

Model Checking with BDDs

Sets as Propositions

- Consider the universe “ABCD”

$[('A', [\sim p_2, \sim p_1]), ('B', [\sim p_2, p_1]),$
 $, ('C', [p_2, \sim p_1]), ('D', [p_2, p_1])]$

- Or the universe [1,5,6,79,13]

$[(1, [\sim p_3, \sim p_2, \sim p_1]), (5, [\sim p_3, \sim p_2, p_1])$
 $, (6, [\sim p_3, p_2, \sim p_1]), (79, [\sim p_3, p_2, p_1])$
 $, (13, [p_3, \sim p_2, \sim p_1])]$

Consider some subsets

`subset "ABCD" "A"`

`~p1 /\ ~p2`

`subset "ABCD" "AC"`

`(~p1 /\ ~p2) \/ (~p1 /\ p2)`

`subset "ABCD" "ACDB"`

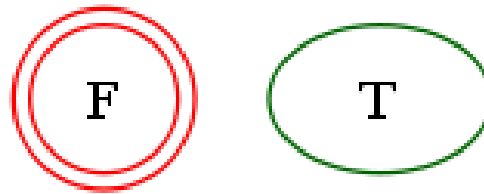
`(~p1 /\ ~p2) \/ (p1 /\ ~p2) \/ (~p1 /\ p2)
 \/ (p1 /\ p2)`

`subset "ABCD" ""`

`Absurd`

And their BDDs

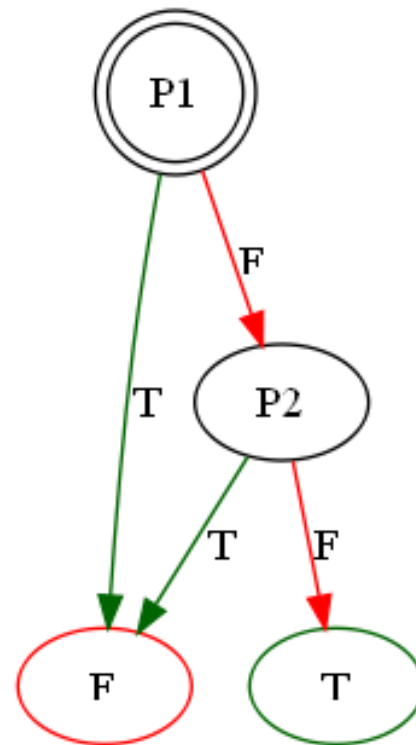
- Set "ABCD" ""
- Absurd



$[('a', [\sim p2, \sim p1]), ('b', [\sim p2, p1]),$
 $('c', [p2, \sim p1]), ('d', [p2, p1])]]$

subset "abcd" "a"

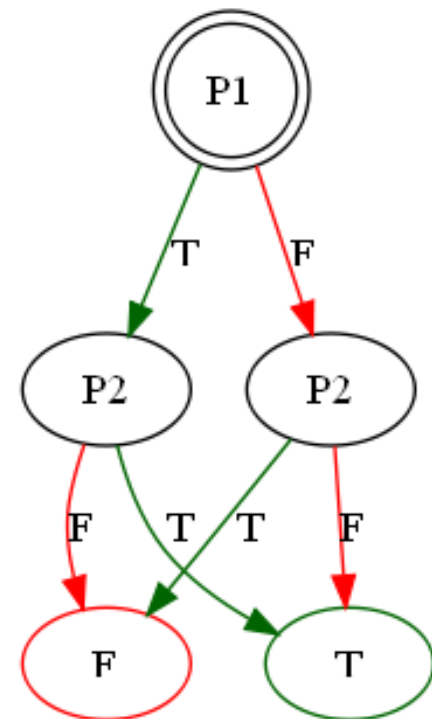
$\sim p1 \wedge \sim p2$



$[('a', [\sim p_2, \sim p_1]), ('b', [\sim p_2, p_1]),$
 $('c', [p_2, \sim p_1]), ('d', [p_2, p_1])]$

subset "abcd" "ad"

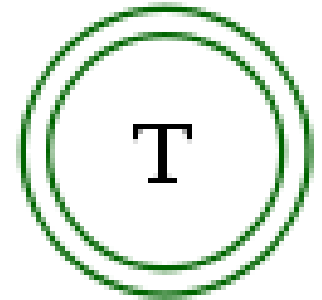
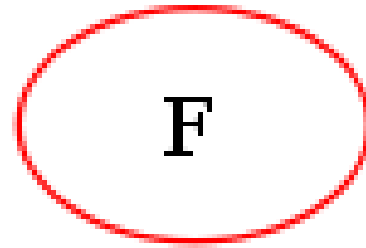
$(\sim p_1 \wedge \sim p_2) \vee$
 $(p_1 \wedge p_2)$



```
[ ('a', [~p2, ~p1]), ('b', [~p2, p1]),  
  ('c', [p2, ~p1]), ('d', [p2, p1]) ]
```

```
subset "abcd" "adbc"
```

```
(~p1 /\ ~p2) \/  
(p1 /\ ~p2) \/  
 (~p1 /\ p2) \/  
(p1 /\ p2)
```



```

initial xs = zip xs (reverse (g 1))
  where n = numbits (length xs)
        g :: Int -> [[Prop Int]]
        g m | m > n = [[]]
        g n = map (LetterP n:) ys ++ map ((NotP
(LetterP n)):) ys
            where ys = (g (n+1))
subset univ set = foldr acc AbsurdP univ
  where acc x prop | elem x set = orOpt (get x) prop
        acc x prop = prop
        mapping = initial univ
        get n = case lookup n mapping of
                  Just literals -> andL literals

```


Membership test

- Represent an element of a set as the singleton subset
- $\text{item univ } x = \text{subset univ } [x]$
- Then membership uses the tautology
- $\{x\} \text{ `elem` } zS \text{ iff } \{x\} == \{x\} \cap zS$

Lift to BDDs

`subsetB x y = p2b (subset x y)`

`itemB x y = p2b(item x y)`

`mem univ x xs = same temp (conj
temp (subsetB univ xs))`

`where temp = (itemB univ x)`

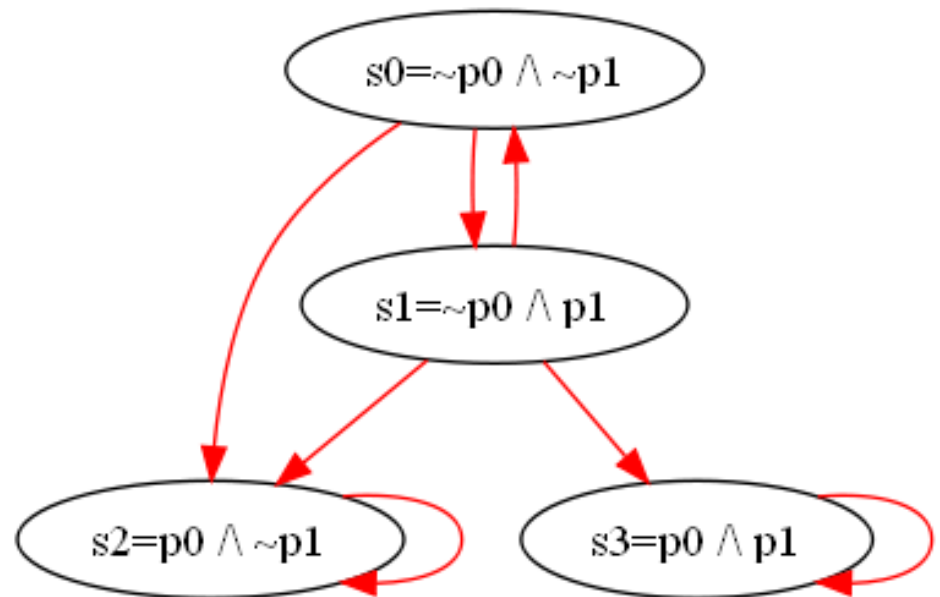
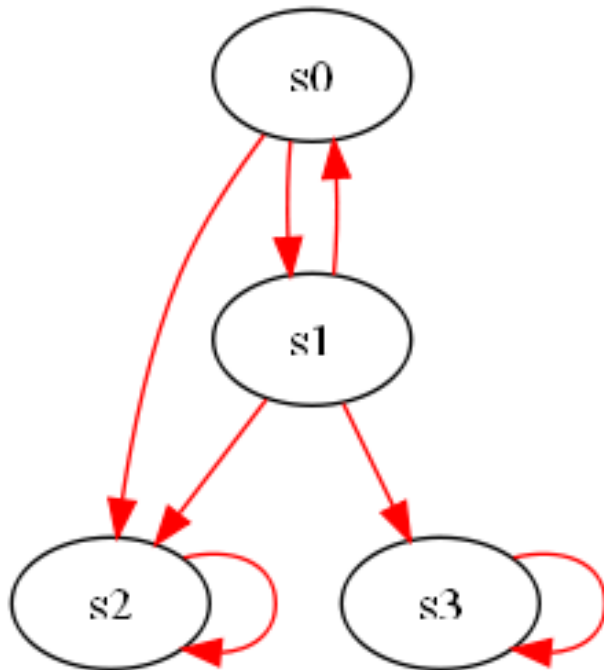
A relation between two items in a set

```
[ ( ' a ' , [ ~p2 , ~p1 ] ) ,  
  ( ' b ' , [ ~p2 , p1 ] ) ,  
  ( ' c ' , [ p2 , ~p1 ] ) ,  
  ( ' d ' , [ p2 , p1 ] ) ]
```

- $R(a,b) = \text{True}$
- $R(b,c) = \text{True}$
- $R(c,c) = \text{True}$
- $R(_,_) = \text{False}$

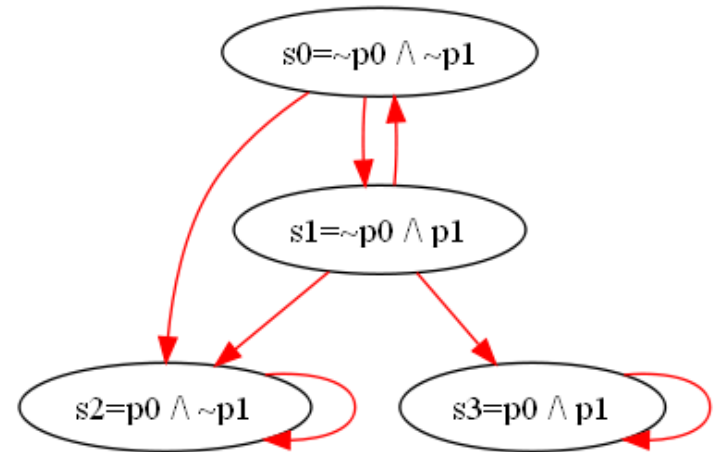
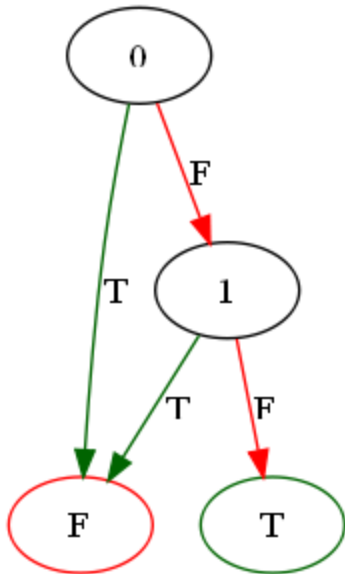
Graph transitions

- Consider the graph and its assignment of states to boolean formula



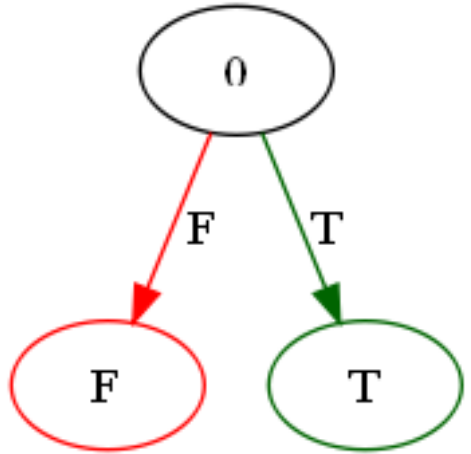
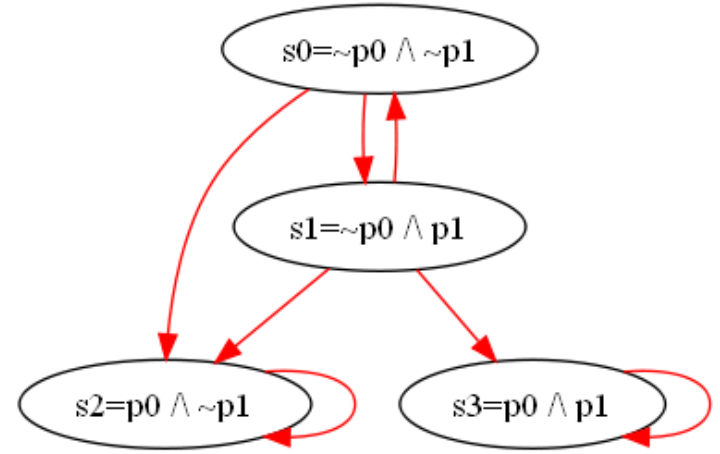
Recall how we represent subsets

- $u1 = [0,1,2,3]$
- $sub1 = \text{subset } u1 [0]$
- $\sim p0 \wedge \sim p1$



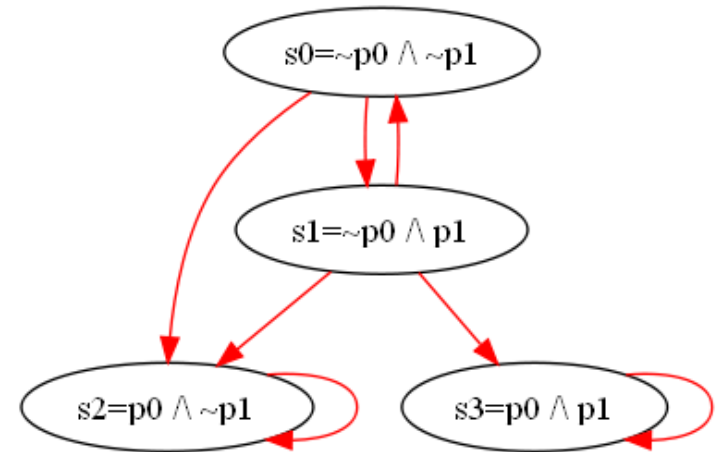
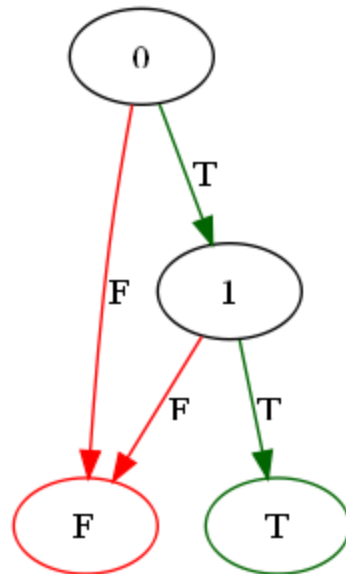
Subset {s2,s3}

- $u1 = [0,1,2,3]$
- $sub2 = subset\ u1\ [2,3]$
- $(p0 \wedge \sim p1) \vee (p0 \wedge p1)$



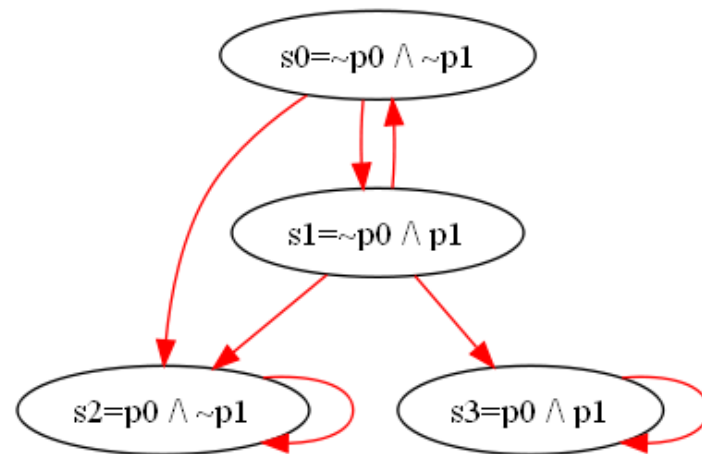
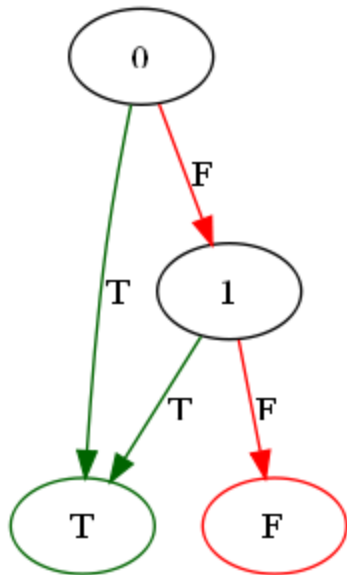
Subset {s3}

- $u1 = [0,1,2,3]$
- $sub3 = subset\ u1\ [3]$
- $p0 \wedge p1$



Subset $\{s1,s2,s3\}$

- $u1 = [0,1,2,3]$
- $sub4 = \text{subset } u1 [1,2,3]$
- $(\sim p0 \wedge p1) \vee (p0 \wedge \sim p1) \vee (p0 \wedge p1)$



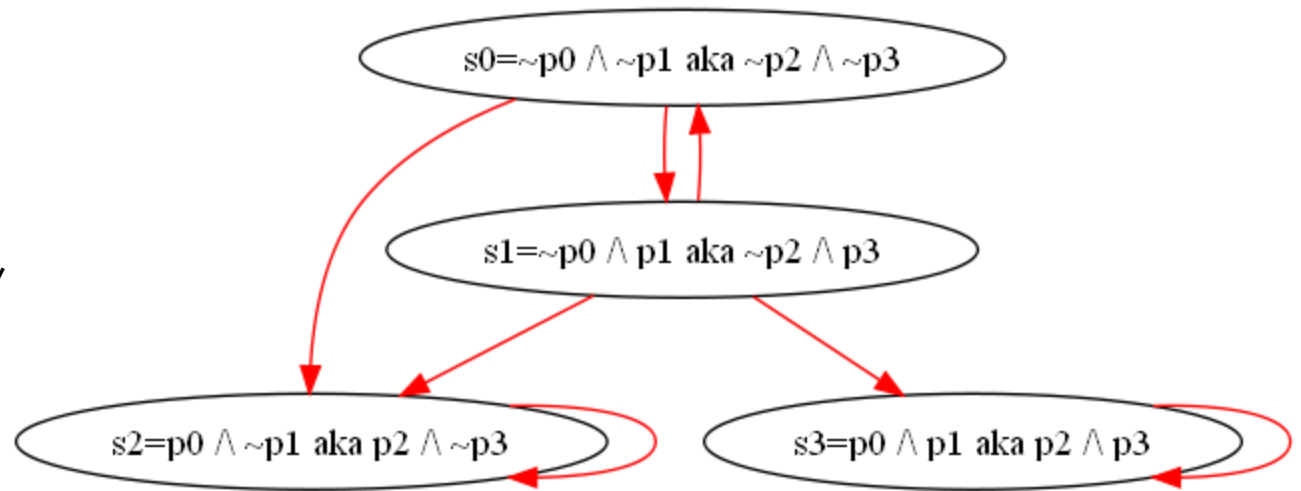
- Introduce new variables p2 and p3 that mirror p0 and p1

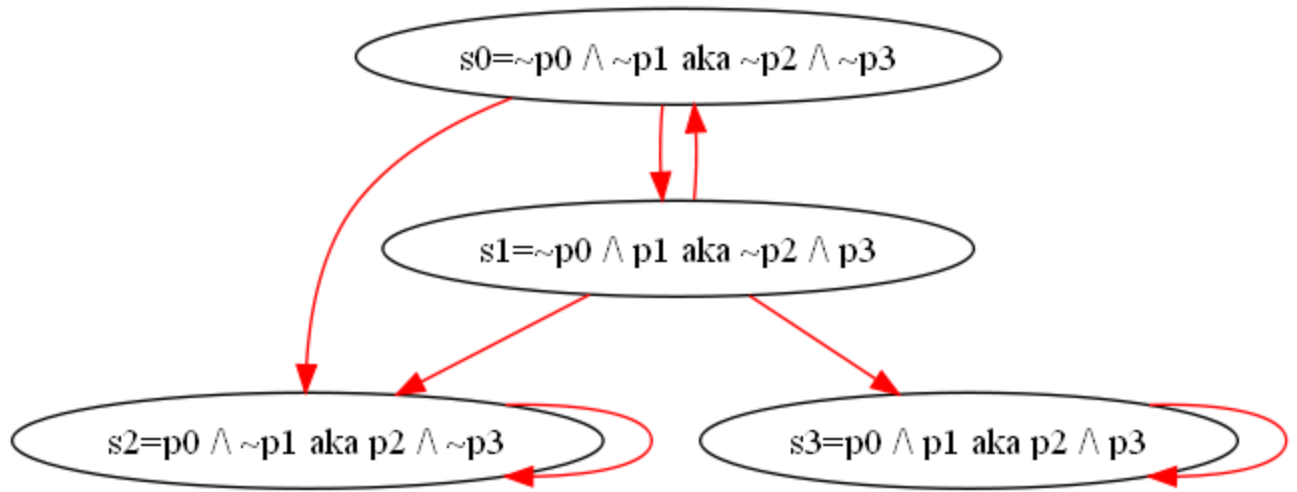
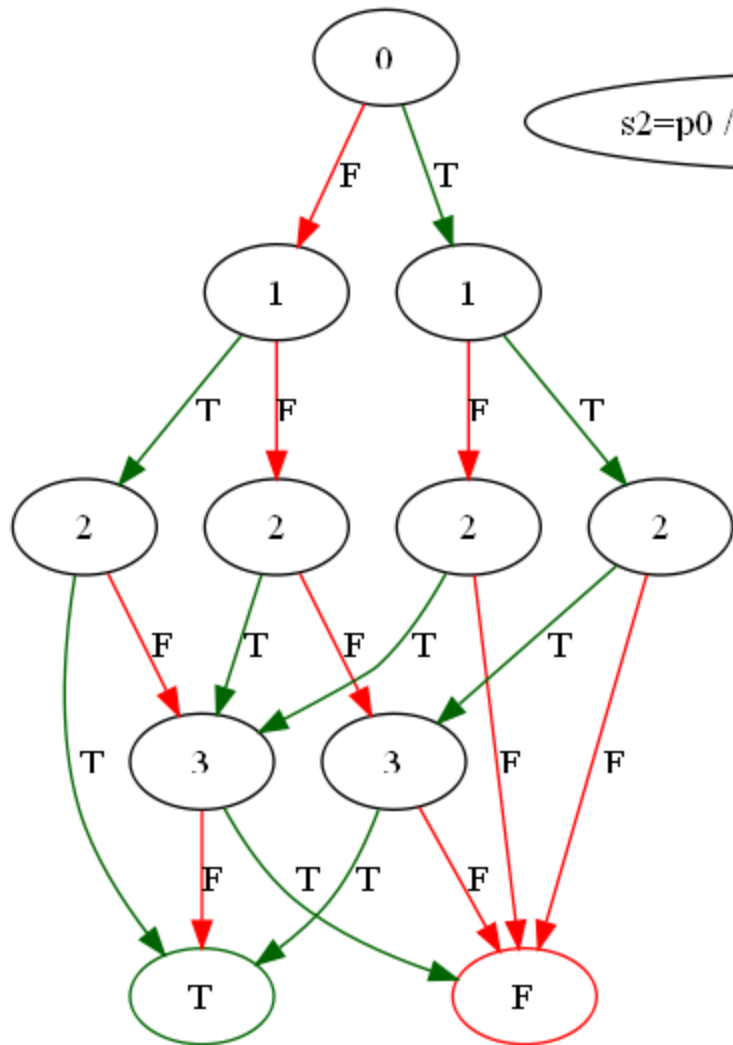
The transition relation

```

prop1 =
(~p0 /\ ~p1 /\ ~p2 /\ p3)
\|
(~p0 /\ ~p1 /\ p2 /\ ~p3)
\|
(~p0 /\ p1 /\ ~p2 /\ ~p3)
\|
(~p0 /\ p1 /\ p2 /\ ~p3)
\|
(p0 /\ ~p1 /\ p2
\|
(~p0 /\ p1 /\ p2
\|
(p0 /\ p1 /\ p2 /

```





As a BDD

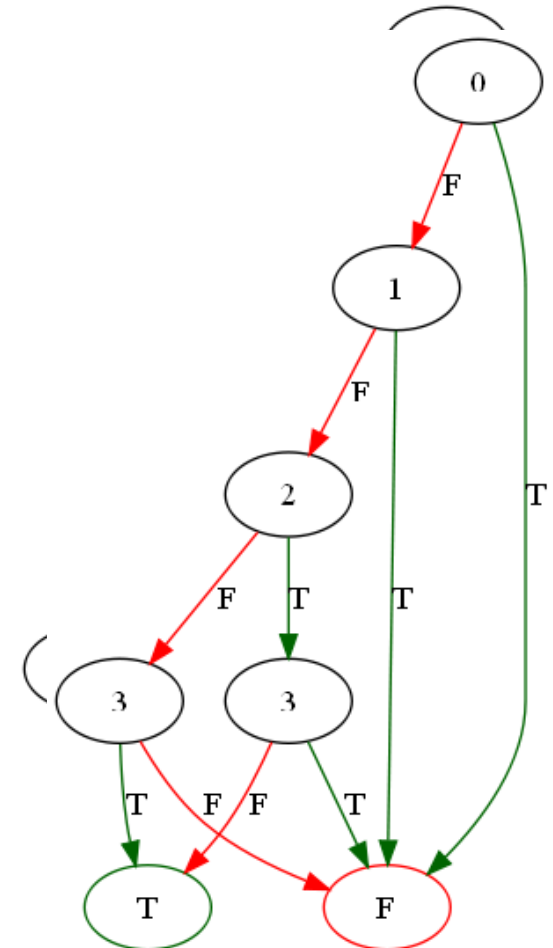
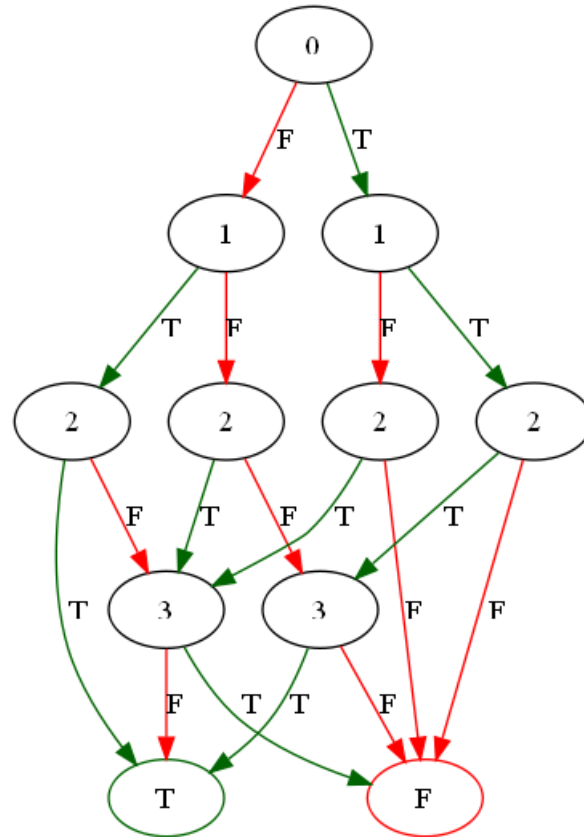
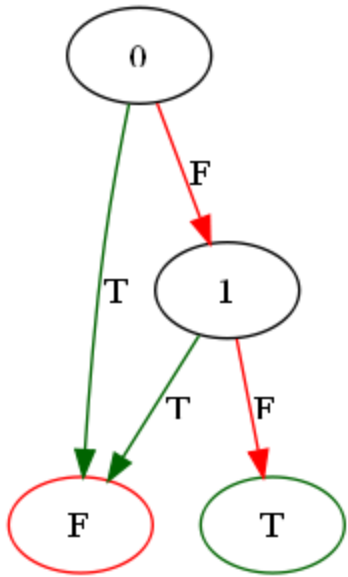
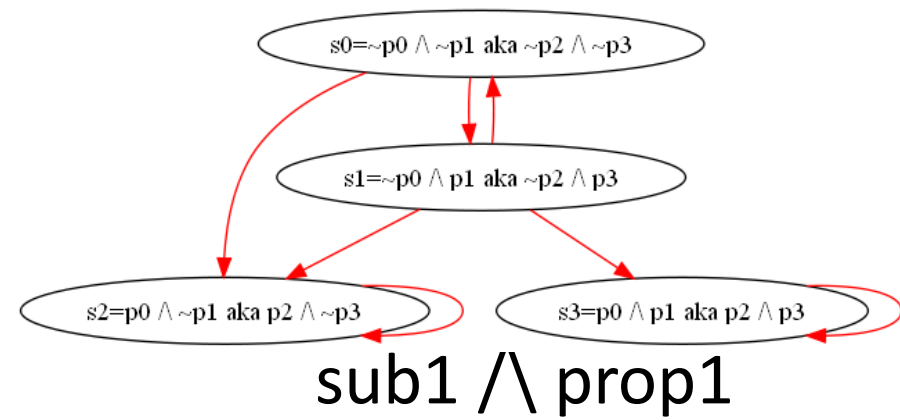
States reachable in one step

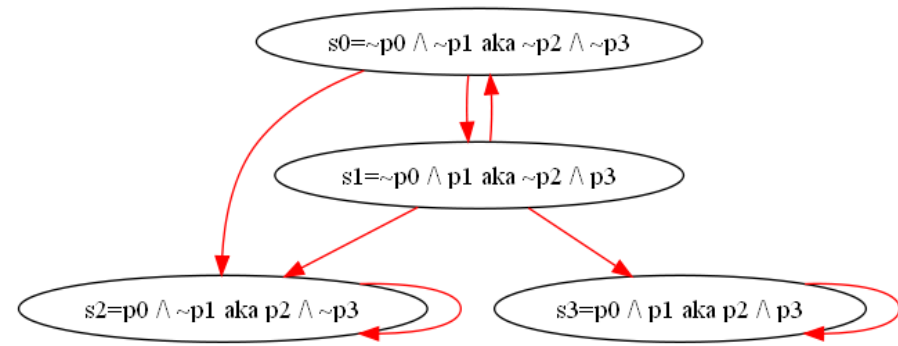
- Let sub be a set of states
- What is reachable in one step?
- $step\ set = prop1 \wedge set$

To take a step conjoin

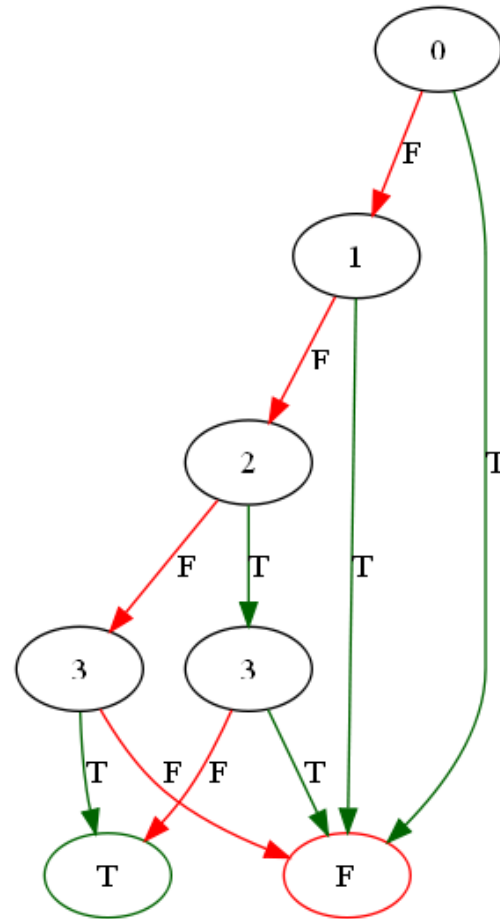
- $sub1 = \{s0\}$

prop1



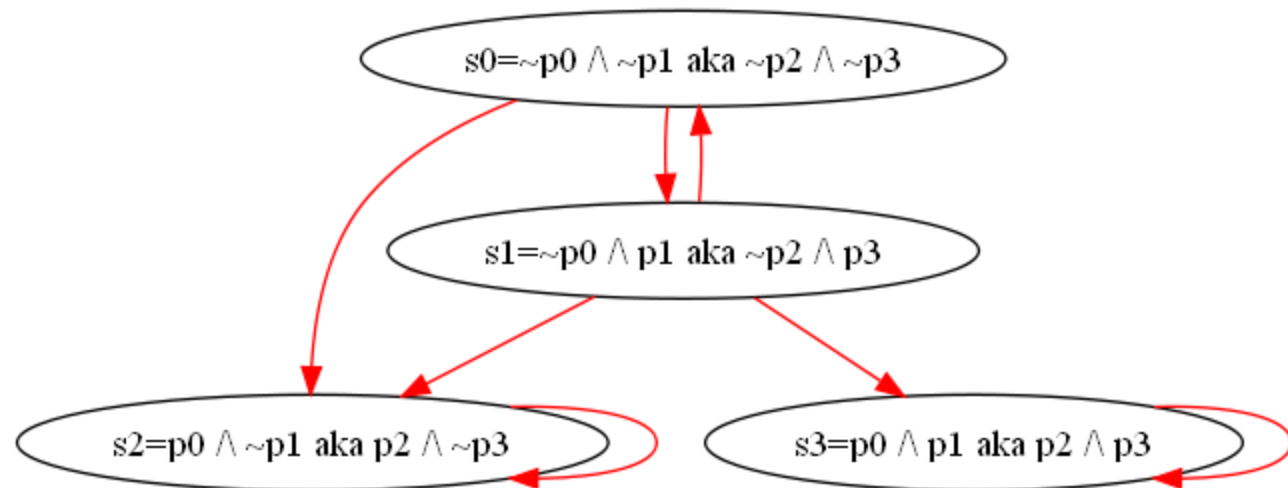


- Note the paths to True
- There are two of them
- Each corresponds to one next state
- The values of p_2 and p_3 tell what states $\{s_1, s_2\}$

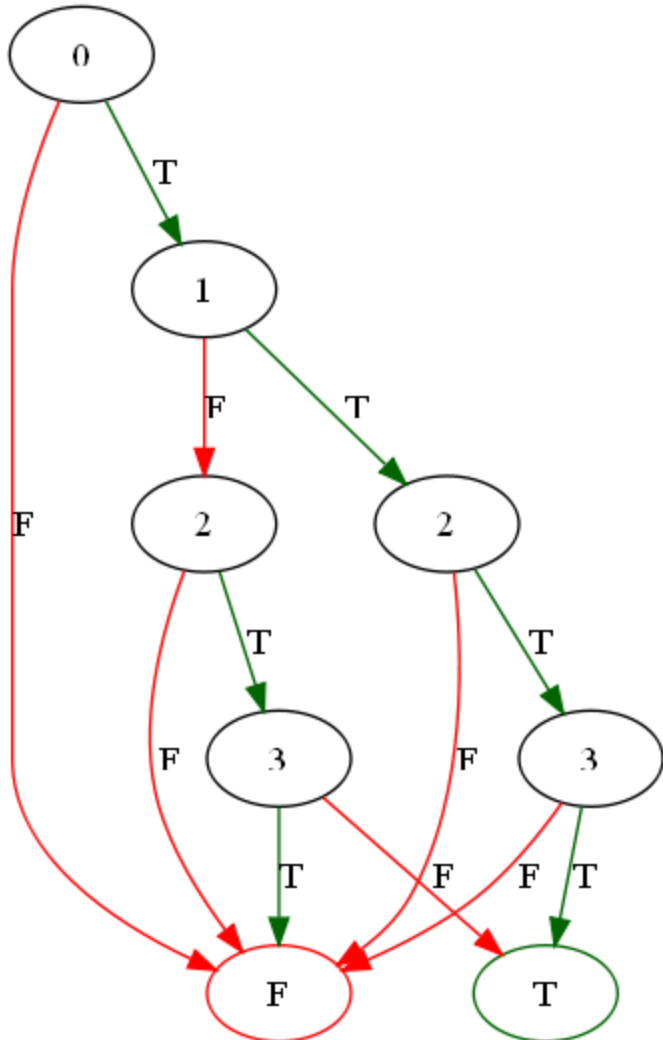


Consider the solutions

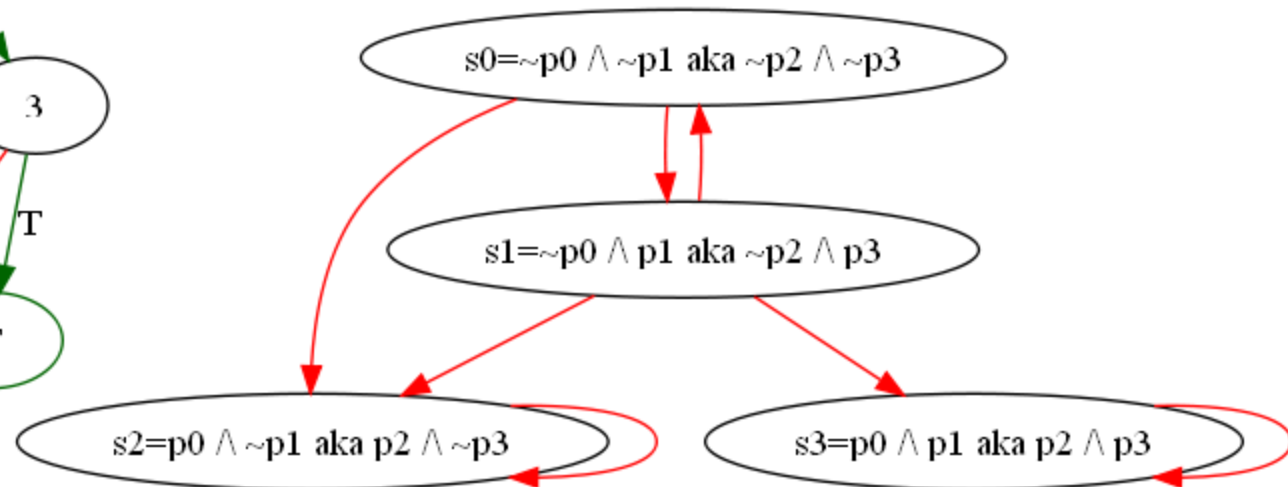
- $[(0, \text{False}), (1, \text{False}), (2, \text{False}), (3, \text{True})]$,
- $[(0, \text{False}), (1, \text{False}), (2, \text{True}), (3, \text{False})]$
- By throwing away the assignments to p_0 and p_1 , and by renaming p_2 to p_0 , and p_3 to p_1 , we get to solutions
- $\sim p_0 \wedge p_1$
- $p_0 \wedge \sim p_1$
- Cooresponding
- to the states
- $\{s_1, s_2\}$



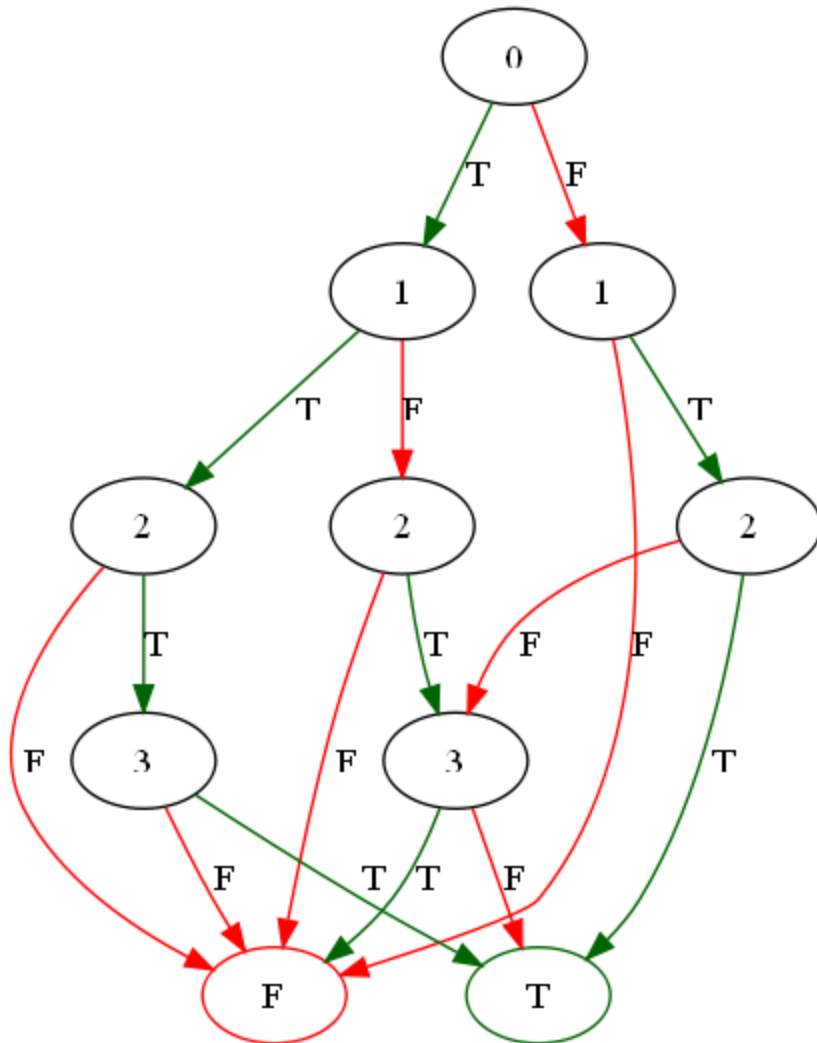
Start at the set {s2,s3}



- $[(0, \text{True}), (1, \text{False}), (2, \text{True}), (3, \text{False})]$
- $[(0, \text{True}), (1, \text{True}), (2, \text{True}), (3, \text{True})]$
- $p_0 \wedge \sim p_1$
- $p_0 \wedge p_1$



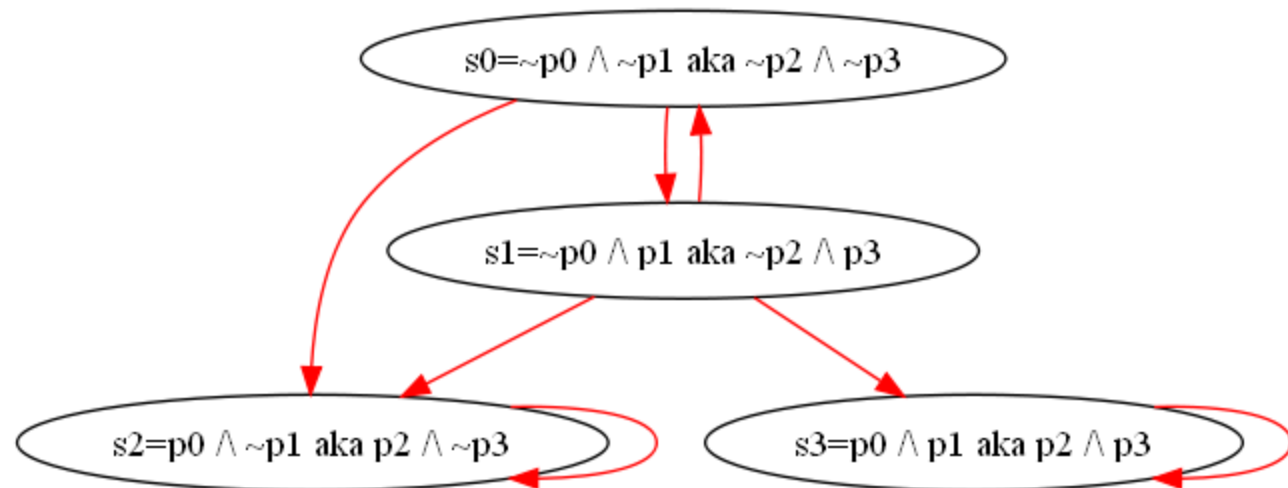
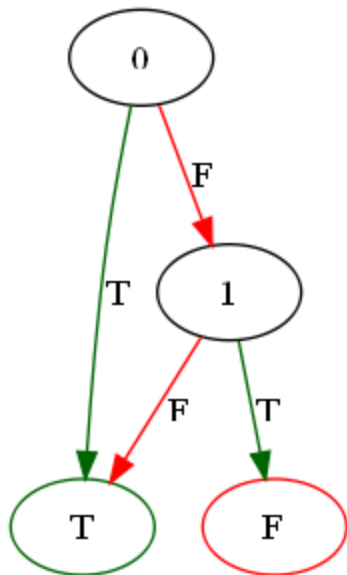
Start at the subset $\{s1,s2,s3\}$



- $\sim p0 \wedge \sim p1$
- $p0$
- $p0 \wedge \sim p1$
- $p0 \wedge p1$

$$p_0 \vee (p_0 \wedge p_1) \vee (p_0 \wedge \sim p_1) \vee (\sim p_0 \wedge \sim p_1)$$

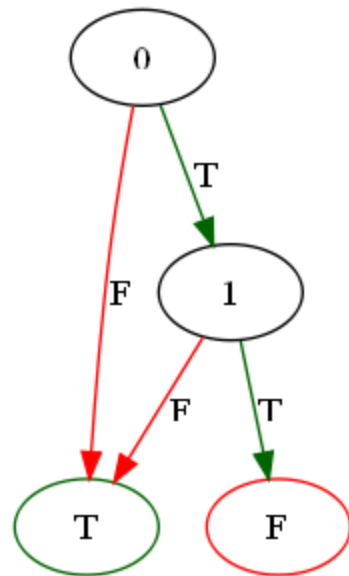
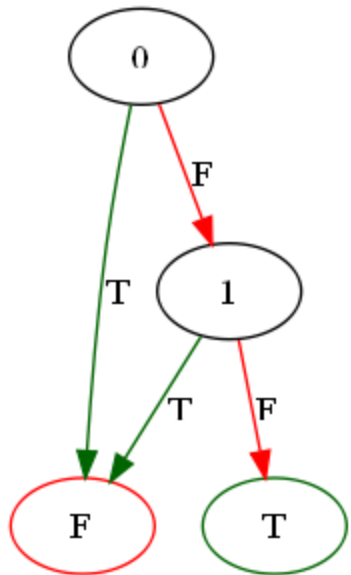
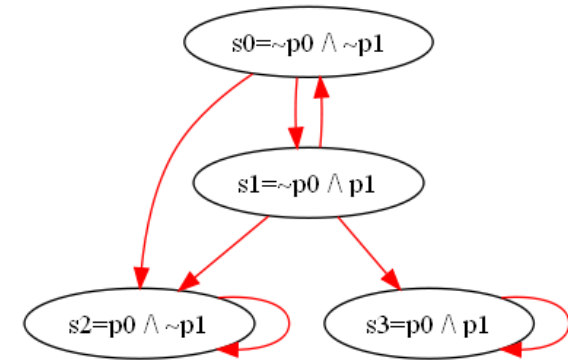
- This corresponds to the BDD
- Which is every state except s_2 , which is exactly what can be reached from $\{s_1, s_2, s_3\}$



To take multiple steps

- Compute the states reachable in 1 step
- Union in the starting states
- And repeat

- $sub1 = \{s0\}$
- *LectureBDD> pnG (p2b sub1)
- *LectureBDD> pnG (p2b (step sub1))
- *LectureBDD> pnG (p2b (step (step sub1)))
- *LectureBDD> pnG (p2b (step (step (step sub1))))



- $\text{sub3} = \{s3\}$
- *LectureBDD> pnG (p2b sub3)
- *LectureBDD> pnG (p2b (step sub3))

