

# 7 steps to adding a new solver

The implementation of Funlog

# Multiple solvers

- Funlog incorporates several solvers
  - SAT solver
  - SMT solver
  - Mathematical programming solver
- Earlier version have also incorporated a max-Sat solver and a narrowing solver.

# Modality

- A Funlog program has two kinds of modalities
  - Evaluation
  - Search

```
dim i10#Int = [0,1,2,3,4,5,6,7,8,9]
dim colors#String = ["Red","Blue","Green","Yellow"]
```

```
graph = [(1,2),(2,3),(3,4),(4,5),
         (5,1),(1,6),(2,7),(3,8),
         (4,9),(5,0),(6,8),(7,9),
         (8,0),(9,6),(0,7)]
```

```
edges = set #(i10,i10) graph
-- color = toSet colors
```

```
pairs = [(1,"Red"),(2,"Red"),(3,"Blue")]
tim = set #(i10,colors) pairs
```

```
same r =
  let f(x,y,c) -> i10(x), i10(y), colors(c).
      f(x,y,c) <- $r(x,c), edges(x,y), $r(y,c).
  in f
```

```
justNodes r = $({ (n) <- $r(n,c)})
```

```
exists coloring: set #(i10,colors) none .. fullRel #(i10,colors)
where none (same coloring)
      && full (justNodes coloring)
find Many 4
by SAT
```

# Overloading and Staging

- A primitive function or a user defined function can be used in both modalities
- We think of these functions being overloaded (i.e. having two (or more) separate implementations.
- Internally we use staging to decide which implementation to use.

# The Boolean Class

```
-- Something acts like a Boolean if it supports  
-- these operations
```

```
class Show b => Boolean b where  
  true  :: b  
  false :: b  
  isTrue  :: b -> Bool  
  isFalse :: b -> Bool  
  conj :: b -> b -> b      -- conjunction  
  disj :: b -> b -> b      -- disjunction  
  neg  :: b -> b          -- negation  
  imply :: b -> b -> b    -- implication
```

# Number class

```
class NumLike t where
  liftI :: Int -> t
  liftR :: Rational -> t
  (+) :: t -> t -> t
  (*) :: t -> t -> t
```

# Comparisons

```
class(NumLike t, BoolLike b)
```

```
  => Compare t b where
```

```
(<=) :: t -> t -> b
```

```
(==) :: t -> t -> b
```

# Sets or Relations

class SetLike s where

create:: Dim a -> [a] -> s a

select:: (a -> Bool) -> s a -> s a

union:: s a -> s a -> s a

proj3of3:: s (a,b,c) -> s c



# A solver

- Each solver answers questions over some kind of data.
- Each solver chooses a second representation type for the data it knows how to solve.
- The representation internally represents (part of) the input to a solver.

# 7 steps

1. Choose a representation type
2. Overload the primitive functions for the “real” type and the “representation” type
3. Describe how to initialize unknown existentially quantified data in the representation type.
4. Stage the search modality code
5. Execute the staged code to generate constraints
6. Format the constraints as input to a solver
7. Instantiate the solution back into a “real” piece of data.

# SAT (Finite Set) Solvers. Step 1

```
data SAT =  
    VarP Int  
  | FalseP  
  | TruthP  
  | AndP SAT SAT  
  | OrP SAT SAT  
  | ImpliesP SAT SAT  
  | NotP SAT
```

# SAT (Finite Set) Solvers. Step 2

```
instance BoolLike SAT where
  true  = TruthP
  false = FalseP
  (&&)  = AndP
  liftB True  = TruthP
  liftB False = FalseP
```

# SAT (Finite Set) Solvers. Step 2

```
instance BoolLike b => SetLike (BitVector b)
  where
    create d xs = ...
    select p (BV d xs) =
      BV d [(x, liftB (p x)) | (x,b) <- xs]
    proj3of3 (BV (D3 _ _ d) xs) = ...
    join (BV (D2 a b) xs) (BV (D2 _ c) ys) =
      ...
```

# Example SAT problem

```
n = 3
dim size#Int = [1..n]
dim node#Char = "abcdefg"

-- cover(i,j,x,y,m,n) A block of size i*j, at point (x,y), covers squares (m,n)
cover = set #(size,size,size,size,size,size)
  [ (i,j,x,y,x+m,y+n)
    | i <- size
    , j <- size
    , x <- size, i+x <= n+1
    , y <- size, j+y <= n+1
    , m <- [0..i-1]
    , n <- [0..j-1]]

rect = set #(node,size,size) [('a',1,1),('b',1,2),('c',1,3),('d',2,1),('e',2,2),('f',3,2)]

possible(nm,i,j,x,y,m,n) -> node(nm),size(i),size(j),size(x),size(y),size(m),size(n).
possible(nm,i,j,x,y,m,n) <- rect(nm,i,j), cover(i,j,x,y,m,n).

exists sol : set #(node,size,size,size,size,size,size) none .. possible
  where full $( { (m,n) <- sol(nm,i,j,x,y,m,n) } )    && -- every pair (m,n) is covered
            $( sol(nm,i,j,x,y,m,n) | nm -> (x,y) )    && -- Each rect is used at most once
            $( sol(nm,i,j,x,y,m,n) | (m,n) -> nm )    -- each pair is covered once.

  find Abstract
  by SAT

placement(nm,x,y) -> node(nm), size(x), size(y) .
placement(nm,x,y) <- sol(nm,i,j,x,y,m,n).
```

# The concrete cover relation

(Int#3, Int#3, Int#3, Int#3, Int#3, Int#3)

{ (1,1,1,1,1,1) (1,1,1,2,1,2) (1,1,1,3,1,3) (1,1,2,1,2,1)  
(1,1,2,2,2,2) (1,1,2,3,2,3) (1,1,3,1,3,1) (1,1,3,2,3,2)  
(1,1,3,3,3,3) (1,2,1,1,1,1) (1,2,1,1,1,2) (1,2,1,2,1,2)  
(1,2,1,2,1,3) (1,2,2,1,2,1) (1,2,2,1,2,2) (1,2,2,2,2,2)  
(1,2,2,2,2,3) (1,2,3,1,3,1) (1,2,3,1,3,2) (1,2,3,2,3,2)  
(1,2,3,2,3,3) (1,3,1,1,1,1) (1,3,1,1,1,2) (1,3,1,1,1,3)  
(1,3,2,1,2,1) (1,3,2,1,2,2) (1,3,2,1,2,3) (1,3,3,1,3,1)  
(1,3,3,1,3,2) (1,3,3,1,3,3) (2,1,1,1,1,1) (2,1,1,1,2,1)  
(2,1,1,2,1,2) (2,1,1,2,2,2) (2,1,1,3,1,3) (2,1,1,3,2,3)  
(2,1,2,1,2,1) (2,1,2,1,3,1) (2,1,2,2,2,2) (2,1,2,2,3,2)

```
possible(nm,i,j,x,y,m,n) <- rect(nm,i,j), cover(i,j,x,y,m,n).
rect = set #(node,size,size)
[('a',1,1),('b',1,2),('c',1,3),('d',2,1),('e',2,2),('f',3,2)]
```

```
exp> possible
```

```
(Char#7,Int#3,Int#3,Int#3,Int#3,Int#3,Int#3)
```

```
{('a',1,1,1,1,1,1) ('a',1,1,1,2,1,2) ('a',1,1,1,3,1,3)
 ('a',1,1,2,1,2,1) ('a',1,1,2,2,2,2) ('a',1,1,2,3,2,3)
 ('a',1,1,3,1,3,1) ('a',1,1,3,2,3,2) ('a',1,1,3,3,3,3)
 ('b',1,2,1,1,1,1) ('b',1,2,1,1,1,2) ('b',1,2,1,2,1,2)
 ('b',1,2,1,2,1,3) ('b',1,2,2,1,2,1) ('b',1,2,2,1,2,2)
 ('b',1,2,2,2,2,2) ('b',1,2,2,2,2,3) ('b',1,2,3,1,3,1)
 ('b',1,2,3,1,3,2) ('b',1,2,3,2,3,2) ('b',1,2,3,2,3,3)
 ('c',1,3,1,1,1,1) ('c',1,3,1,1,1,2) ('c',1,3,1,1,1,3)}
```



# SAT (Finite Set) Step 3

```
exists sol : set #(node,size,size,size,size,size,size) none .. Possible
```

```
exp> sol
```

```
(Char#7,Int#3,Int#3,Int#3,Int#3,Int#3,Int#3)
```

```
{('a',1,1,1,1,1,1)=p1 ('a',1,1,1,2,1,2)=p2 ('a',1,1,1,3,1,3)=p3  
 ('a',1,1,2,1,2,1)=p4 ('a',1,1,2,2,2,2)=p5 ('a',1,1,2,3,2,3)=p6  
 ('a',1,1,3,1,3,1)=p7 ('a',1,1,3,2,3,2)=p8 ('a',1,1,3,3,3,3)=p9  
 ('b',1,2,1,1,1,1)=p10 ('b',1,2,1,1,1,2)=p11 ('b',1,2,1,2,1,2)=p12  
 ('b',1,2,1,2,1,3)=p13 ('b',1,2,2,1,2,1)=p14 ('b',1,2,2,1,2,2)=p15  
 ('b',1,2,2,2,2,2)=p16 ('b',1,2,2,2,2,3)=p17 ('b',1,2,3,1,3,1)=p18  
 ('b',1,2,3,1,3,2)=p19 ('b',1,2,3,2,3,2)=p20 ('b',1,2,3,2,3,3)=p21  
 ('c',1,3,1,1,1,1)=p22 ('c',1,3,1,1,1,2)=p23 ('c',1,3,1,1,1,3)=p24  
 ('c',1,3,2,1,2,1)=p25 ('c',1,3,2,1,2,2)=p26 ('c',1,3,2,1,2,3)=p27
```

# SAT (Finite Set) Step 4

```
col r n = $( { k <- r(i,j,k), j =#size $n } )  
row r n = $( { k <- r(i,j,k), i =#size $n } )  
box r n = $( { k <- r(i,j,k), square(i,j,$n) } )
```

```
exists grid : set #(size,size,digit) none .. full  
  where  and [ full (col grid i) | i <- size ] &&  
         and [ full (row grid i) | i <- size ] &&  
         and [ full (box grid i) | i <- size ] &&  
         $(grid(i,j,n),grid(i,j,m) -> n =#digit m)  
find Many (setToArray grid)  
  by SAT
```

# SAT (Finite Set) Step 5

where full  $\$( \{ (m,n) \leftarrow \text{sol}(nm,i,j,x,y,m,n) \} )$       &&    --  
every pair  $(m,n)$  is covered

$\$( \text{sol}(nm,i,j,x,y,m,n) \mid nm \rightarrow (x,y) )$       &&    --  
Each rect is used at most once

$\$( \text{sol}(nm,i,j,x,y,m,n) \mid (m,n) \rightarrow nm )$       --  
each pair is covered once.

Abstract where clause

$(p1 \ \backslash / \ p10 \ \backslash / \ p22 \ \backslash / \ p31 \ \backslash / \ p43 \ \backslash / \ p59) \ / \backslash$   
 $(p2 \ \backslash / \ p11 \ \backslash / \ p12 \ \backslash / \ p23 \ \backslash / \ p33 \ \backslash / \ p44 \ \backslash / \ p47 \ \backslash / \ p60 \ \backslash / \ p65)$   
   $/ \backslash$   
 $(p3 \ \backslash / \ p13 \ \backslash / \ p24 \ \backslash / \ p35 \ \backslash / \ p48 \ \backslash / \ p66) \ / \backslash$   
 $(p4 \ \backslash / \ p14 \ \backslash / \ p25 \ \backslash / \ p32 \ \backslash / \ p37 \ \backslash / \ p45 \ \backslash / \ p51 \ \backslash / \ p61) \ / \backslash \dots$

# SAT (Finite Set) Step 6

```
p cnf 70 559
1 10 22 31 43 59 0
2 11 12 23 33 44 47 60 65 0
3 13 24 35 48 66 0
4 14 25 32 37 45 51 61 0
5 15 16 26 34 39 46 49 52 55 62 67 0
6 17 27 36 41 50 56 68 0
7 18 28 38 53 63 0
8 19 20 29 40 54 57 64 69 0
9 21 30 42 58 70 0
-1 -2 0
-1 -3 0
-1 -4 0
...
```

# SAT (Finite Set) Step 7

SAT

```
-1 2 -3 -4 -5 -6 -7 -8 -9 -10 -11 -12 -13 14 15 -16 -17 -18 -19 -  
 20 -21 -22 -23 -24 -25 -26 -27 28 29 30 -31 -32 -33 -34 35 36 -  
 37 -38 -39 -40 -41 -42 43 -44 -45 -46 -47 -48 -49 -50 -51 -52 -  
 53 -54 -55 -56 -57 -58 -59 -60 -61 -62 -63 -64 -65 -66 -67 -68  
-69 -70 0
```

```
exp> sol
```

```
(Char#7,Int#3,Int#3,Int#3,Int#3,Int#3,Int#3)  
{('a',1,1,1,1,1,1)=p1 ('a',1,1,1,2,1,2)=p2  
( 'a',1,1,1,3,1,3)=p3  
( 'a',1,1,2,1,2,1)=p4 ('a',1,1,2,2,2,2)=p5  
( 'a',1,1,2,3,2,3)=p6  
( 'a',1,1,3,1,3,1)=p7 ('a',1,1,3,2,3,2)=p8  
( 'a',1,1,3,3,3,3)=p9  
( 'b',1,2,1,1,1,1)=p10 ('b',1,2,1,1,1,2)=p11  
( 'b',1,2,1,2,1,2)=p12
```

```
exp> sol
```

```
(Char#7,Int#3,Int#3,Int#3,Int#3,Int#3,Int#3)  
{('a',1,1,1,2,1,2) ('b',1,2,2,1,2,1) ('b',1,2,2,1,2,2)  
( 'c',1,3,3,1,3,1) ('c',1,3,3,1,3,2) ('c',1,3,3,1,3,3)  
( 'd',2,1,1,3,1,3) ('d',2,1,1,3,2,3) ('e',2,2,1,1,1,1)}
```

# mathematical programming problem

```
sum [] = 0
sum [x] = x
sum (x:xs) = x + sum xs

and [] = True
and [x] = x
and (x:xs) = x && and xs

-----

-- Production minimization problem

data Factory = A | B | C
data Store = NYC | ATL | LA

pairs = #(Factory,Store)

ship = array pairs [2,3,5,3,2,1,3,4,2]
sales = array Store [230,140,300]

exists prod: Array #(Factory,Store) Int
  where sum[ prod.(A,s) | s <- Store ] <= 150 &&
        and [ prod.(f,s) >= 0 | (f,s) <- pairs ] &&
        and [ sales.s == sum [prod.(f,s) | f <- Factory]
              | s <- Store ]
  find Min sum[ prod.(f,s) * ship.(f,s)
              | (f,s) <- pairs ]
  by IP
```

# Math Prog Step 1

## Representation types

```
type PolyNom n = [(String,n)]

-- The meaning of an MExp is
-- a PolyNomial with an additive
-- constant. The polynomial may
-- have no polynomial terms

data MExp n = Term (PolyNom n) n

data Rel n
  = RANGE (PolyNom n) (Range n)
  | TAUT   -- True
  | UNSAT  -- False
deriving Eq
```

# Step 2

```
instance Num n => NumLike (Mexp n) where
  liftI n = Term [] n
  (+) = plusM (+)
```

```
plusM (Term [] a) (Term [] b) = Term [] (a+b)
plusM (Term [] a) (Term ys b) = Term ys (a+b)
plusM (Term xs a) (Term [] b) = Term xs (a+b)
plusM (Term xs a) (Term ys b) = Term (mergeP (+) xs ys) (a+b)
```

```
-- in a Sum, if the same indexed variables appears twice,
-- we add the coefficients
```

```
mergeP :: (Num n, Eq n) => (n -> n -> n) -> PolyNom n -> PolyNom n -> PolyNom n
mergeP f [] ys = ys
mergeP f xs [] = xs
mergeP f ((x,n):xs)((y,m):ys)=
  case compare x y of
    EQ -> case (f n m) of
      0 -> mergeP f xs ys
      i -> (x,i):mergeP f xs ys
    LT -> (x,n):
      mergeP f xs ((y,m):ys)
    GT -> (y,m):
      mergeP f ((x,n):xs) ys
```



# Step 2

```
instance Num n => BoolLike [Rel n] where
  true = Taut
  false = UnSat
  (&&) = AndP
  liftB True = Taut
  liftB False = UnSat
```

```
andM xs ys = help (sort xs) (sort ys)
  where help (UNSAT:_) ys = [UNSAT]
        help xs (UNSAT:_) = [UNSAT]
        help (TAUT: xs) ys = help xs ys
        help xs (TAUT: ys) = help xs ys
        help [] ys = ys
        help xs [] = xs
        help (RANGE x a:xs) (RANGE y b:ys)
          | x==y = RANGE x (intersectRange a b):help xs ys
        help (RANGE x a:xs)(ys@(RANGE y b:_))
          | x < y = RANGE x a:(help xs ys)
        help (xs@(RANGE x a:_))(RANGE y b:ys)
          | x > y = RANGE y b:(help xs ys)
```

# Step 2

```
instance (Num n) =>
```

```
  Compare (Mexp n) [Rel n] where
```

```
    (<=) = lteqM
```

```
lteqM (Term [] a) (Term [] b) =
```

```
  if (a <= b) then [TAUT] else [UNSAT]
```

```
lteqM (Term [] a) (Term xs b) =
```

```
  [RANGE xs (Range (LtEQ(a-b)) PlusInf)]
```

```
lteqM (Term xs a) (Term [] b) =
```

```
  [RANGE xs (Range MinusInf (LtEQ (b-a)))]
```

```
lteqM (Term xs a) (Term ys b) =
```

```
  [RANGE (mergeP (+) xs (negPoly ys))
```

```
    (Range MinusInf (LtEQ (b-a)))]
```

# Math Prog Step 3

```
exists prod: Array #(Factory,Store) Int
```

Abstract values

```
  prod =  
  NYC ATL LA  
+----+----+----+  
A|`a  |`b  |`c  |  
+----+----+----+  
B|`d  |`e  |`f  |  
+----+----+----+  
C|`g  |`h  |`i  |  
+----+----+----+
```

- The array is concrete, but the values are abstract

# Math Prog Step 4

```
exists prod: Array #(Factory,Store) Int
where sum [ prod.(A,s)
           | s <- Store ] <= (liftI 150) &&
and [ prod.(f,s) >= (liftI 0)
     | (f,s) <- pairs ] &&
and [ sales.s == sum [prod.(f,s)
                     | f <- Factory]
     | s <- Store ]
```

# Math Prog Step 5

Abstract where clause

```
[ 0 <= [ ("a",1) ] < +Inf ,  
  -Inf < [ ("a",1), ("b",1), ("c",1) ] <= 150 ,  
  230 <= [ ("a",1), ("d",1), ("g",1) ] <= 230 ,  
  0 <= [ ("b",1) ] < +Inf ,  
  140 <= [ ("b",1), ("e",1), ("h",1) ] <= 140 ,  
  0 <= [ ("c",1) ] < +Inf ,  
  300 <= [ ("c",1), ("f",1), ("i",1) ] <= 300 ,  
  0 <= [ ("d",1) ] < +Inf ,  
  0 <= [ ("e",1) ] < +Inf ,  
  0 <= [ ("f",1) ] < +Inf ,  
  0 <= [ ("g",1) ] < +Inf ,  
  0 <= [ ("h",1) ] < +Inf ,  
  0 <= [ ("i",1) ] < +Inf ]
```

# Math Prog Step 6

```

NAME          prod
ROWS
  N  COST
  L  R2
  E  R3
  E  R5
  E  R7
COLUMNS
  a      COST      2
  a      R2        1
  a      R3        1
  b      COST      3
  b      R2        1
  b      R5        1
  c      COST      5
  c      R2        1
  c      R7        1
  d      COST      3
  d      R3        1
  e      COST      2
  e      R5        1
  f      COST      1
  f      R7        1
  g      COST      3
  g      R3        1
  h      COST      4
  h      R5        1
  i      COST      2
  i      R7        1
RHS
  RHS    R2        150
  RHS    R3        230
  RHS    R5        140
  RHS    R7        300
BOUNDS
LO BND1  a         0
LO BND1  b         0
LO BND1  c         0
LO BND1  d         0
LO BND1  e         0
LO BND1  f         0
LO BND1  g         0
LO BND1  h         0
LO BND1  i         0
ENDATA

```

```

maximize (2a + 3b + 5c + 3d + 2e
+ 1f + 3g + 4h + 2i) where
0 <= (1a) < +Inf
-Inf < (1a + 1b + 1c) <= 150
230 <= (1a + 1d + 1g) <= 230
0 <= (1b) < +Inf
140 <= (1b + 1e + 1h) <= 140
0 <= (1c) < +Inf
300 <= (1c + 1f + 1i) <= 300
0 <= (1d) < +Inf
0 <= (1e) < +Inf
0 <= (1f) < +Inf
0 <= (1g) < +Inf
0 <= (1h) < +Inf
0 <= (1i) < +Inf

```

0	a	150	0
1	b	0	2
2	c	0	5
3	d	80	0
4	e	140	0
5	f	300	0
6	g	0	0
7	h	0	2
8	i	0	1

	NYC	ATL	LA
A	`a	`b	`c
B	`d	`e	`f
C	`g	`h	`i

# Math Prog Step 7

	NYC	ATL	LA
A	150	0	0
B	80	140	300
C	0	0	0