

# Trees

## •Today's Topics

- Trees
- Kinds of trees - branching factor
- functions over trees
- patterns of recursion - the fold for trees
- Arithmetic expressions
- Infinite trees

# Trees

- **Trees are important data structures in computer science**
- **Trees have interesting properties**
  - They usually are finite, but unbounded in size
  - Sometimes contain other types inside
  - Sometimes the things contained are polymorphic
  - differing “branching factors”
  - different kinds of leaf and branching nodes
- **Lots of interesting things can be modeled by trees**
  - lists (linear branching)
  - arithmetic expressions
  - parse trees (for languages)
- **In a lazy language it is possible to have infinite trees**

# Examples

```
data List a = Nil | MkList a (List a)
```

```
data Tree a = Leaf a | Branch (Tree a) (Tree a)
```

```
data IntegerTree = IntLeaf Integer  
                | IntBranch IntegerTree IntegerTree
```

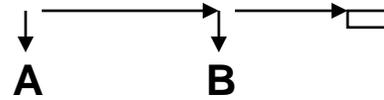
```
data SimpleTree = SLeaf  
               | SBranch SimpleTree SimpleTree
```

```
data InternalTree a = ILeaf  
                   | IBranch a (InternalTree a)  
                               (InternalTree a)
```

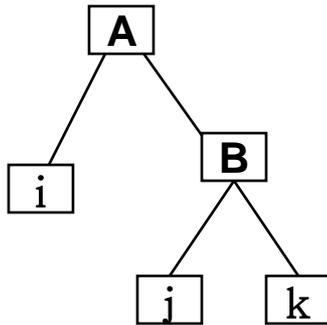
```
data FancyTree a b = FLeaf a  
                   | FBranch b (FancyTree a b)  
                               (FancyTree a b)
```

# Match up the trees

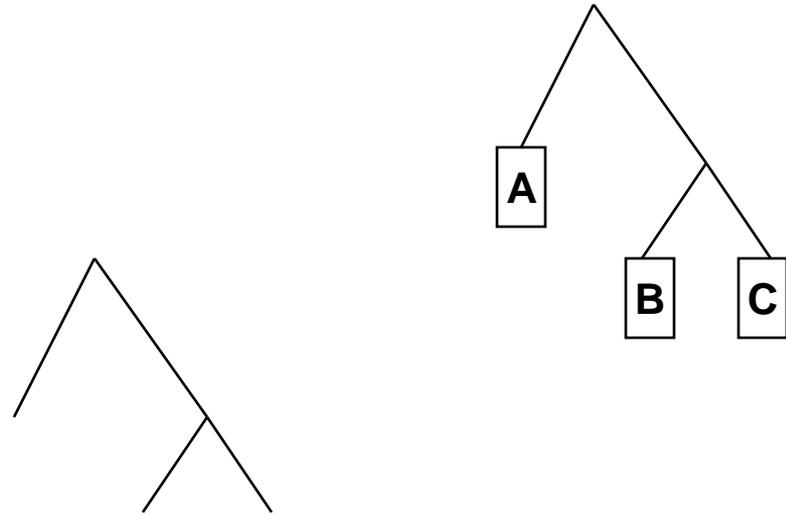
- IntegerTree



- Tree

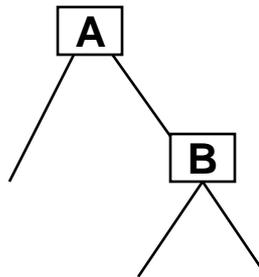


- SimpleTree

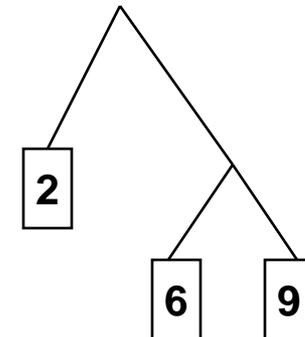


- List

- InternalTree



- FancyTree



# Functions on Trees

- Transforming one kind of tree into another

```
mapTree :: (a->b) -> Tree a -> Tree b
mapTree f (Leaf x)           = Leaf (f x)
mapTree f (Branch t1 t2) = Branch (mapTree f t1)
                               (mapTree f t2)
```

- Collecting the items in a tree

```
fringe :: Tree a -> [a]
fringe (Leaf x)           = [x]
fringe (Branch t1 t2) = fringe t1 ++ fringe t2
```

- what kind of information is lost using `fringe`?

# More functions

```
treeSize                :: Tree a -> Integer
treeSize (Leaf x)       = 1
treeSize (Branch t1 t2) = treeSize t1 + treeSize t2
```

```
treeHeight              :: Tree a -> Integer
treeHeight (Leaf x)     = 0
treeHeight (Branch t1 t2) = 1 + max (treeHeight t1)
                                   (treeHeight t2)
```

# Capture the pattern of recursion

```
foldTree :: (a -> a -> a) -> (b -> a) -> Tree b -> a
foldTree bf lf (Leaf x)           = lf x
foldTree bf lf (Branch t1 t2) =
    bf (foldTree bf lf t1) (foldTree bf lf t2)
```

```
mapTree2 f = foldTree Branch (Leaf . f)
```

```
fringe2 = foldTree (++) (\ x -> [x])
```

```
treeSize2 = foldTree (+) (const 1)
```

```
treeHeight2 = foldTree (\ x y -> 1 + max x y)
                (const 0)
```

# Flattening Trees

```
data Tree a
  = Leaf a | Branch (Tree a) (Tree a)

flatten :: Tree a -> [a]
flatten (Leaf x) = [x]
flatten (Branch x y) = flatten x ++ flatten y
```

**What is the complexity of flattening a deep fully filled out tree?**

# Flattening with accumulating parameter

```
data Tree a
  = Leaf a | Branch (Tree a) (Tree a)
```

```
flatten :: Tree a -> [a]
```

```
flatten t = flat t []
```

```
flat (Leaf x) xs = x:xs
```

```
flat (Branch a b) xs = flat a (flat b xs)
```

# Arithmetic Expressions

```
data Expr2 = C2 Float
           | Add2 Expr2 Expr2
           | Sub2 Expr2 Expr2
           | Mul2 Expr2 Expr2
           | Div2 Expr2 Expr2
```

- using infix constructor functions

```
data Expr = C Float
          | Expr :+: Expr
          | Expr :- Expr
          | Expr :* Expr
          | Expr :/ Expr
```

Infix constructor operators start with a colon (:), just like constructor functions start with an upper case letter

# Example uses

```
e1 = (C 10 :+: (C 8 :/ C 2)) :* (C 7 :- C 4)
```

```
evaluate :: Expr -> Float
```

```
evaluate (C x) = x
```

```
evaluate (e1 :+: e2) = evaluate e1 + evaluate e2
```

```
evaluate (e1 :- e2) = evaluate e1 - evaluate e2
```

```
evaluate (e1 :* e2) = evaluate e1 * evaluate e2
```

```
evaluate (e1 :/ e2) = evaluate e1 / evaluate e2
```

```
Main> evaluate e1
```

```
42.0
```

# Infinite Trees

- Can we make an Expr tree that represents the infinite expression:  $1 + 2 + 3 + 4 \dots$

```
sumFromN n = C n :+: (sumFromN (n+1))
sumAll = sumFromN 1
```

```
add1 (C n) = C (n+1)
add1 (x :+: y) = add1 x :+: add1 y
add1 (x :- y) = add1 x :- add1 y
add1 (x :* y) = add1 x :* add1 y
add1 (x :/ y) = add1 x :/ add1 y

sumAll2 = C 1 :+: (add1 sumAll2)
```

# Observing Infinite Trees

- We can observe an infinite tree by printing a finite prefix of it. We need a `take`-like function for trees.

```
showE 0 _ = "..."  
showE n (C m) = show m  
showE n (x :+: y) = "(" ++ (showE (n-1) x) ++ "+"  
                        ++ (showE (n-1) y) ++ ")"
```

```
Main> showE 5 sumAll2  
"(1.0+(2.0+(3.0+(4.0+(...+...)))))"  
  
Main> showE 5 sumAll  
"(1.0+(2.0+(3.0+(4.0+(...+...)))))"
```