Trees

Today's Topics

- Trees
- Kinds of trees branching factor
- -functions over trees
- -patterns of recursion the fold for trees
- -Arithmetic expressions
- -Infinite trees

Trees

- Trees are important data structures in computer science
- Trees have interesting properties
 - They usually are finite, but unbounded in size
 - Sometimes contain other types inside
 - Sometimes the things contained are polymorphic
 - differing "branching factors"
 - different kinds of leaf and branching nodes
- Lots of interesting things can be modeled by trees
 - lists (linear branching)
 - arithmetic expressions
 - parse trees (for languages)
- In a lazy language it is possible to have infinite trees

Examples

```
data List a = Nil | MkList a (List a)
data Tree a = Leaf a | Branch (Tree a) (Tree a)
data IntegerTree = IntLeaf Integer
                   IntBranch IntegerTree IntegerTree
data SimpleTree = SLeaf
                 | SBranch SimpleTree SimpleTree
data InternalTree a = ILeaf
                      IBranch a (InternalTree a)
                                 (InternalTree a)
data FancyTree a b = FLeaf a
                     FBranch b (FancyTree a b)
                                (FancyTree a b)
```

Cse536 Functional Programming

Match up the trees



Functions on Trees

Transforming one kind of tree into another

Collecting the items in a tree

fringe :: Tree a -> [a]
fringe (Leaf x) = [x]
fringe (Branch t1 t2) = fringe t1 ++ fringe t2

what kind of information is lost using fringe?

More functions

Capture the pattern of recursion

foldTree :: $(a \rightarrow a \rightarrow a) \rightarrow (b \rightarrow a) \rightarrow Tree b \rightarrow a$ foldTree b l (Leaf x) = l x foldTree b l (Branch t1 t2) = b (foldTree b l t1) (foldTree b l t2)

mapTree2 f = foldTree Branch (Leaf . f)

```
fringe2 = foldTree (++) (\land x \rightarrow [x])
```

```
treeSize2 = foldTree (+) (const 1)
```

```
treeHeight2 = foldTree (\ x \ y \ -> 1 + max \ x \ y)
(const 0)
```

Flattening Trees

data Tree a = Leaf a | Branch (Tree a) (Tree a) flatten :: Tree a -> [a] flatten (Leaf x) = [x] flatten (Branch x y) = flatten x ++ flatten y

What is the complexity of flattening a deep fully filled out tree?

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Flattening with accumulating parameter

data Tree a

= Leaf a | Branch (Tree a) (Tree a)

```
flatten :: Tree a -> [a]
Flatten t = flat t []
```

```
flat (Leaf x) xs = x:xs
Flat (Branch a b) = flat a (flat b xs)
```

Arithmetic Expressons

data Expr2 = C2 Float

- | Add2 Expr2 Expr2
- Sub2 Expr2 Expr2
- Mul2 Expr2 Expr2
- | Div2 Expr2 Expr2
- using infix constructor functions/

data Expr = C Float

- Expr :+ Expr
- | Expr :- Expr
- Expr :* Expr
 - Expr :/ Expr

Infix constructor operators start with a colon (:), just like constructor functions start with an upper case letter

Example uses

```
e1 = (C 10 :+ (C 8 :/ C 2)) :* (C 7 :- C 4)
```

evaluate	:: Expr ->	Float			
evaluate	(C x) = x				
evaluate	(e1 :+ e2)	= evaluate	e1 +	evaluate	e2
evaluate	(e1 :- e2)	= evaluate	e1 -	evaluate	e2
evaluate	(e1 :* e2)	= evaluate	e1 *	evaluate	e2
evaluate	(e1 :/ e2)	= evaluate	e1 /	evaluate	e2

Main> evaluate e1 42.0

Infinite Trees

Can we make an Expr tree that represents the infinite expression: 1 + 2 + 3 + 4

```
sumFromN n = C n :+ (sumFromN (n+1))
sumAll = sumFromN 1
```

```
add1 (C n) = C (n+1)
add1 (x :+ y) = add1 x :+ add1 y
add1 (x :- y) = add1 x :- add1 y
add1 (x :* y) = add1 x :* add1 y
add1 (x :/ y) = add1 x :/ add1 y
sumAll2 = C 1 :+ (add1 sumAll2)
```

Observing Infinite Trees

 We can observe an infinite tree by printing a finite prefix of it. We need a take-like function for trees.

```
showE 0 _ = "..."

showE n (C m) = show m

showE n (x :+ y) = "(" ++ (showE (n-1) x) ++ "+"

++ (showE (n-1) y) ++ ")"
```

```
Main> showE 5 sumAll2
"(1.0+(2.0+(3.0+(4.0+(...+...)))))"
Main> showE 5 sumAll
"(1.0+(2.0+(3.0+(4.0+(...+...)))))"
```