

Context Free Pumping Lemma

Some languages are not context free!

Sipser pages 125 - 129

CFL Pumping Lemma

- A *CFL pump* consists of two non-overlapping substrings that can be pumped simultaneously while staying in the language.
- Precisely, two substrings u and v constitute a CFL pump for a string w of L when
 1. $uv \neq \varepsilon$ ($|uv| > 0$, which means that at least one of u or v is not empty)
 2. And we can write $w = xuyvz$
so that for every $i \geq 0$
 3. $xu^i y v^i z \in L$

Pumping Lemma

- Let L be a CFL. Then there exists a number n (depending on L) such that every string w in L of length greater than n contains a CFL pump.
- Moreover, there exists a CFL pump such that (with the notation as above), $|uyv| \leq n$.
- For example, take $L = \{0^i1^i \mid i \geq 0\}$: there are no (RE) pumps in any of its strings, but there are plenty of CFL pumps.

The pumping Lemma Game

- We want to prove L is not context-free. For a proof, it suffices to give a winning strategy for this game.
 1. The demon first plays n .
 2. We respond with $w \in L$ such that $|w| \geq n$.
 3. The demon factors w into five substrings, $w=xuyvz$, with the proviso that $uv \neq \varepsilon$ and $|uyv| \leq n$
 4. Finally, we play an integer $i \geq 0$, and we win if $xu^iyv^iz \notin L$.

Example 1

- We prove that $L = \{0^i 1^i 2^i \mid i \geq 0\}$ is not context-free.
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- In response to the demon's n , we play $w = 0^n 1^n 2^n$.
- The middle segment uyv of the demon's factorization of $w = xuyvz$, cannot have an occurrence of both 0 and 2 (because we can assume $|uyv| \leq n$).
- Suppose 2 does not occur in uyv (the other case is similar).
 1. We play $i = 0$.
 2. Then the total number of 0's and 1's in $w_0 = xyz$ will be smaller than $2n$,
 3. while the number of 2's in w_0 will be n .
 4. Thus, $w_0 \notin L$.

Example 2

- Let L be the set of all strings over $\{0,1\}$ whose length is a perfect square.
 1. The demon plays n
 2. We respond with $w = 0^{n^2}$
 3. The demon plays a factorization $0^{n^2} = xuyvz$ with $1 \leq |uyv| \leq n$.
 4. We play $i=2$.
 5. The length of the resulting string $w_2 = xu^2yv^2z$ is between n^2+1 and n^2+n .
 6. In that interval, there are no perfect squares, so $w_2 \notin L$.

Proof of the pumping lemma

- Strategy in several steps
 1. Define fanout
 2. Define height yield
 3. Prove a lemma about height yield
 4. Apply the lemma to prove pumping lemma

Fanout – max rhs length

- Let $\text{fanout}(G)$ denote the maximal length of the rhs of any production in the grammar G .
- E.g. For the Grammar
 - $S \rightarrow S S$
 - $S \rightarrow (S)$
 - $S \rightarrow \varepsilon$
- The fanout is 3

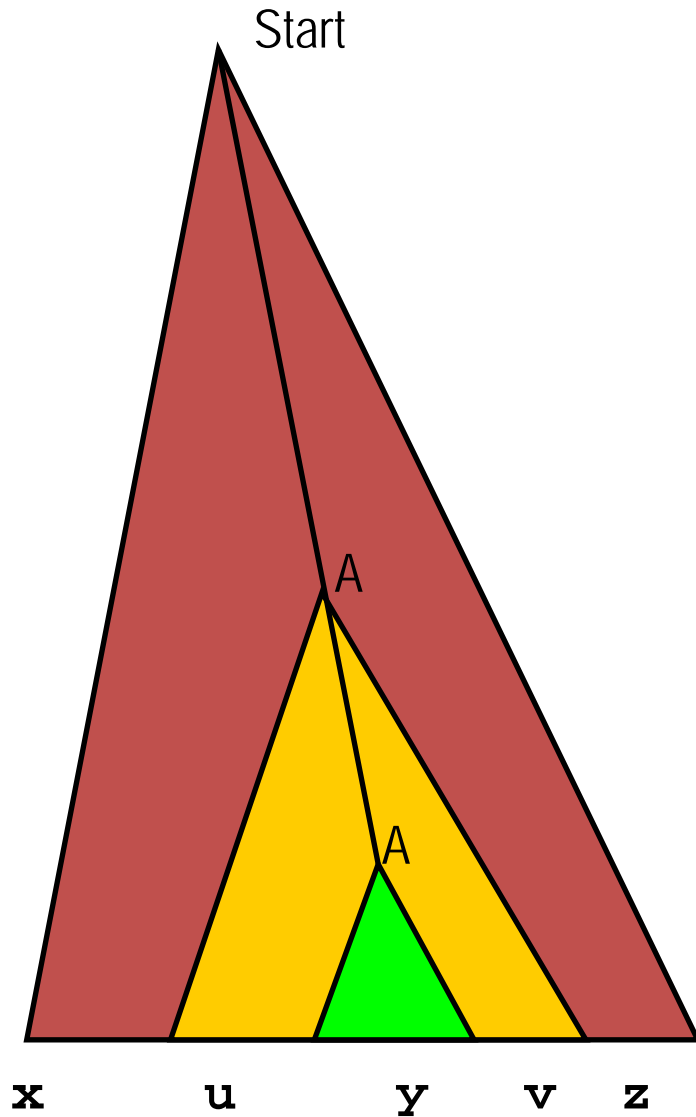
Height Yield

- The proof of Pumping Lemma depends on this simple fact about parse trees.
- The *height* of a tree is the maximal length of any path from the root to any leaf.
- The yield of a parse tree is the string it represents (the terminals from a left-to-right in-order walk)
- **Lemma.** If a parse tree of G has height h , and the $\text{fanout}(G)=b$, then its yield has size at most b^h
 - Conversely if a string has length $> b^h$ (i.e. it has length $b^h + 1$ or more) then its parse tree is at least $h+1$ deep
- **Proof.** Induction on h
- qed

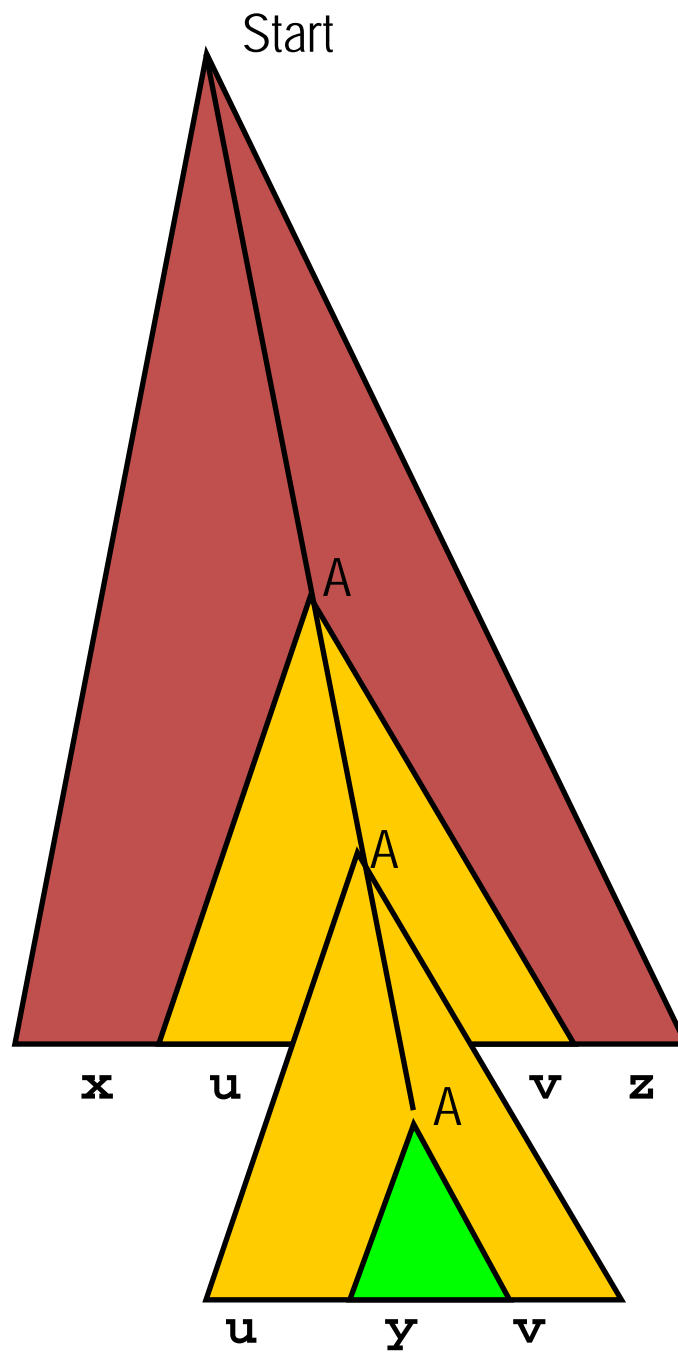
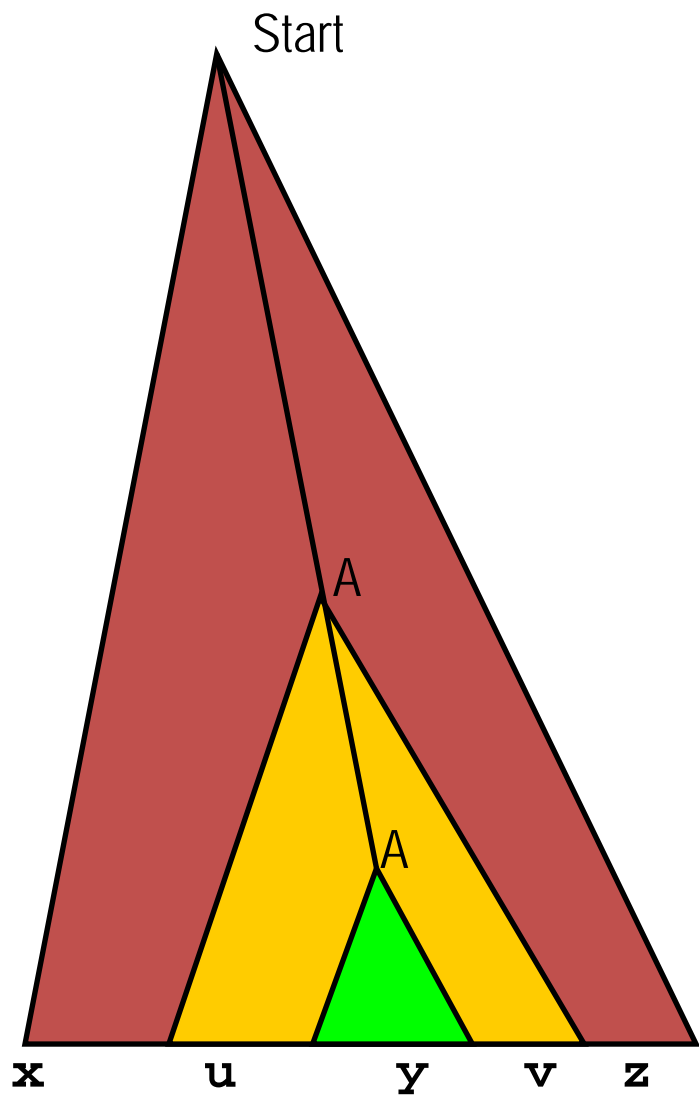
The actual Proof

- The constant n for the grammar G is $\text{fanout}(G)^{|V|}$ where V is the set of variables of G .
- Suppose $w \in L(G)$ and $|w| \geq n$.
- Take a parse tree of w with the smallest possible number of nodes.
- By the Height-Yield Lemma, any parse tree of w must have height $\geq |V|$.
- Therefore, there must be two occurrences of the same variable on a path from root to a leaf.
- Consider the last two occurrences of the same variable (say A) on that path.
- They determine a factorization of the yield $w = xuyvz$ as in the picture on the next slide

Diagram



- We have
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- $S \Rightarrow^* xAz$
- $A \Rightarrow^* uAv$
- $A \Rightarrow^* y$
- so clearly $S \Rightarrow^* xu^i y v^i z$ for any $i \geq 0$.



Proof continued

- We also need to check that $uv \neq \varepsilon$. Indeed, if $uv = \varepsilon$, we can get a smaller parse tree for the same w by ignoring the productions “between the two A s”. But we have chosen the smallest possible parse tree for w ! Which leads to a Contradiction.
- Finally, we need to check that $|uyv| \leq n$. This follows from the Height-Yield Lemma because the nodes on our chosen path from the first depicted occurrence of A , onward, are labeled with necessarily distinct variables.
- qed