Strict vs. Relaxed FIFO Queues

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relaxed FIFO queue implementations

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out-of-order treatment  
of queue elements  
possible

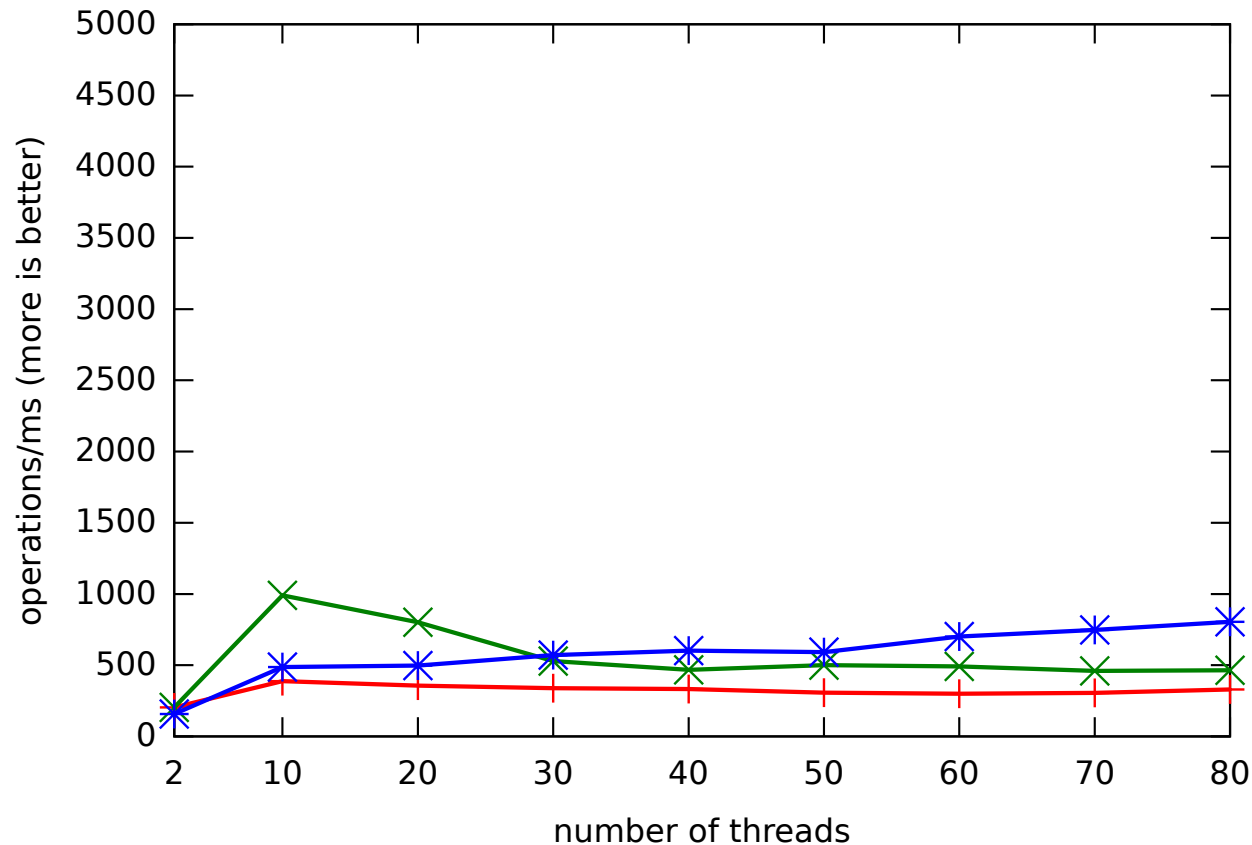
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LB +  
MS x  
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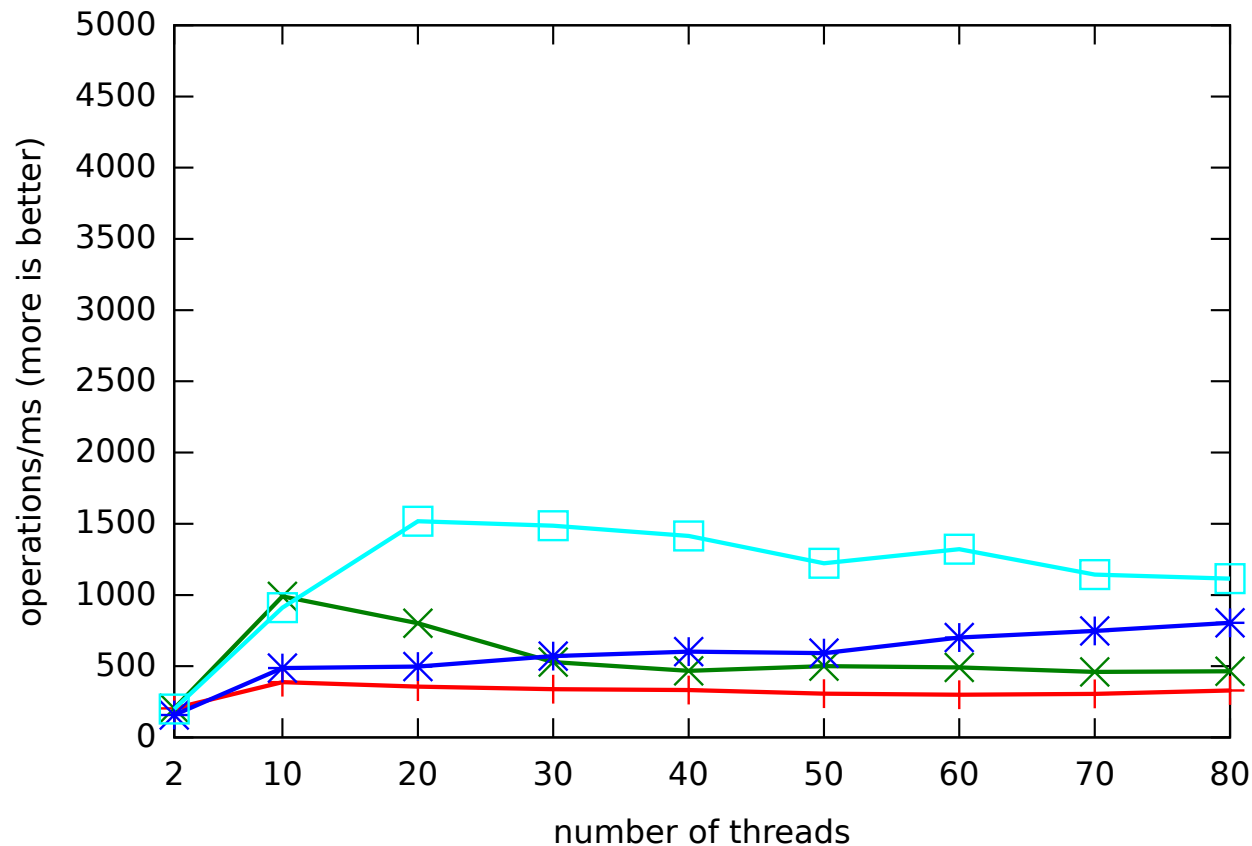
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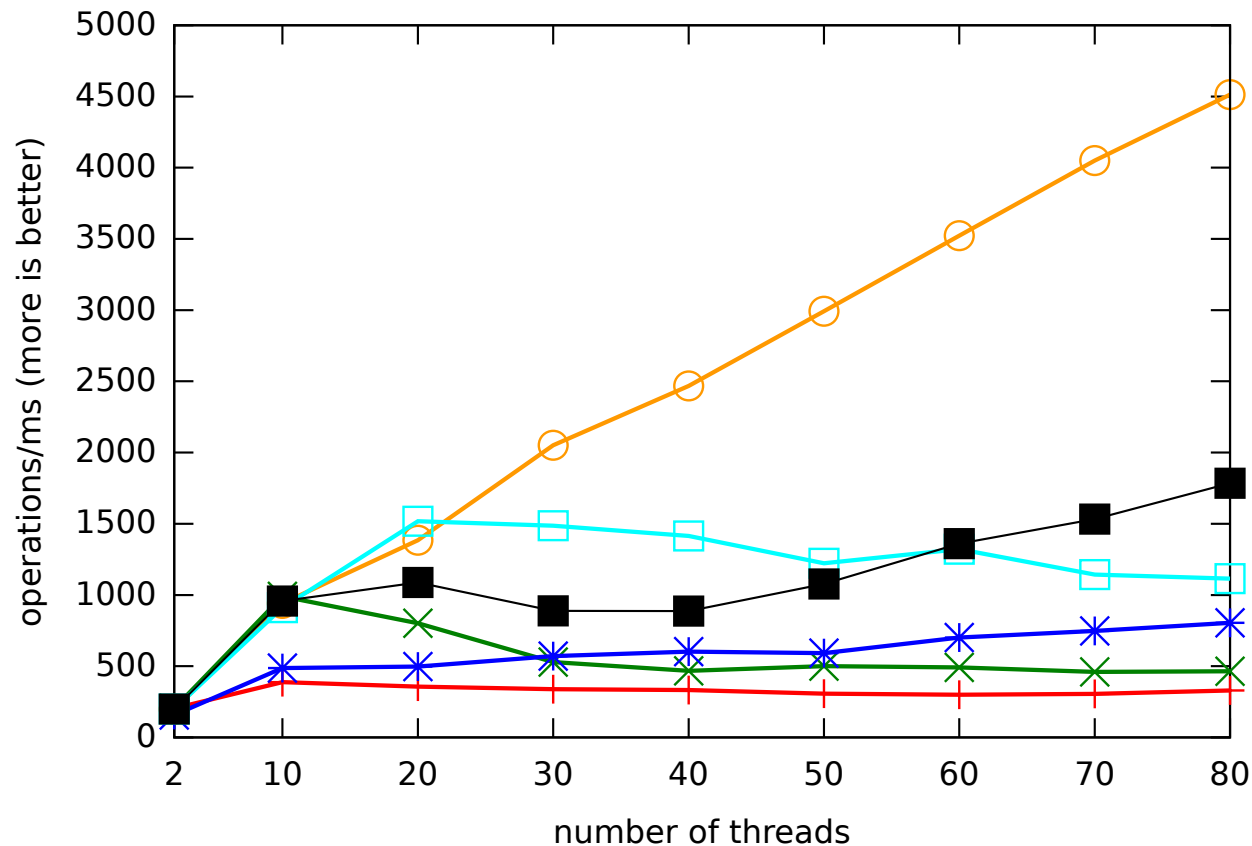
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2RA Scal (p=80) —○—



# How FIFO are Relaxed FIFO Queues?

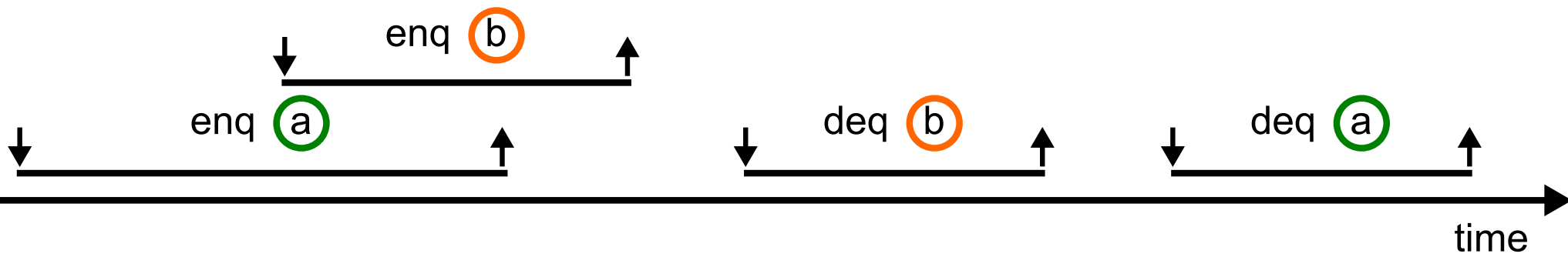
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  - ◆ No applications for relaxed FIFO queues.

# How FIFO are Relaxed FIFO Queues?

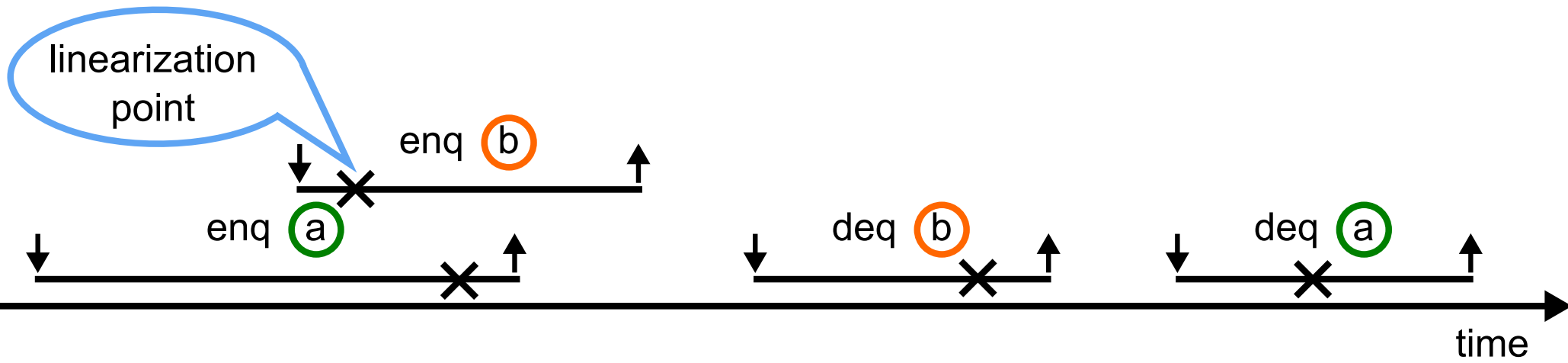
- ▶ Some people say relaxed FIFO queues are not enough FIFO.
  - ◆ No applications for relaxed FIFO queues.

**We say relaxed FIFO queue implementations can be even more FIFO than strict FIFO queue implementations.**

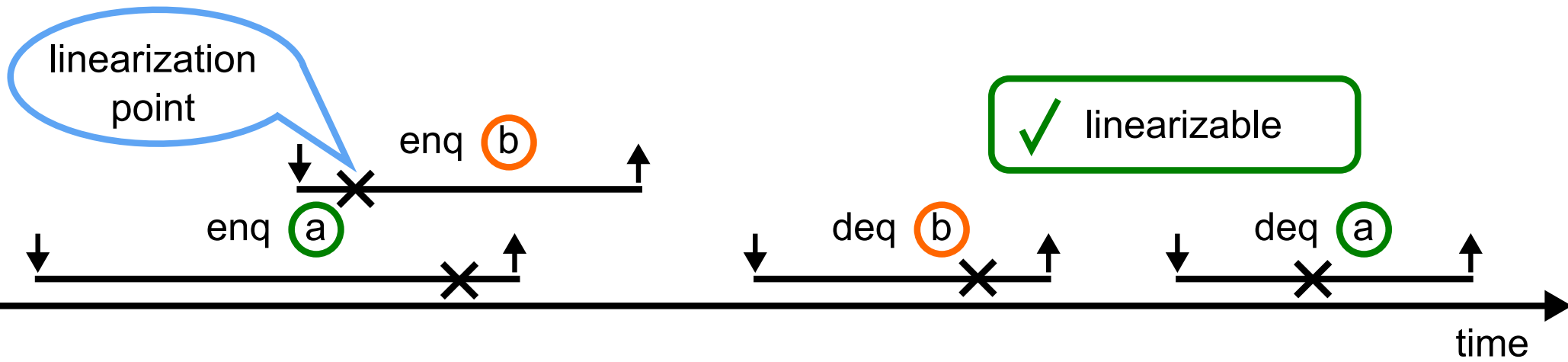
# Example



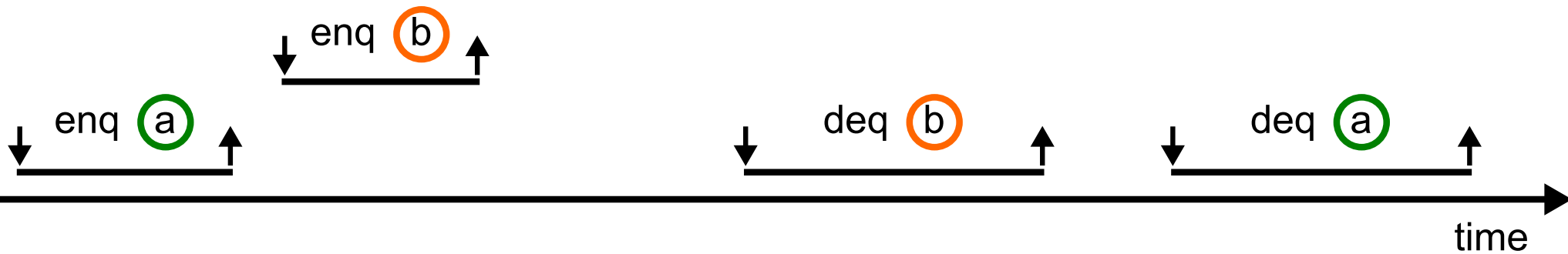
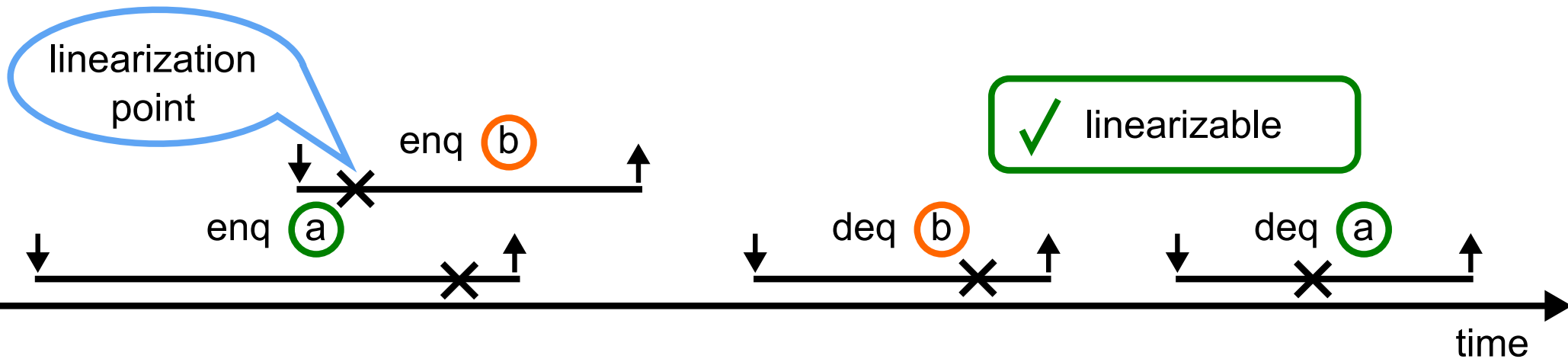
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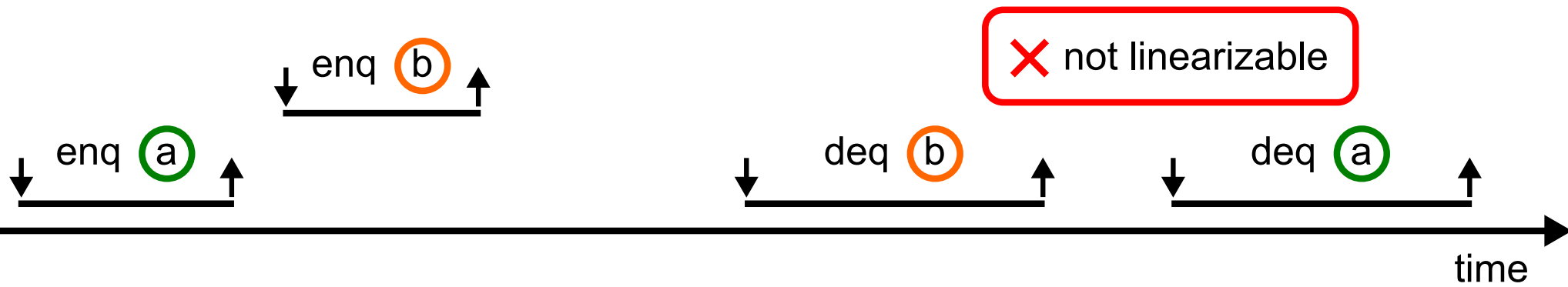
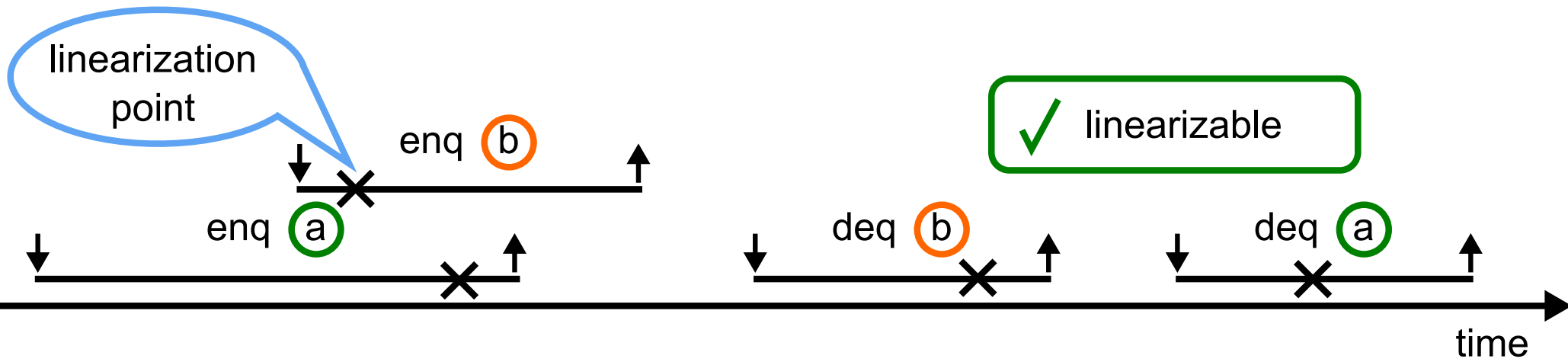
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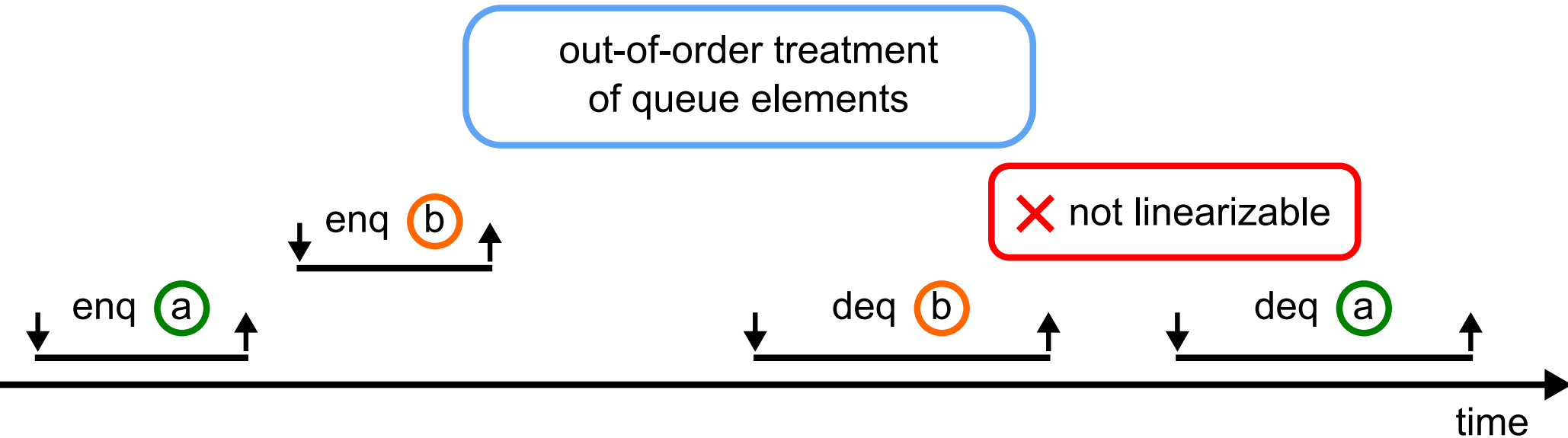
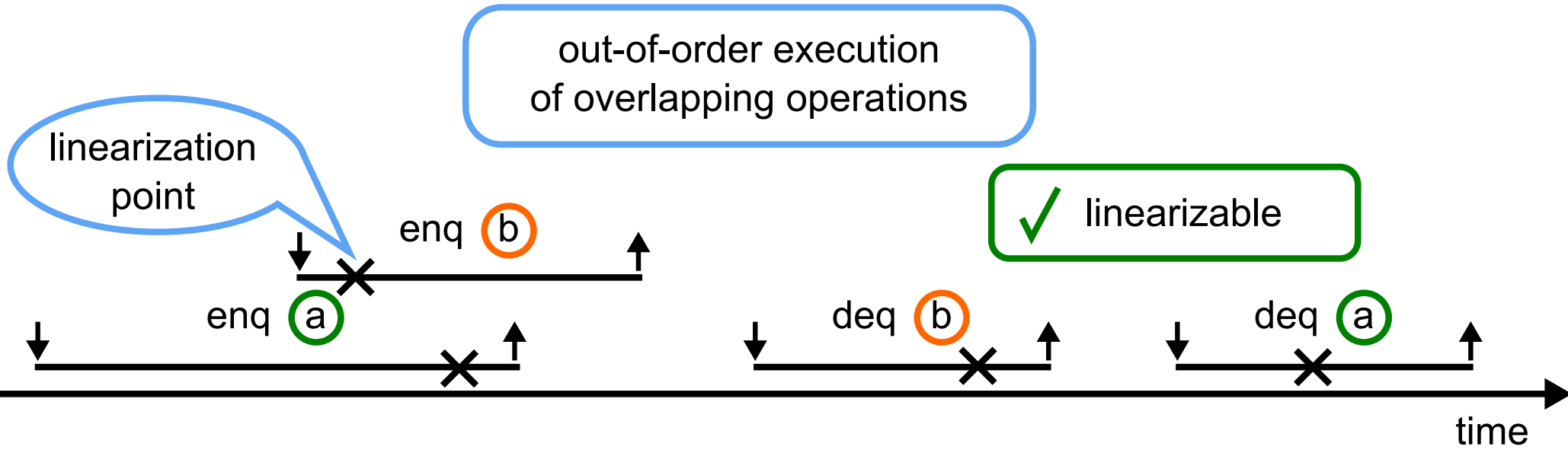
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# Key Idea

- 1.) Record concurrent histories of various FIFO queue implementations.
- 2.) Analyze these concurrent histories using only the invocation times of operations.

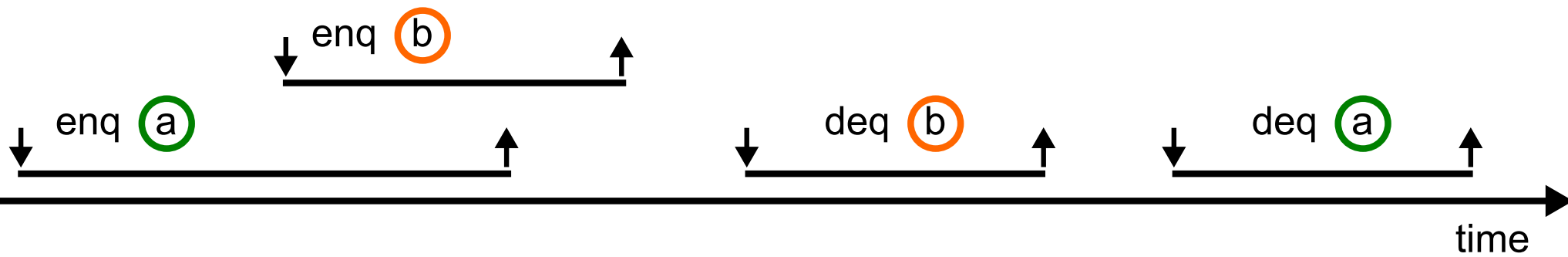
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  - ▶ Independent of the execution time of operations

# Element-Fairness

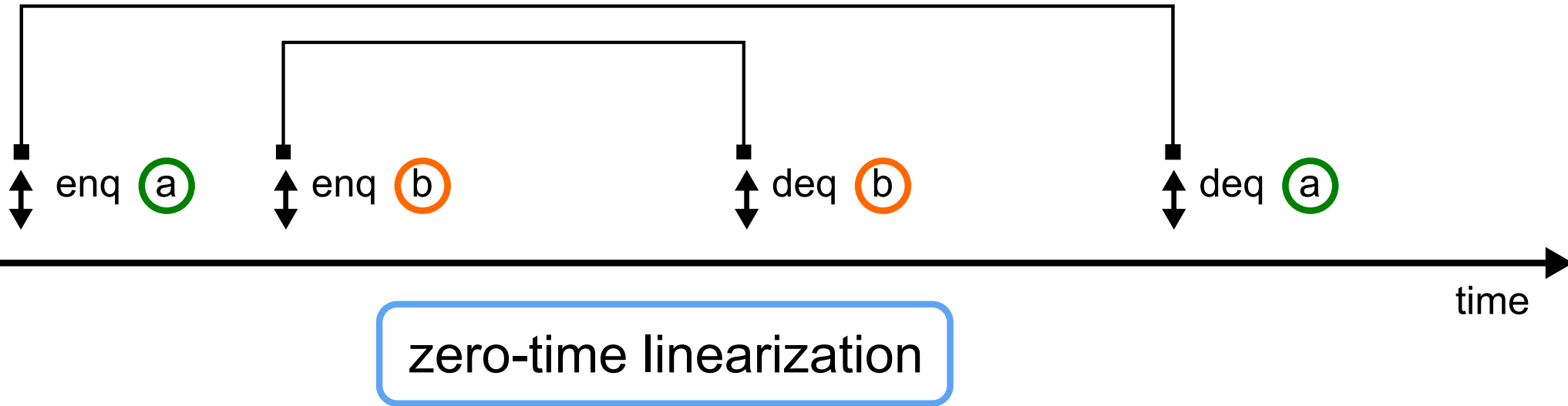


# Element-Fairness



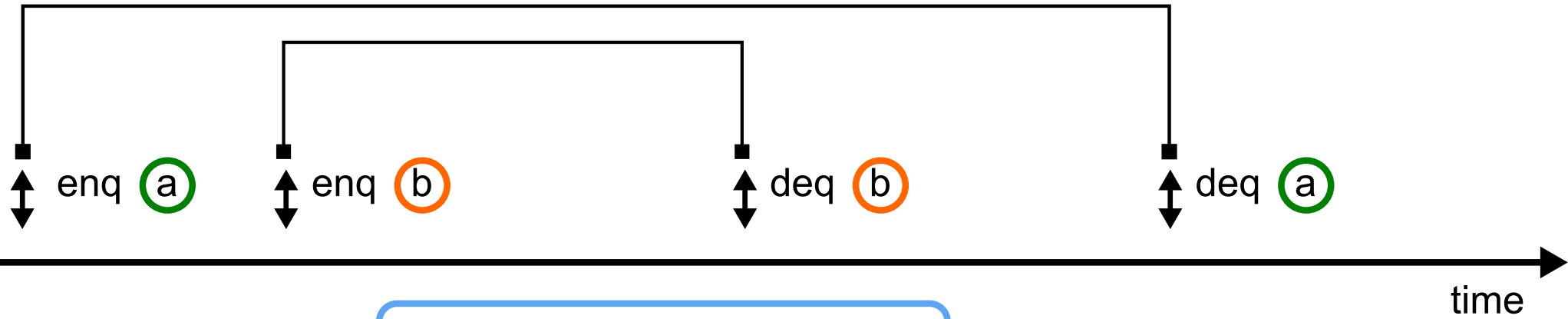
# Element-Fairness

element **(b)** overtakes element **(a)**



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zero-time linearization

Definition

element-fairness =  
number of overtakings in the zero-time linearization

# Experiments

all threads do in parallel

```
for 10.000 iterations  
{  
    enqueue unique element  
  
    dequeue element  
  
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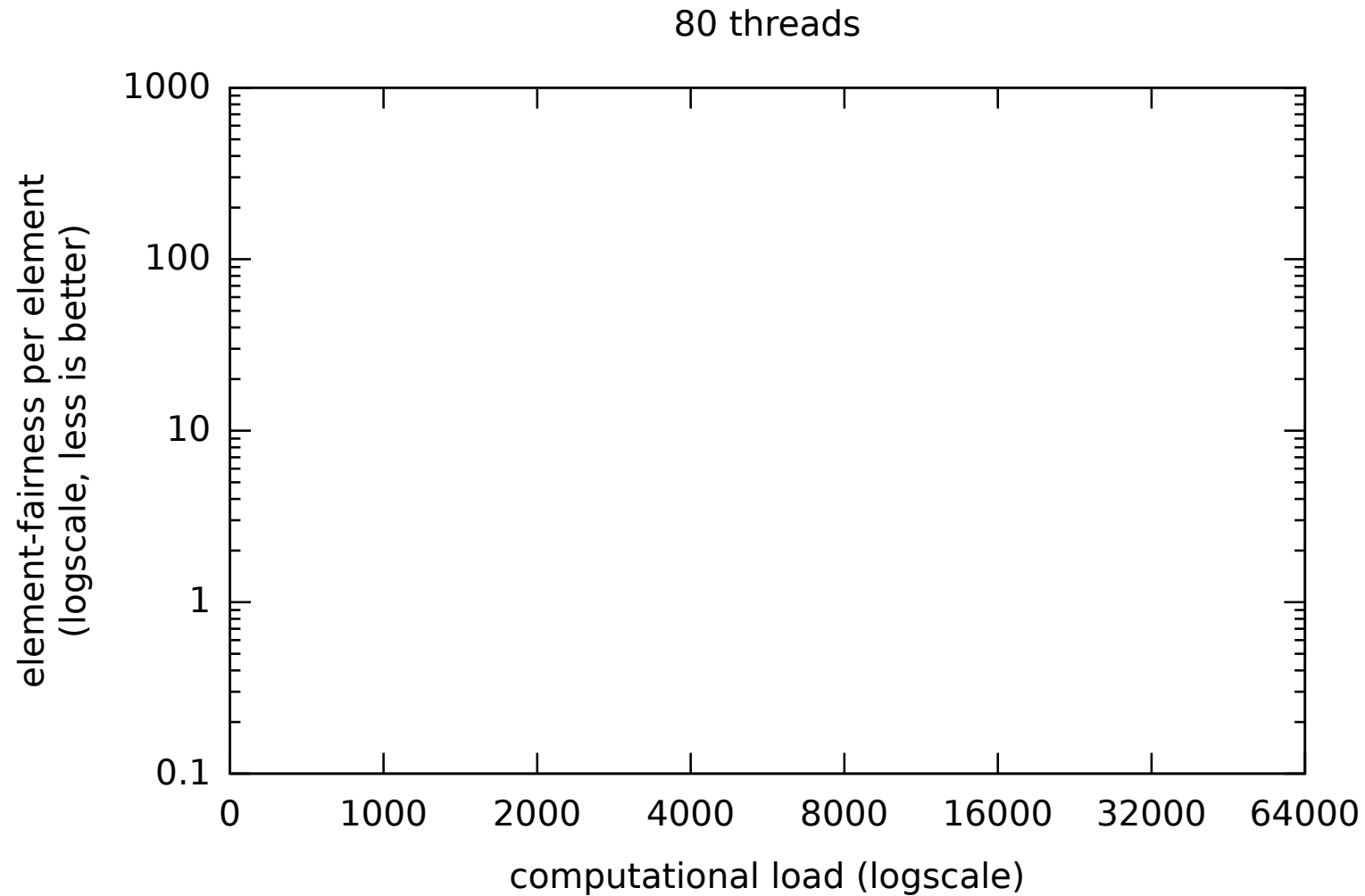
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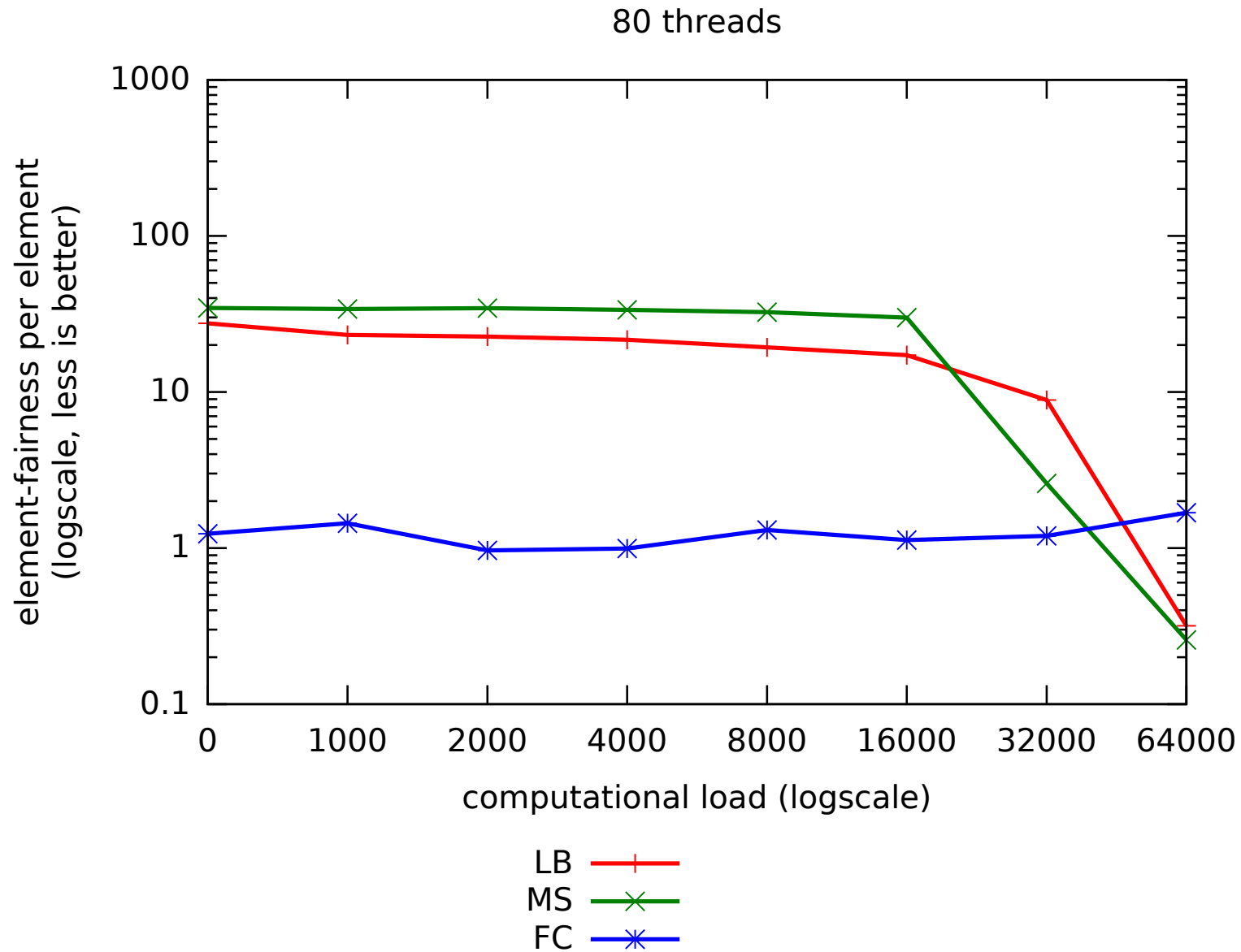
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- ▶ No dequeues in the first 200 iterations to avoid empty checks
- ▶ No enqueues in the last 200 iterations to empty the queue.

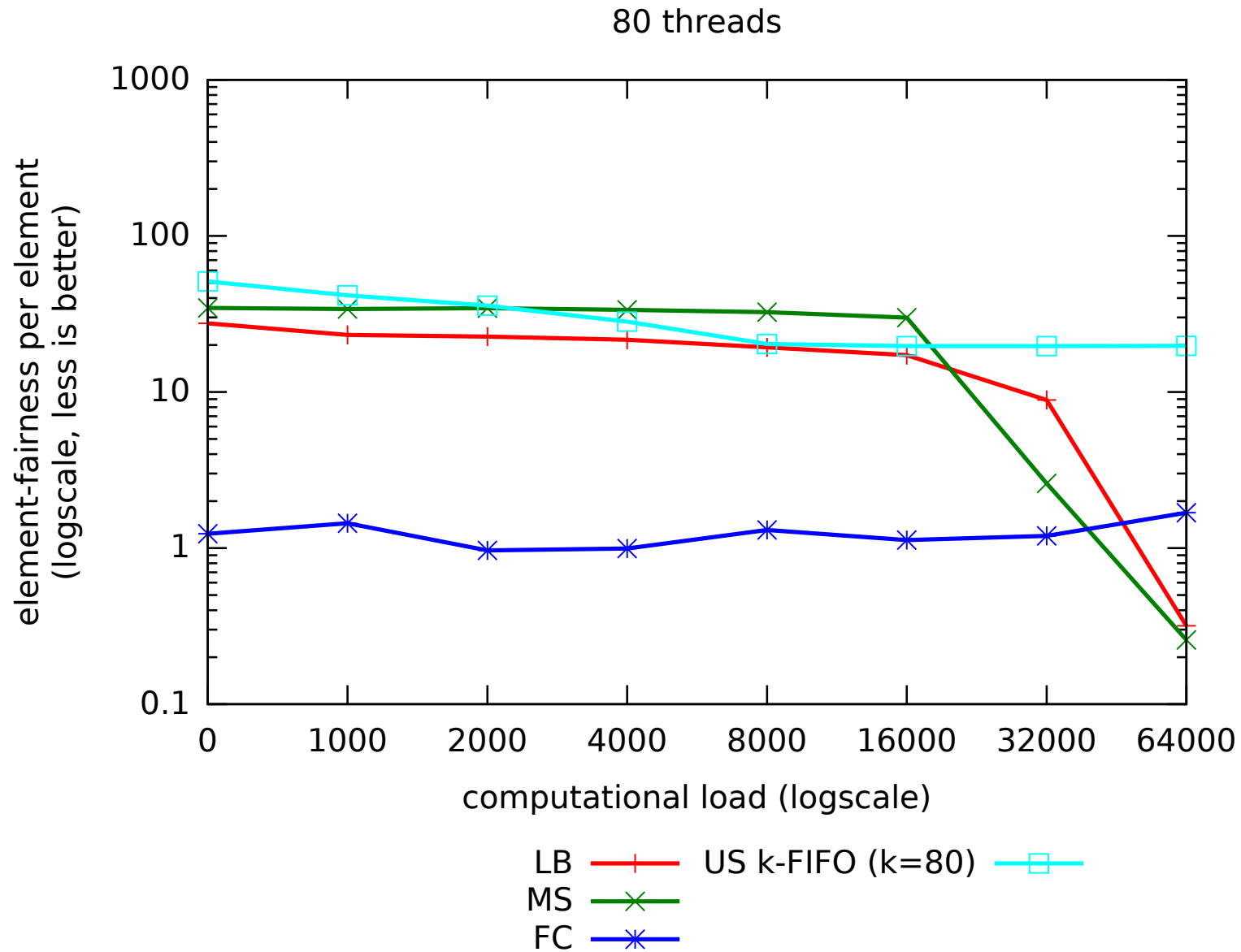
# Element-Fairness per Element



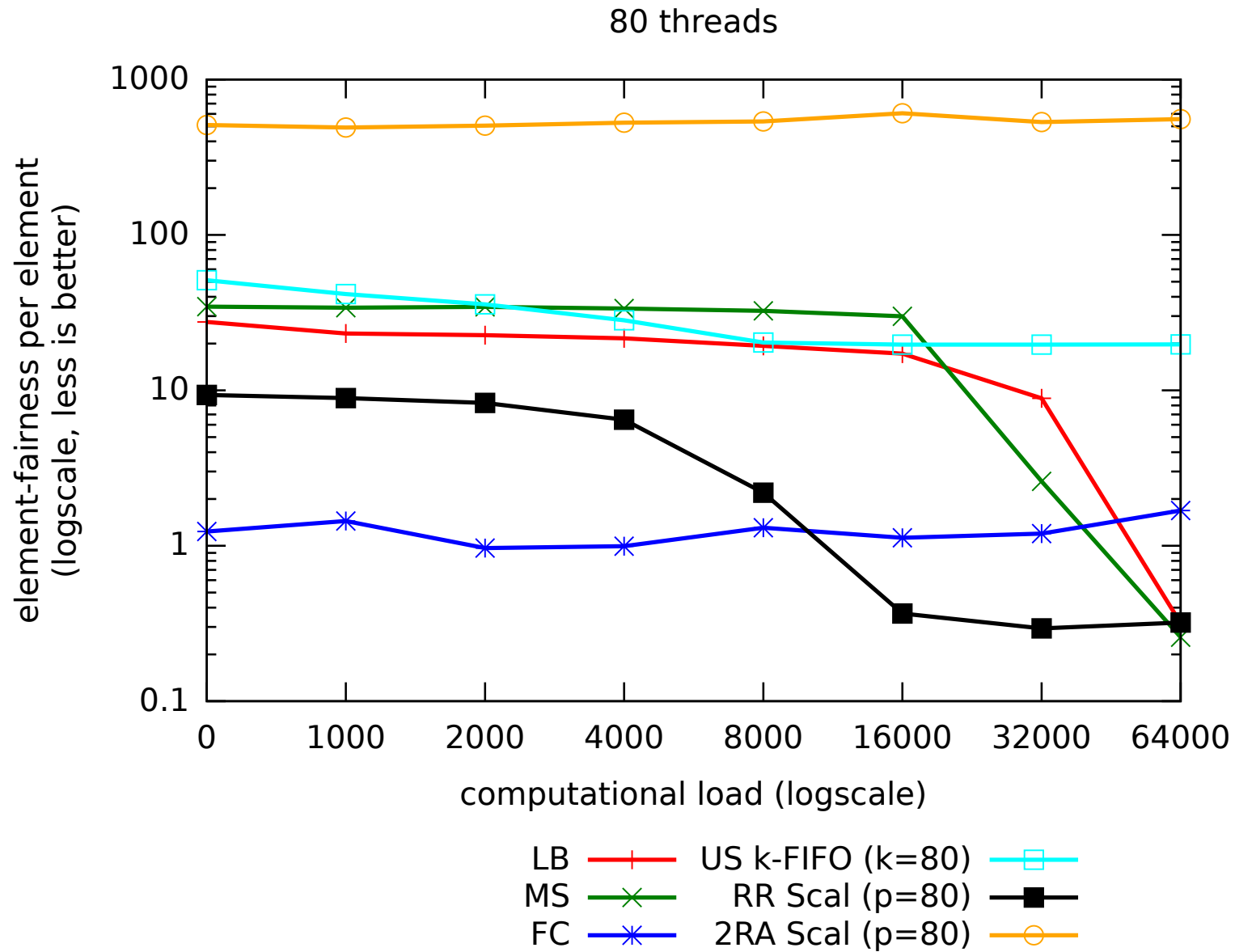
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# Operation-Fairness

- ▶ Measure the out-of order execution of single operations

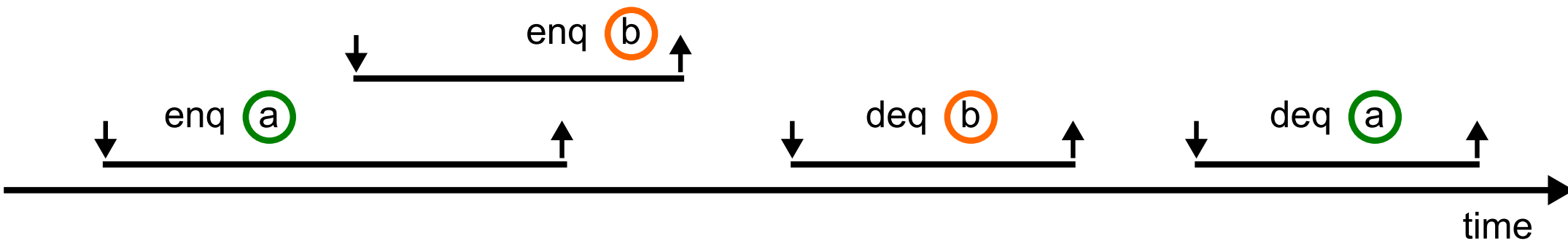
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- ▶ Measure the out-of order execution of single operations
- ▶ Observation: Linearization points induce a strict order on the queue operations



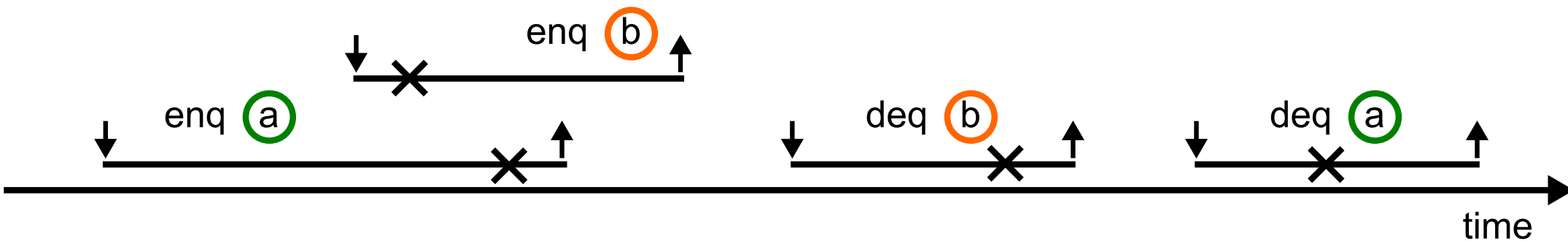
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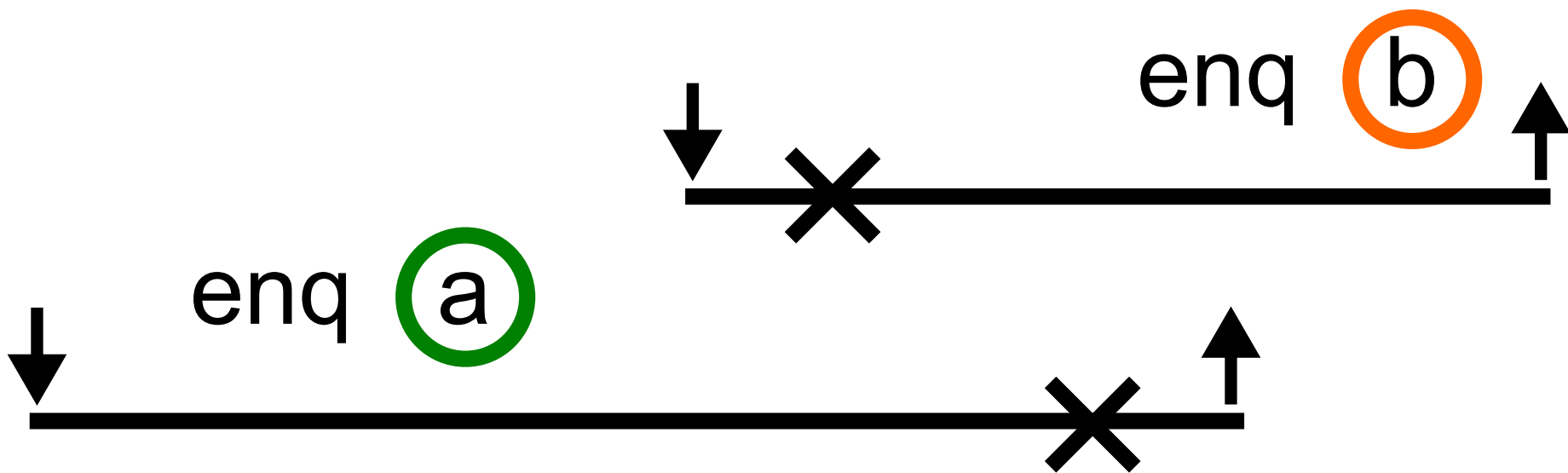
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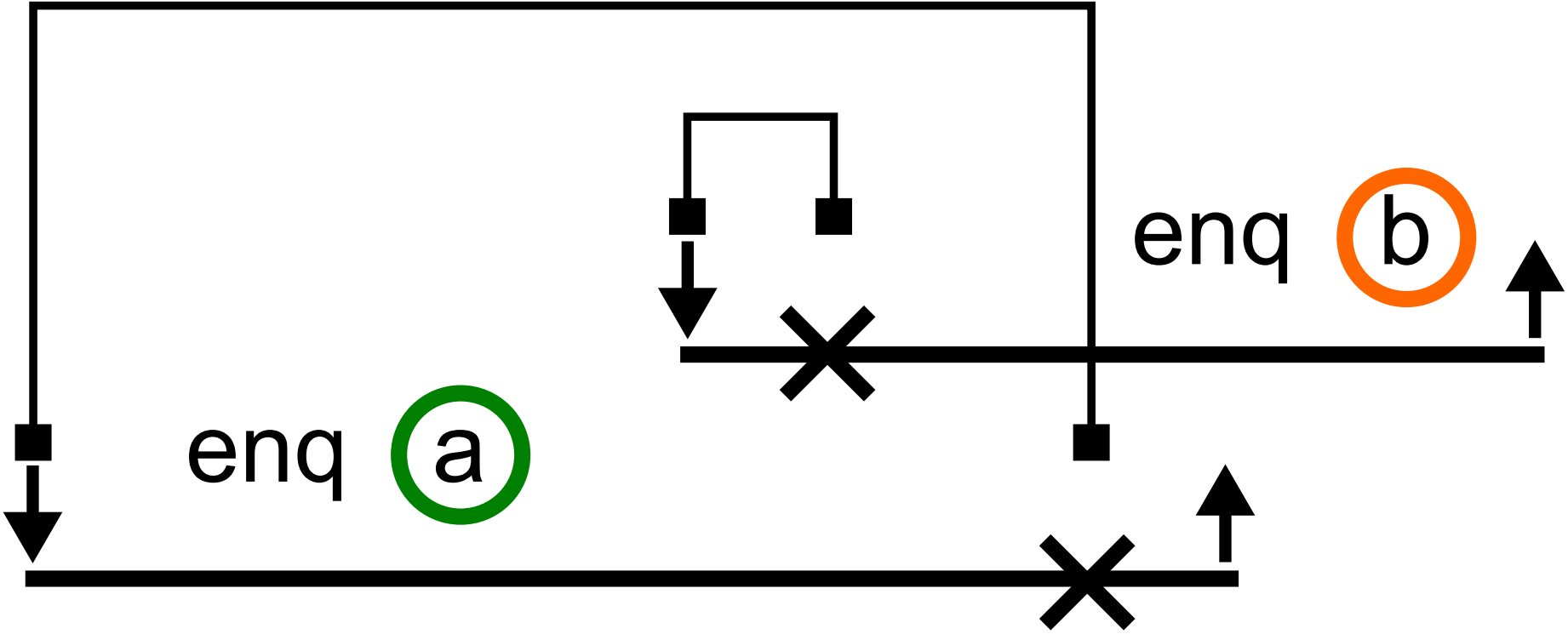
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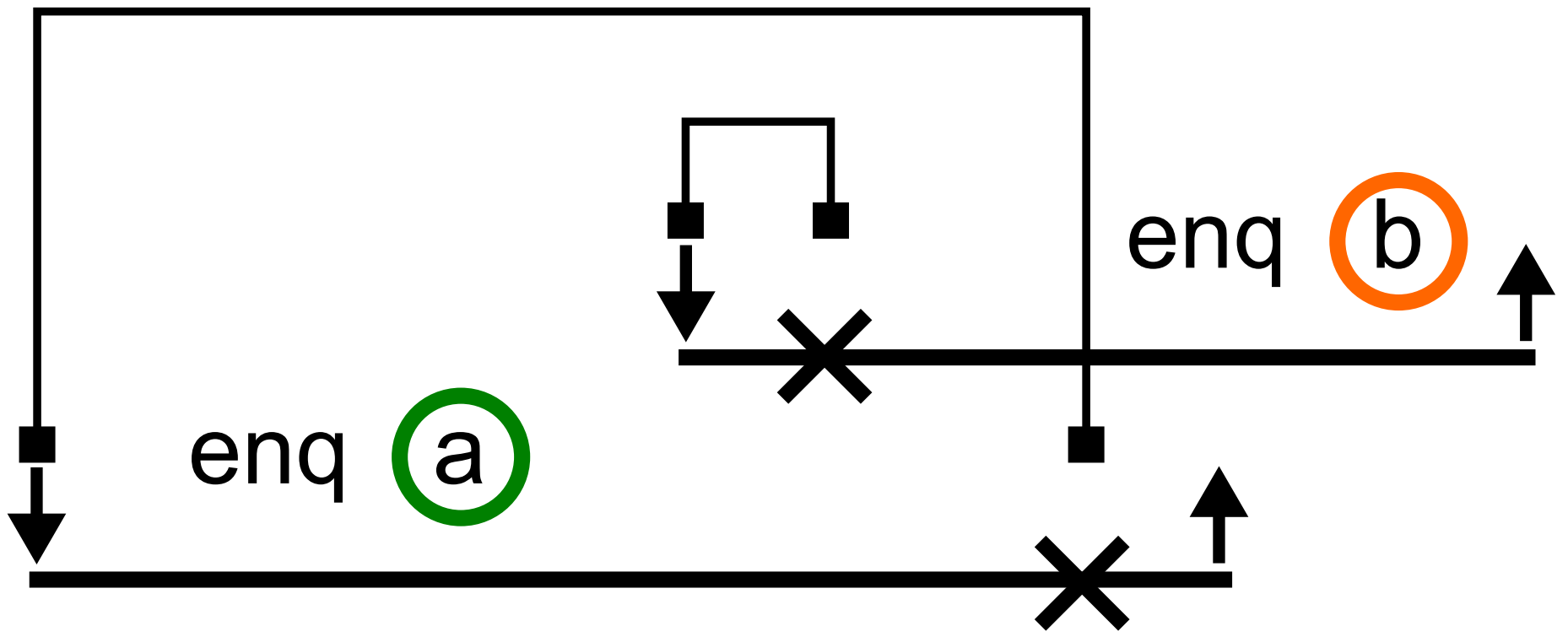




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Definition

operation-fairness =  
number of overtakings in a concurrent history

# Operation-Age and Operation-Lateness

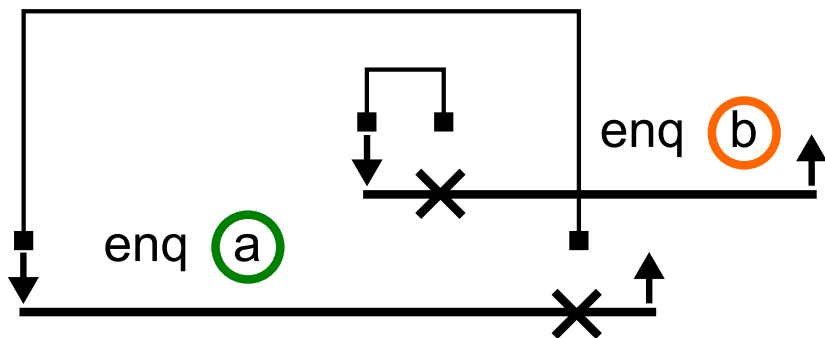
## Definition

$\text{age}(\text{op}) =$   
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$\text{age}(\text{enq } \textcircled{a}) = 0$

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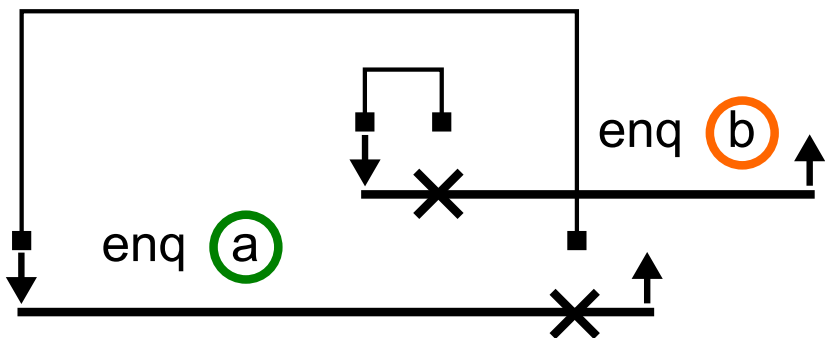
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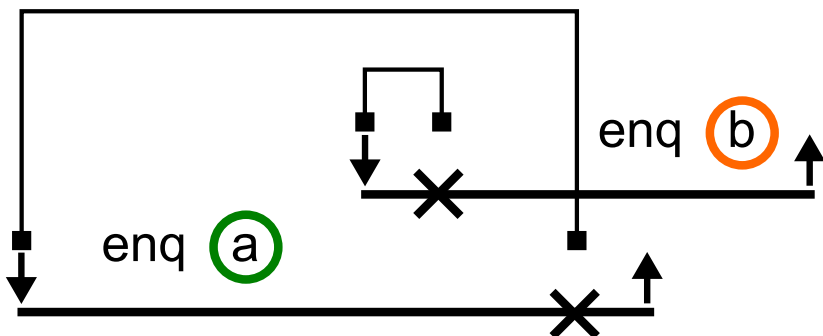
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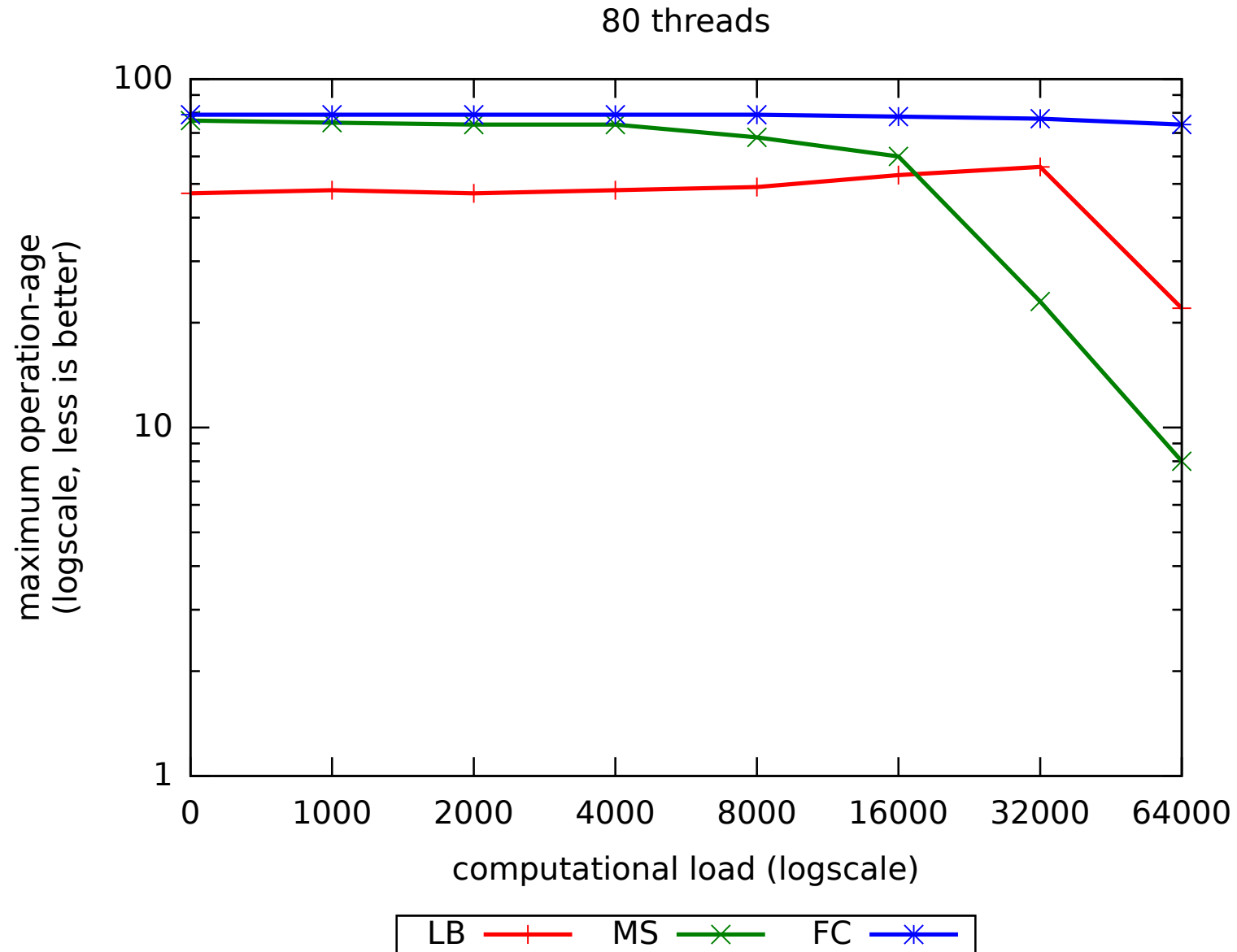
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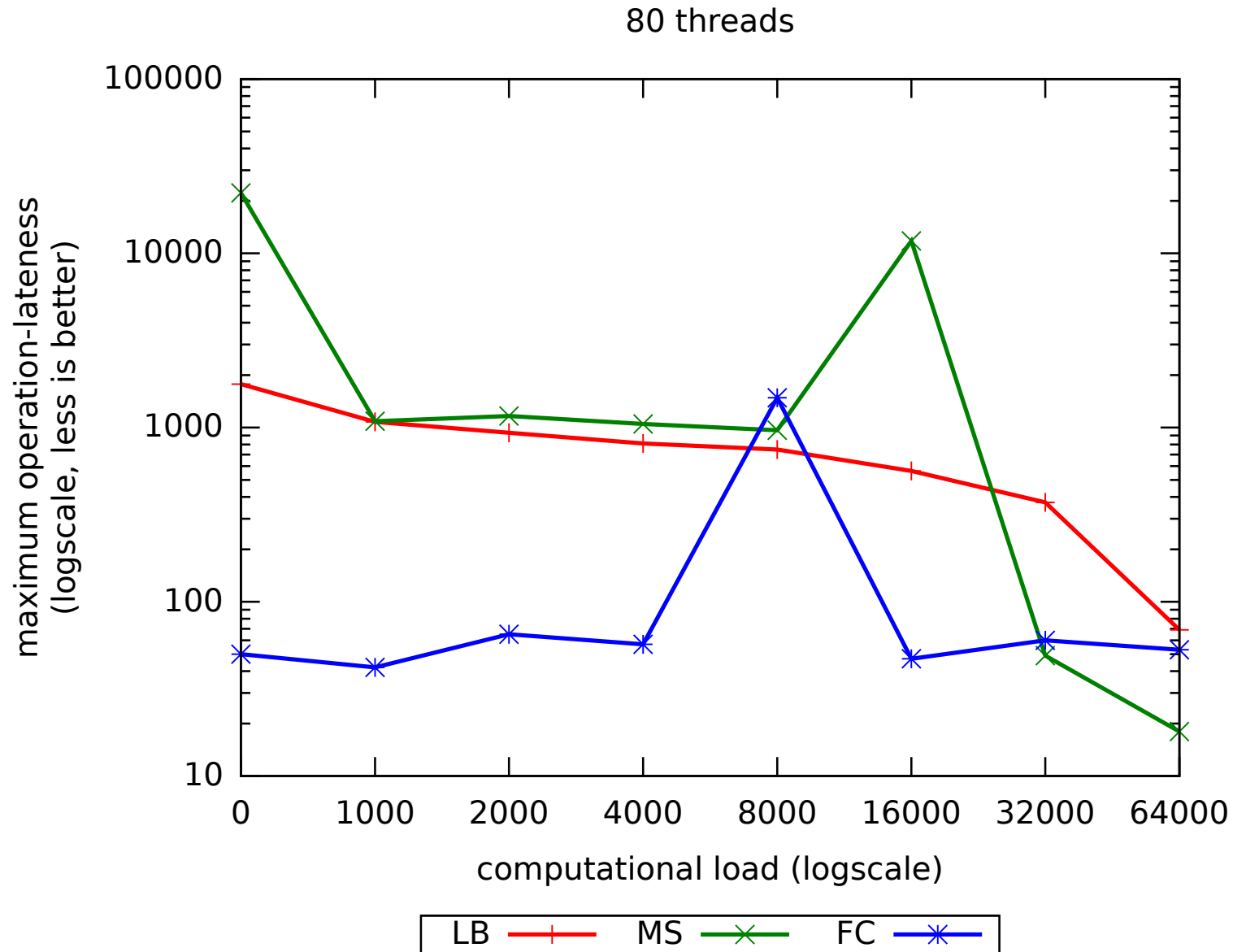
one thread does

dequeue all elements sequentially

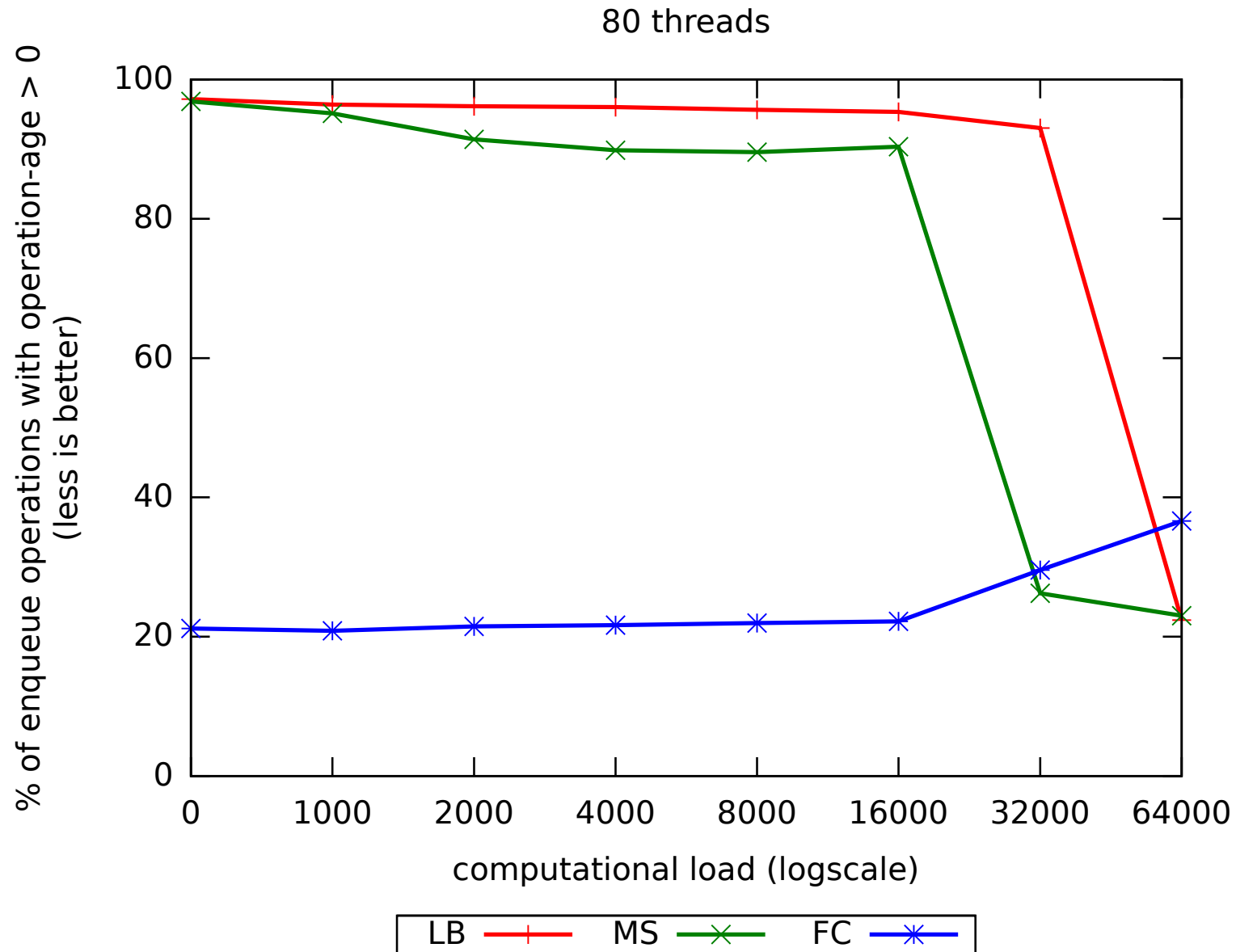
# Maximum Operation-Age



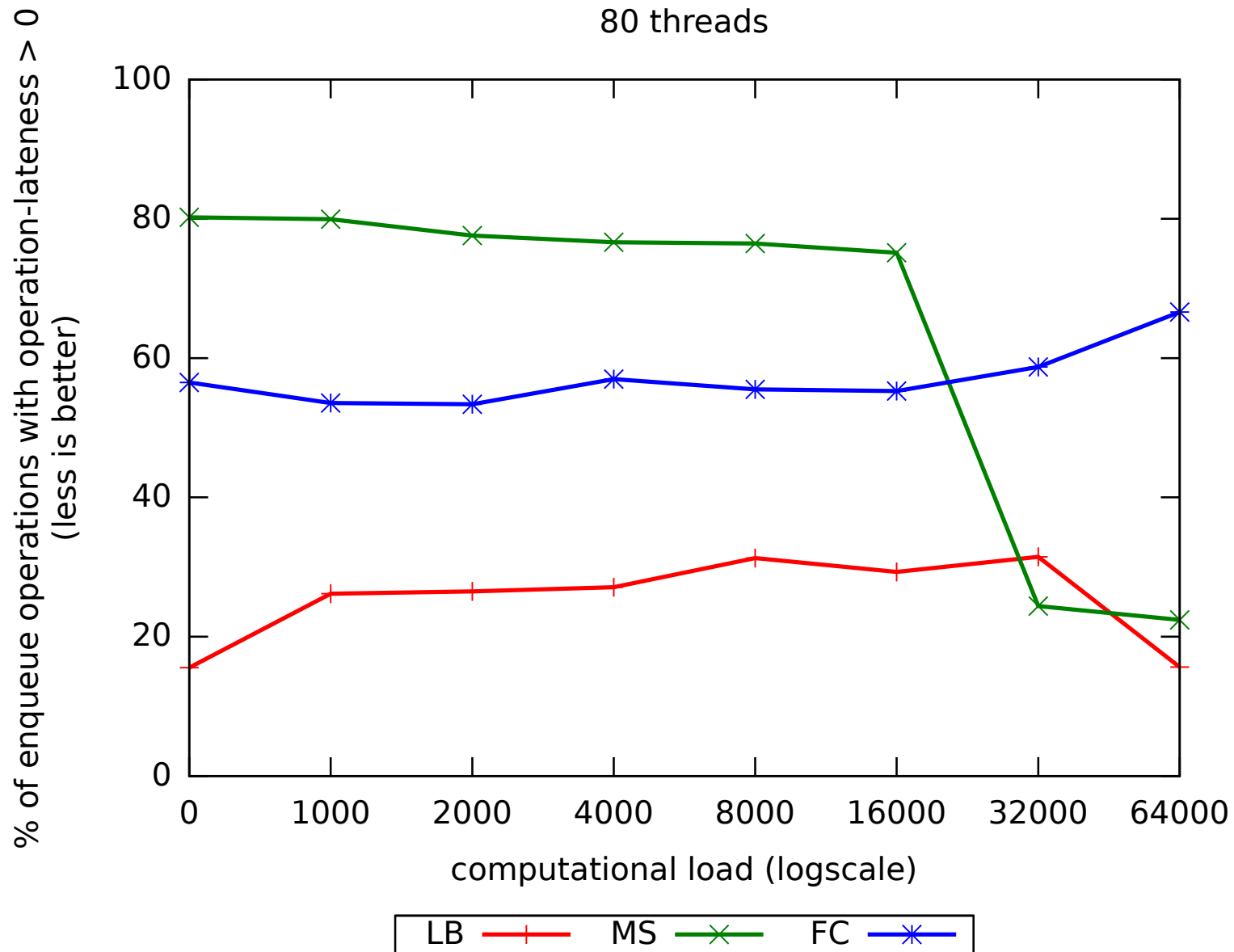
# Maximum Operation-Latency



# Number of Overtaking Operations



# Number of Overtaken Operations





# Conclusion

- ▶ We introduced metrics to compare the behavior of various FIFO queue implementations.
  - ◆ Relaxed implementation can appear more FIFO than strict implementations.
- ▶ Future work
  - ◆ Measure operation-fairness of relaxed FIFO queue implementations.
  - ◆ Use element-fairness to analyze implementation of other data structures, e.g. stacks.

# Thank You

For more information about the queue implementations see  
<http://scal.cs.uni-salzburg.at/>

Additional measurement results can be seen on  
<http://scal.cs.uni-salzburg.at/races12/>