



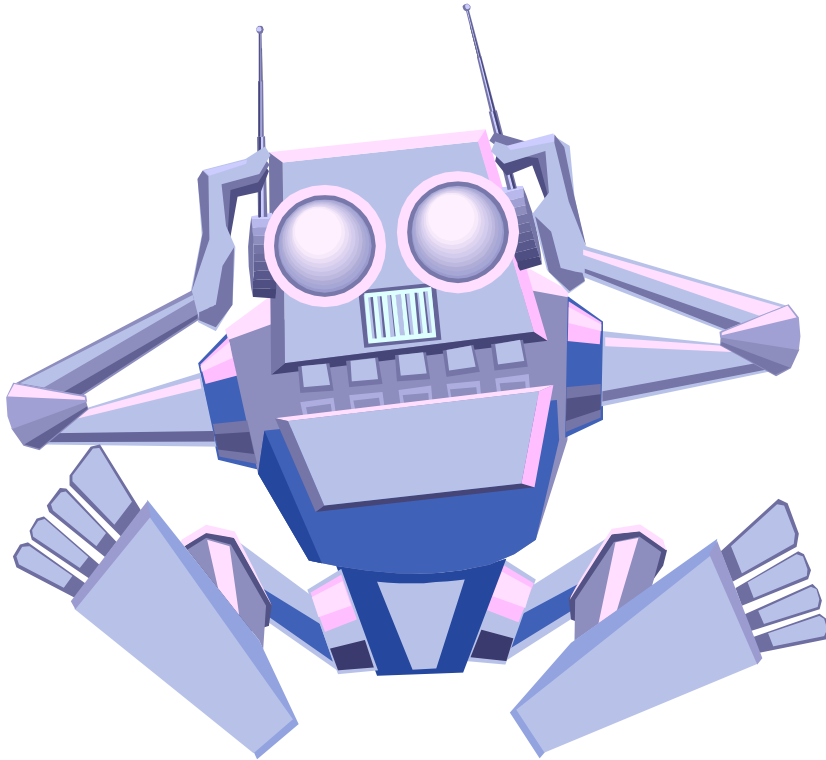
10001

Cellular Automata III



Based mostly on lectures by Dr.
Richard Spillman

OUTLINE



What an advanced class!

- ➔ Random Number Generation
- ➔ A Cellular Automata Cipher System
- ➔ Computation on Cellular Automata

Random Number Generation

- ⇒ Random numbers are required for a **wide range of applications**
 - ⇒ Cryptology
 - ⇒ Testing
 - ⇒ Modeling and Simulation
 - ⇒ Genetic Algorithms
 - ⇒ . . .
- ⇒ Yet, **true random numbers** are very difficult to find
 - ⇒ Computer based random number generators are really **pseudo-random** number generators because they eventually repeat

Current Random Number Generators

- ⇒ There are **two typical approaches** to random number generation
 - ⇒ Use of a mathematical relationship
 - ⇒ Use of a linear feedback shift register (LFSR)
- ⇒ A common mathematical relationship is of the form **$x' = (ax + b) \bmod n$**

$$x' = (11x + 17) \bmod 61$$

Seed

Random Numbers

13

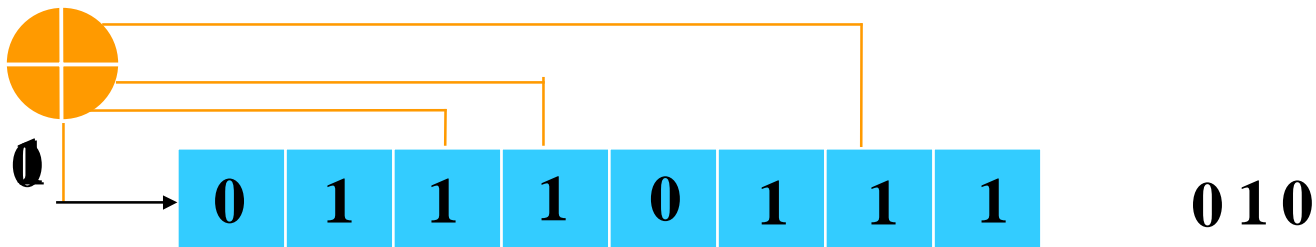
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Linear Feedback Shift Register

- ⇒ A LFSR is a hardware random number generator
 - ⇒ A shift register holds a data word and can shift it to the left or right one bit position on each clock pulse



Add feedback

Load a seed

A CA Random Generator

- ⇒ The study of random bit generation in cellular automata began with the work of Steven Wolfram
 - ⇒ He first suggested a 1-d binary cellular automata with a $k=2$ neighborhood and rule 30

$$q_i(t+1) = q_{i-1}(t) \text{ XOR } (q_i(t) \text{ OR } q_{i+1}(t))$$

Example

⇒ Consider a simple 1-d CA with an initial seed:

Select one cell to provide
a random binary bit

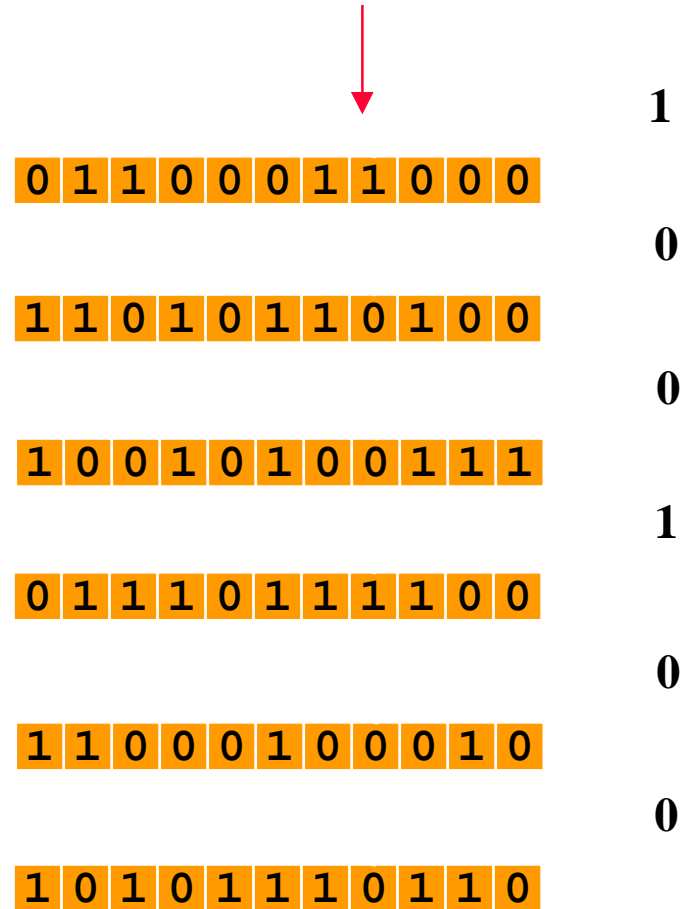
$$q_i(t+1) = q_{i-1}(t) \text{ XOR } (q_i(t) \text{ OR } q_{i+1}(t))$$

Questions:

When will the patterns
begin to repeat?

What is the best seed?

Are there better rules?



GOAL

- ⇒ The goal of a random number generator is to produce **as long** a sequence of random numbers (bits) **as possible** before the pattern begins to repeat itself
 - ⇒ We are looking for the **set of CA parameters** that produce a **maximum length cycle**
 - ⇒ Wolfram found that a **single nonzero element seed** produced the maximum length cycle for **any rule**
 - ⇒ Wolfram also found that the longest cycles were produced by **rule 30**

Experimental Data

⇒ Maximum cycle lengths for an n-bit CA using **rule 30** and a **single nonzero bit seed**:

N	Maximum Cycle Length
4	8
7	63
11	154
14	1428
24	185,040
32	2,002,272

Hybrid Cellular Automata (HCA)

- ⇒ **Hybrid Cellular Automata** (also called non-uniform cellular automata) function the same way as uniform cellular automata except that the cellular rules need not be identical for each cell
 - ⇒ One cell may follow **rule 90** and another cell may follow **rule 30**
 - ⇒ This results in a **new level of cellular automata behavior**

Random Numbers w/HCAs

- ⇒ It turns out that HCAs (hybrid cellular automata) produce **maximal-length binary sequences** when the **right rule set** is used
 - ⇒ One of the better choices seems to be a **combination of rule 90 and rule 150**

RULE 150

$$q_i(t+1) = q_{i-1}(t) \text{ XOR } q_i(t) \text{ XOR } q_{i+1}(t)$$

HCA Operation

- ⇒ If more than one rule is active in a cellular automata then which rule is assigned to which cell?
 - ⇒ This is usually solved experimentally (or it could be solved by a **genetic algorithm**)
 - ⇒ For example, it has been found that a four element hybrid cellular automata with rules distributed as **90 150 90 150** produces a maximal length cycle
 - ⇒ The notation for this rule distribution is 0 1 0 1 where rule 150 is 1 and rule 90 is 0

Comparison

⇒ The best rule distribution has been determined for some HCAs:

N	Maximum Cycle Length	
	Uniform	HCA
4	8	15
7	63	127
11	154	2,047
14	1428	16,383
24	185,040	16,777,215
32	2,002,272	$2^{32}-1$

CA-Based Cipher Systems

- ⇒ The need for **security** and **privacy** has become a given in today's internet connected world.
- ⇒ One of the major tools for security is the use of **encryption algorithms** to hide data and messages
- ⇒ Cellular automata can play a significant role in the construction of **new fast, secure and efficient ciphers**

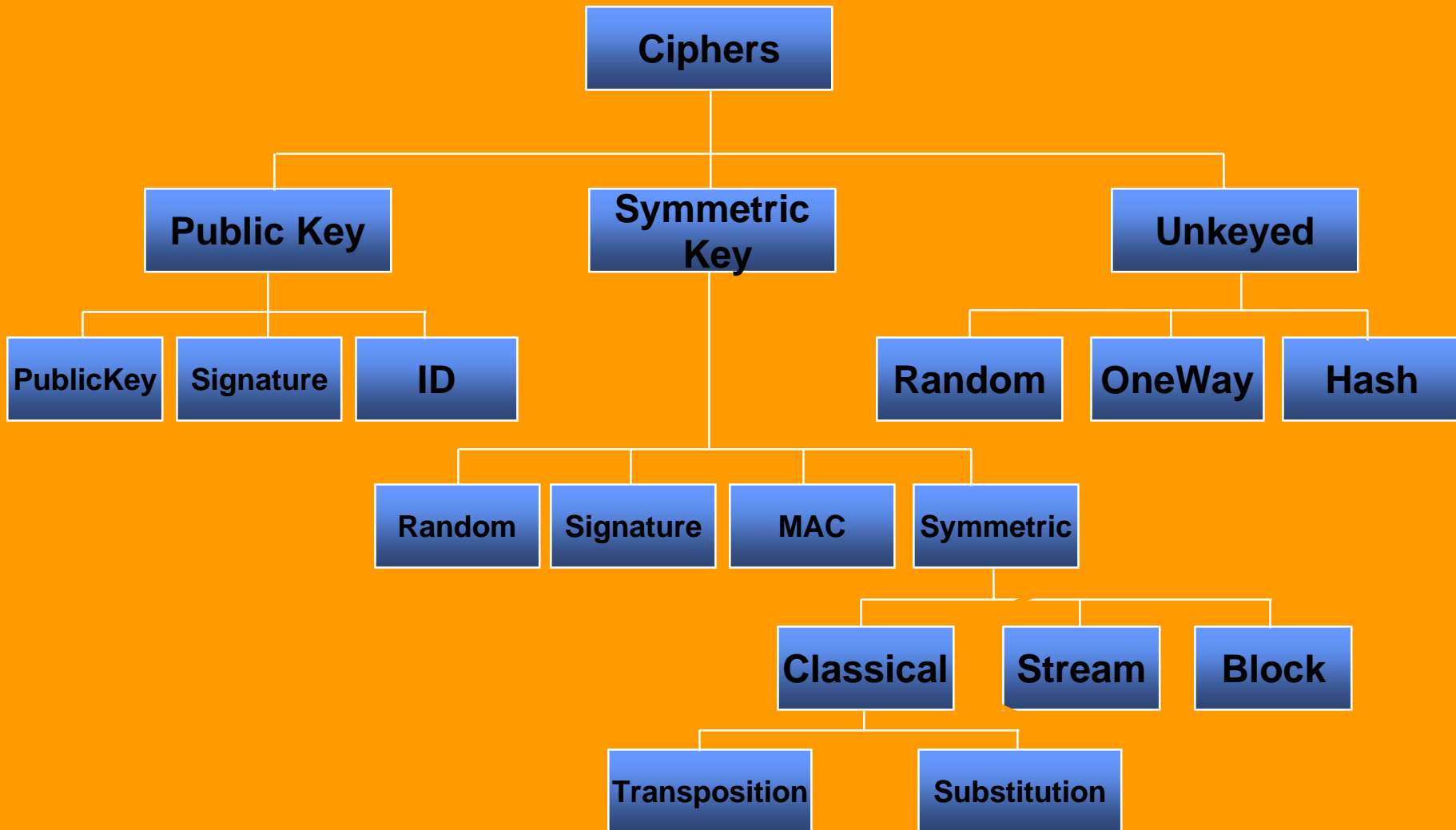
Cipher Systems

- ⇒ The process of disguising a message in such a way as to hide its substance is called *encryption*
 - ⇒ a message is called *plaintext*
 - ⇒ the encrypted message is called *ciphertext*
 - ⇒ the process of turning ciphertext back into plaintext is called *decryption*
- ⇒ The art and science of keeping messages secure is called *cryptology*
 - ⇒ *cryptanalysis* is the *art and science* of breaking ciphertext

A Good Cipher

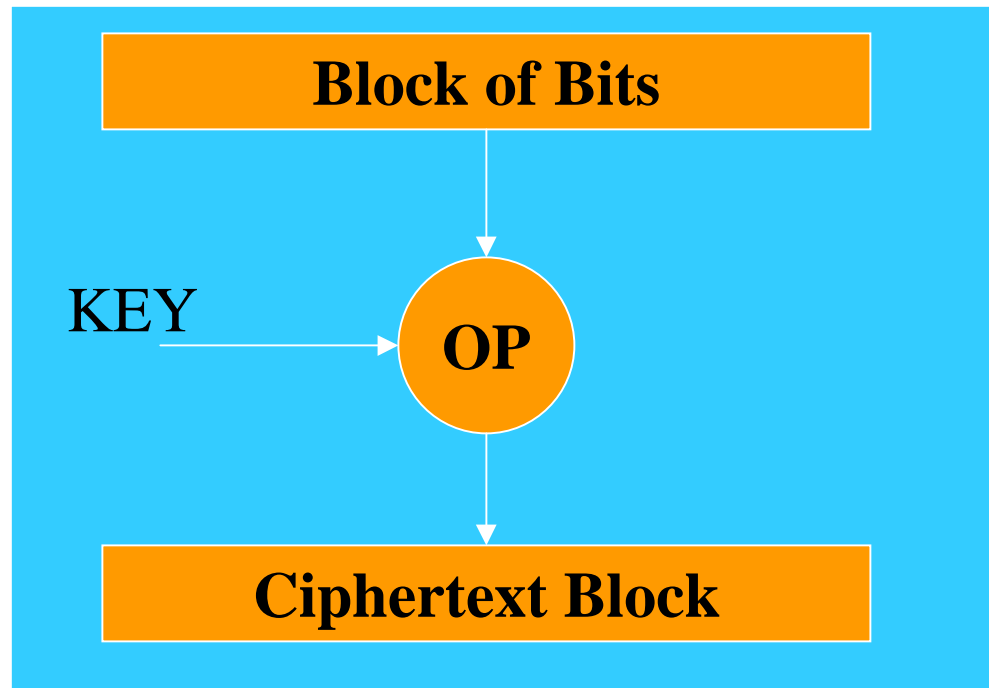
- ⇒ Enciphering and deciphering should be efficient for **all keys**
- it should not take forever to get message.
- ⇒ **Easy to use.** The problem with hard to use cryptosystems is that mistakes tend to be made
- ⇒ The strength of the system should not lie in the **secrecy of your algorithms.**
 - ⇒ The strength of the system should depend the **secrecy of your key.**

Cipher Classification



Block Ciphers

- ⇒ Today's most widely used ciphers are in the class of Block Ciphers
 - ⇒ Define a block of computer bits which represent several characters
 - ⇒ Encipher the complete block at one time



Block Cipher Methods

- ⇒ Plaintext is divided into **fixed length blocks** M_1, M_2, \dots, M_m and each block is transformed into ciphertext so the entire ciphertext is given by C_1, C_2, \dots, C_m
- ⇒ Block size should be **large**
 - ⇒ usually it is 64 bits

CA Block Cipher

- ⇒ The problem with any random number generator including a CA-based random number generator is that eventually **it will cycle back to the beginning**
- ⇒ Hence, the design of a random number generator required finding a rule which made the **cycle as long as possible**
- ⇒ For a cipher, we want to take advantage of the cycle feature

CA Cycles

- ⇒ For a cipher, **find a CA rule** that:
 - ⇒ Produces a short cycle of length $2r$
 - ⇒ The **rule becomes the key**
- ⇒ There are three rules called **fundamental transformations** that meet the necessary requirements for a block cipher – rules **51, 153, and 195**

$$\mathbf{195: } q_i(t+1) = \mathbf{NOT}(q_{i-1}(t) \mathbf{XOR} q_i(t))$$

$$\mathbf{153: } q_i(t+1) = \mathbf{NOT}(q_{i+1}(t) \mathbf{XOR} q_i(t))$$

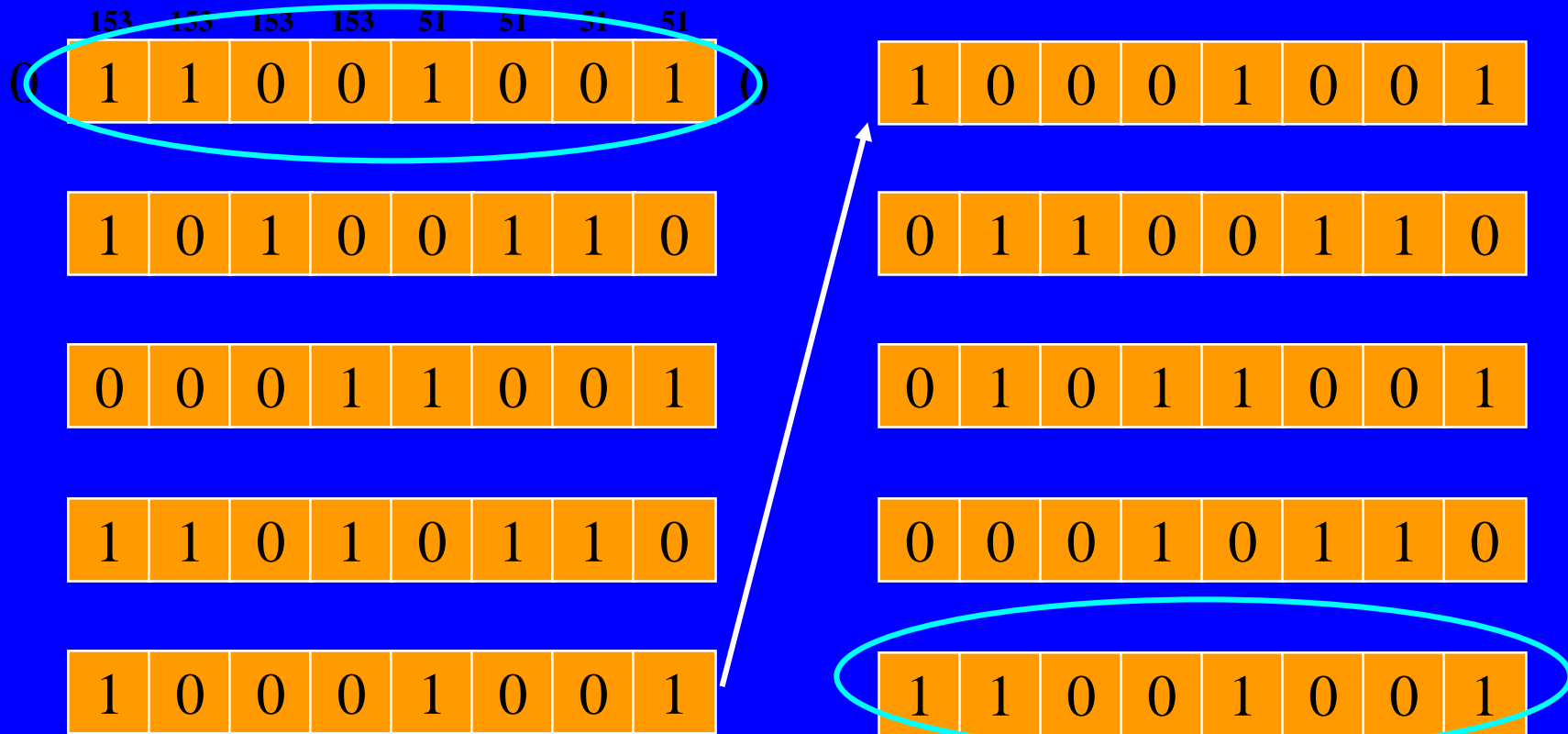
$$\mathbf{51: } q_i(t+1) = \mathbf{NOT}(q_i(t))$$

Hybrid CA

- ⇒ A cellular automata cipher requires a **hybrid system** with some combination of the 3 fundamental rules.
- ⇒ For example, consider a null boundary, 8-bit CA with rules applied to the cells in this order:
 - ⇒ (153,153,153,153,51,51,51,51)
 - ⇒ This has a cycle of length 8, so run it for 4 clock ticks and send the result
 - ⇒ The result is decoded by completing the cycle by running the ciphertext for an additional 4 clock ticks

Example of a hybrid system with some combination of the 3 fundamental rules.

- Start with an 8-cell CA with null boundaries
 - That is, each end is connected to 0 instead of each other



Key Space

- ⇒ For a CA Block Cipher, the **key space** is the space of fundamental transformations
- ⇒ For the example on the prior slide **other possible transformation patterns** include:
 - ⇒ (195, 195, 195, 195, 51, 51, 51, 51)
 - ⇒ (51, 51, 153, 153, 153, 153, 51, 51)
 - ⇒ (51, 153, 153, 153, 153, 153, 153, 51)

Stream Cipher

- ⇒ Consider the plaintext as a sequence of bits
- ⇒ Generate a random sequence of bits for the key
- ⇒ Form the ciphertext by a bit by bit XOR of the *plaintext with the ciphertext*

plaintext:	1	0	0	1	0	1	1	0	0	1	1	1	0	1	0	1	0	0	1
key:	0	0	1	1	1	0	1	0	1	0	0	1	1	0	0	1	1	0	0
ciphertext:	1	0	1	0	1	1	0	0	1	1	1	0	1	1	0	0	1	0	1
key:	0	0	1	1	1	0	1	0	1	0	0	1	1	0	0	1	1	0	0
plaintext:	1	0	0	1	0	1	1	0	0	1	1	1	0	1	0	1	0	0	1

Problem: How do we recover the plaintext from knowledge of the ciphertext and key?

Problem

- ⇒ A **short sequence** of key bits would be easy to remember but not very secure
- ⇒ A **long sequence** of key bits would be secure but hard to remember
- ⇒ **PROBLEM:** How can we generate a long random-appearing sequence of 0's and 1's in way that will insure that everyone who should have access to the plaintext are **able to generate the key** when needed?
- ⇒ **ANSWER:** Construct a random bit generator

Stream Cipher Key

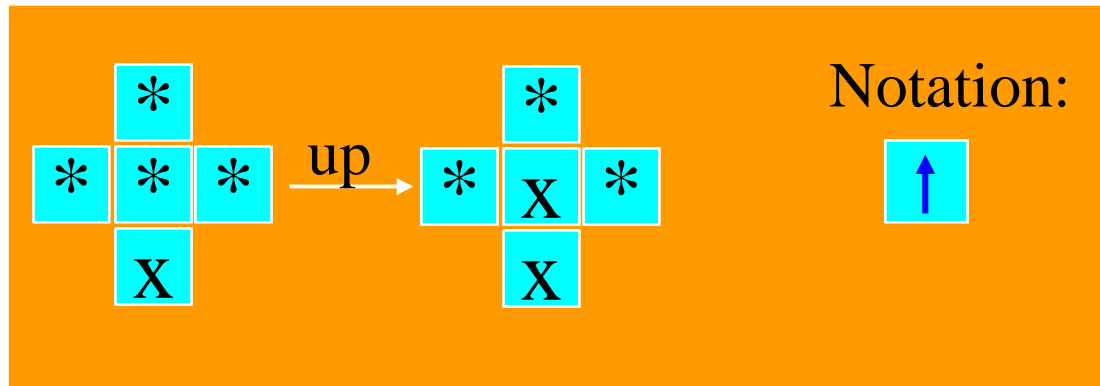
- ⇒ For a CA-based stream cipher, the two parties must **have the same CA** with the **same initial condition**.
- ⇒ Hence, the key is the initial state of the CA
- ⇒ With the key, the person who receives the ciphertext can **start the CA random number generator** and **produce the correct key stream**

Computation with CAs

- ⇒ It is possible to design a complete circuit on a cellular automata
 - ⇒ The equivalent of a digital circuit embedded in the cellular automata in the form of rules
- ⇒ Create a large 2-d binary, von Neumann neighborhood hybrid cellular automata
 - ⇒ It will have pathways for signal transmission (wires)
 - ⇒ It will have functional cells that perform NANDs, XORs, ...

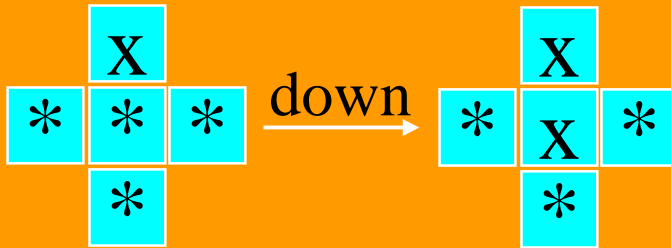
Signal Pathways

- ⇒ A “wire” on our cellular automata consists of a path of connected **propagation cells**
 - ⇒ Each of these cells implements a propagation rule
 - ⇒ There are four propagation rules: **right, left, up, down**
 - ⇒ For cells which are not part of the circuit there is a **NC** (no change) cell

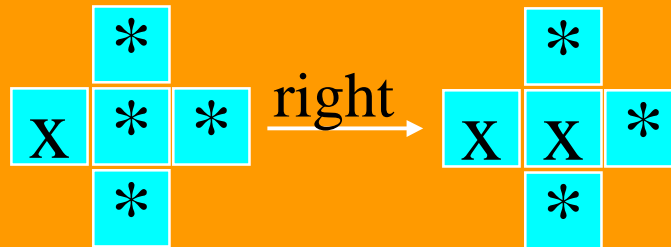


Signal Propagation Rules

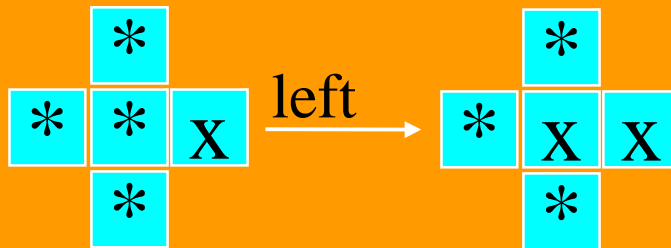
⇒ The other propagation rules are given by:



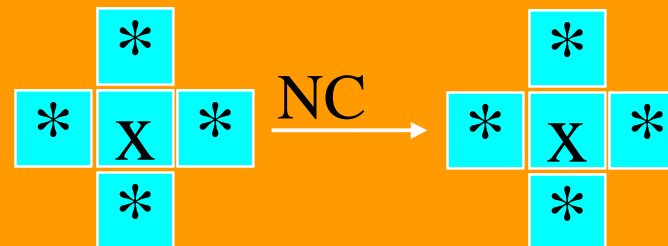
Notation:



Notation:



Notation:

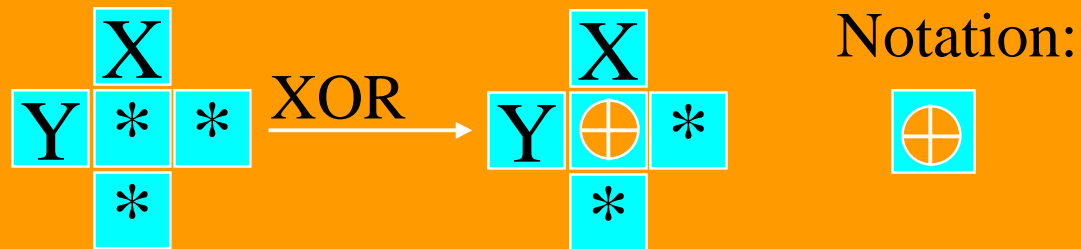
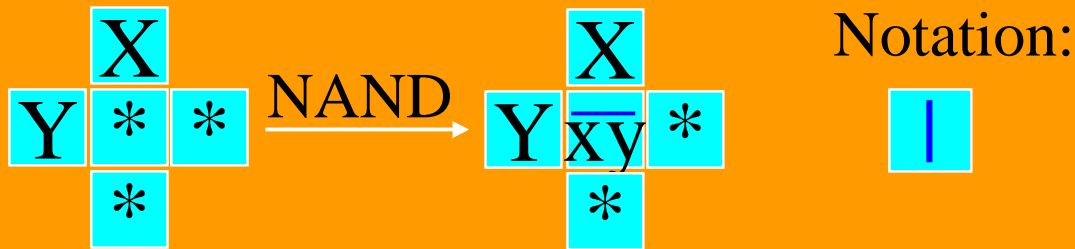


Notation:



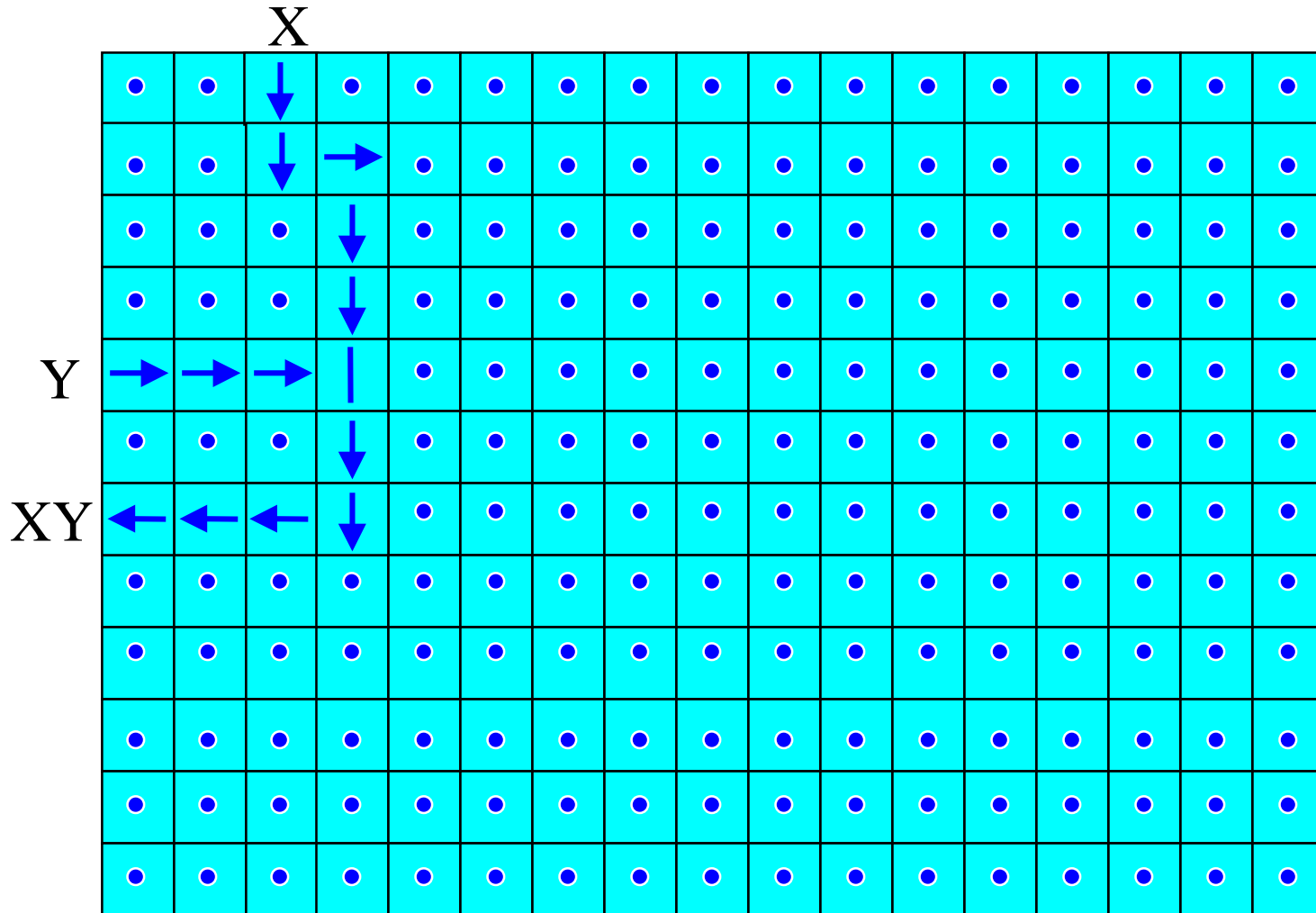
Logic Operations Rules

⇒ Rules for both **NAND** and **XOR** are given by:

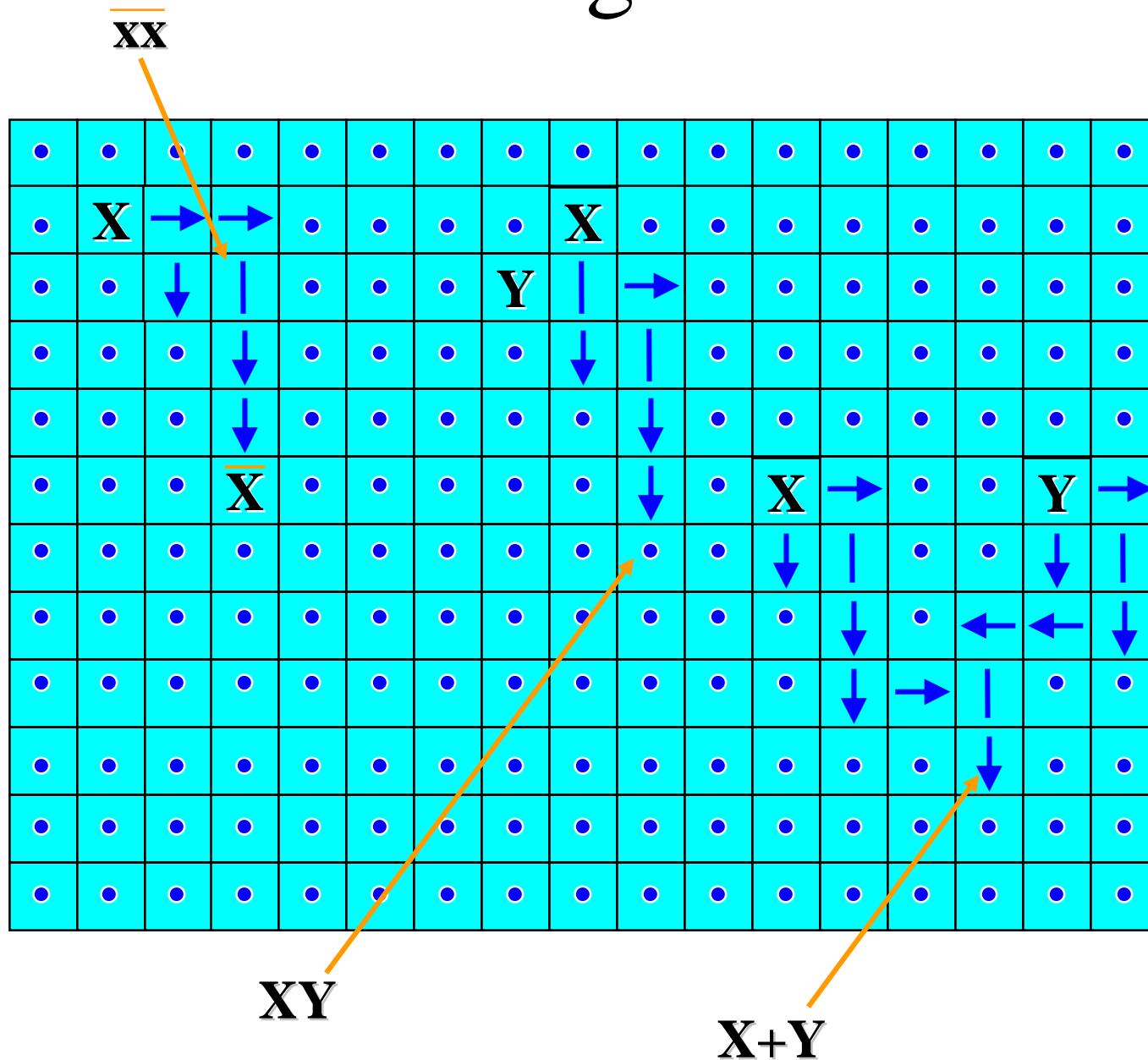


Example

⇒ Given a 2-d cellular automata:

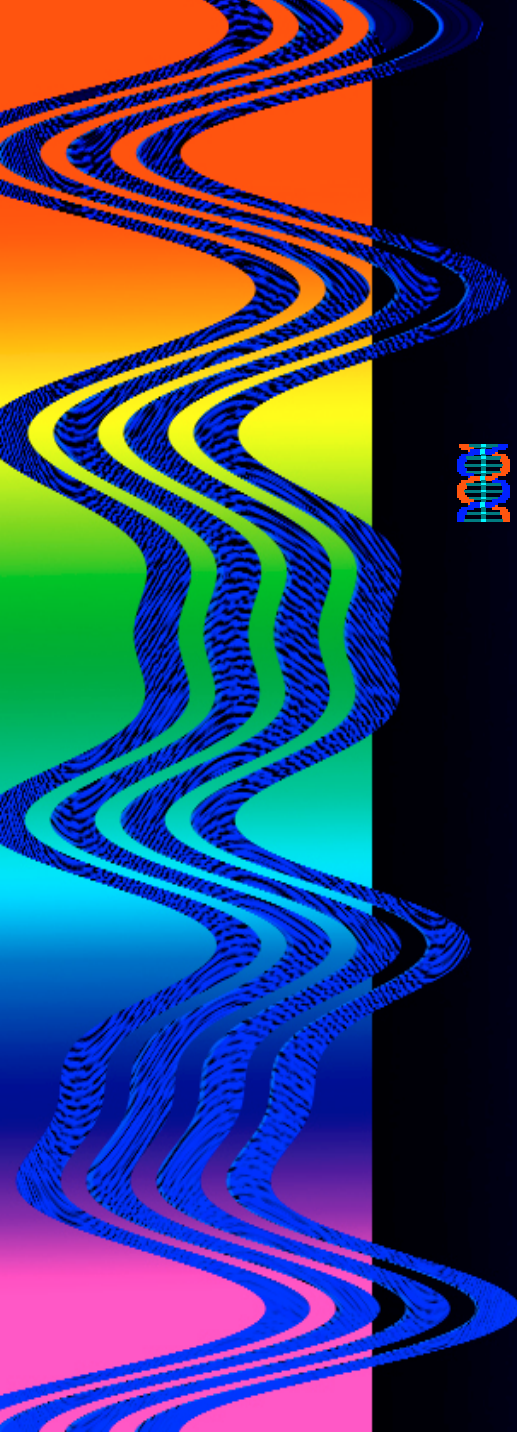


Other Logic Functions



Possible Homeworks

- ⇒ You might want to look at the application of CAs to random number generation
 - ⇒ For a possible senior project, consider other pairs of rules or even 3 or more rules in an HCA
 - ⇒ For a possible advanced homework project, build a genetic algorithm to create a 2 or 3 rule HCA
- ⇒ Construct a simple CA cipher system
- ⇒ Design a simple digital circuit on a CA

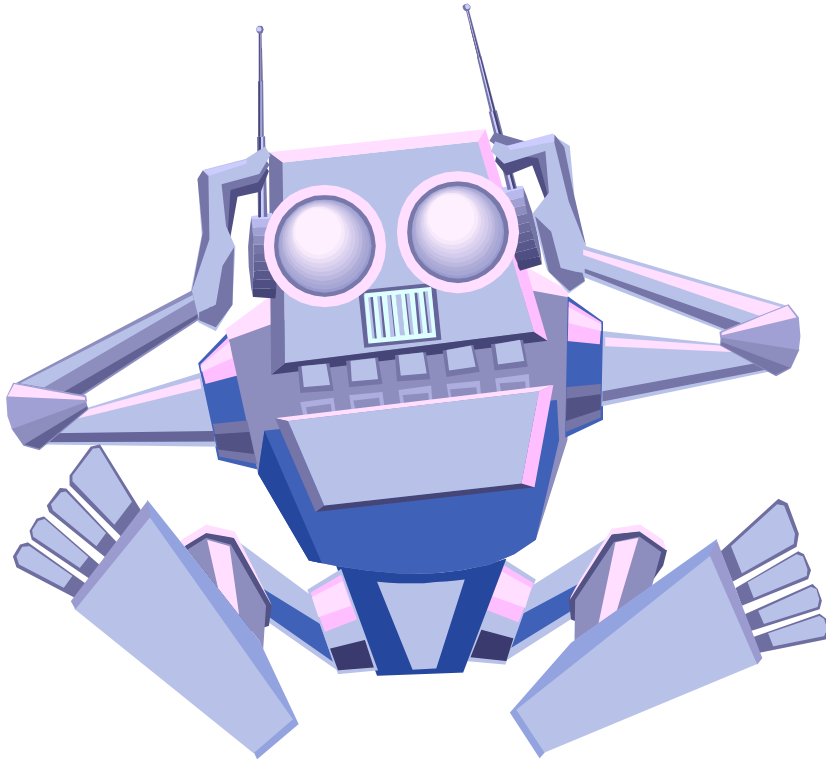


Reversible Cellular Automata



Dr. Richard Spillman

OUTLINE



What an advanced class!

- Reversible Computing
- Reversible Cellular Automata
- Billiard Ball Model of Computation

Reversible Computing

- ⇒ A reversible computing is a **backward deterministic system**
 - ⇒ Each state of the system has at most one predecessor
 - ⇒ The result is that its computation path **can always be retraced**
- ⇒ It turns out that reversible computing does not consume energy or produce heat
- ⇒ **Reversible Computing** is a general Computer Science Idea that is **more general** than Reversible Logic Circuits that we discussed in the past

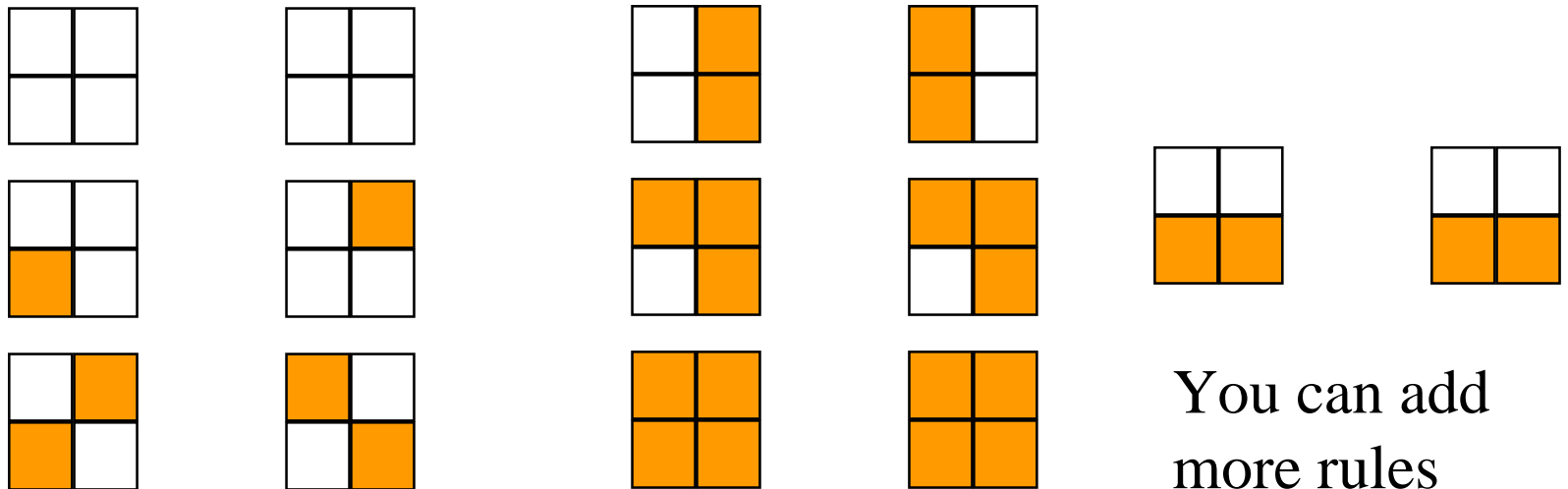
Reversible Cellular Automata

- ⇒ A reversible cellular automata (**RCA**) is a cellular automata for which each state has *at most one predecessor*
 - ⇒ Given any current state it is possible to trace it back to its initial state
- ⇒ An RCA can be implemented that does not require any cooling or energy (in theory)
 - ⇒ The rules could be implemented using Fredkin gates

Example of RCA Rules

⇒ For our example class of RCA, the **neighborhood** will consist of 2x2 blocks of cells in a 2-d cellular automata

⇒ The transformation rules are:



Blocking

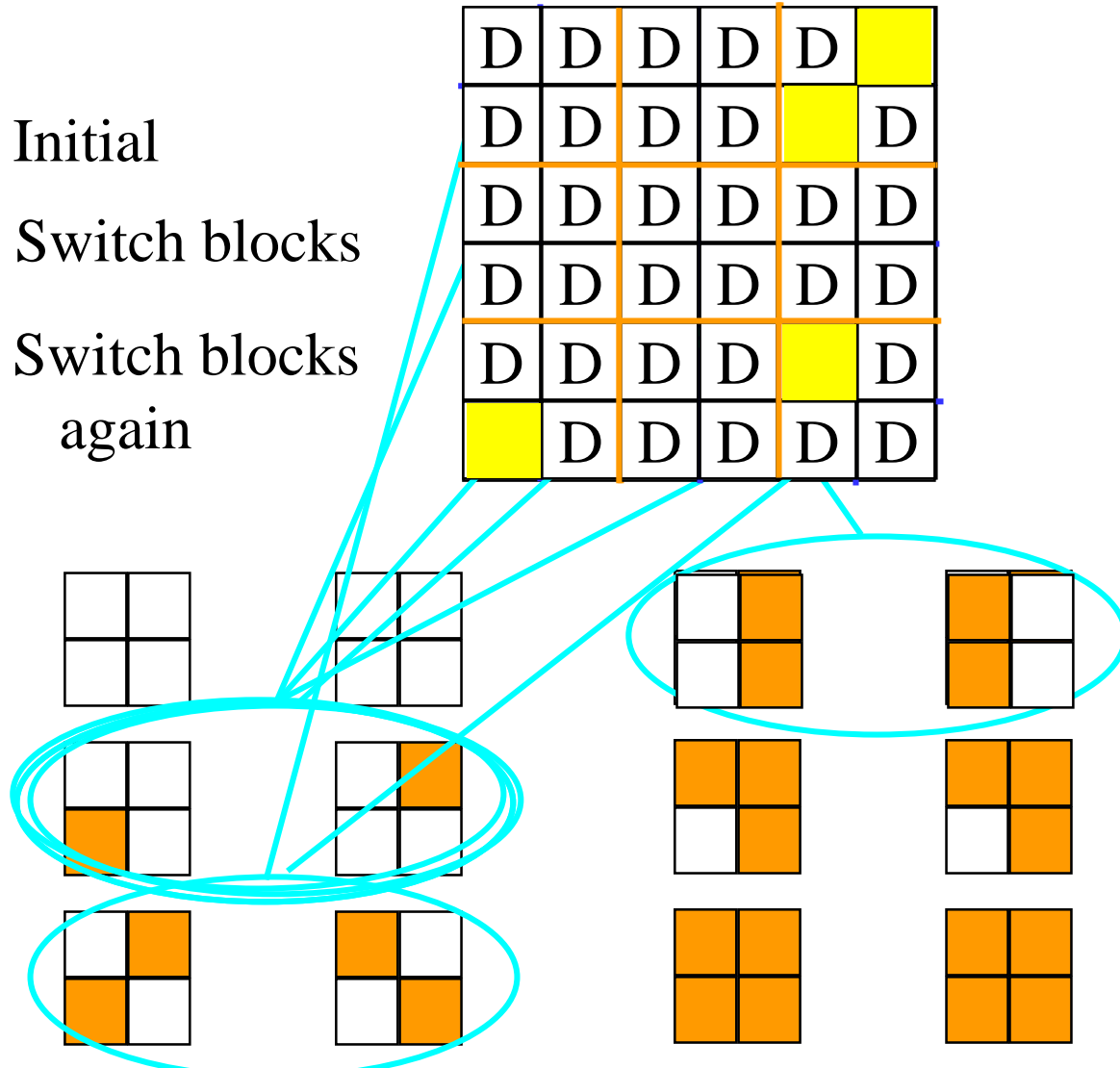
- ⇒ With this new neighborhood definition comes **another alteration** in standard cellular automata structure
 - ⇒ The 2x2 cells are **blocked** into **two different sets**
 - ⇒ The rules are applied in an **alternating cycle** to the **two different blocks**

D	D	D	D	D	D
D	D	D	D	D	D
D	D	D	D	D	D
D	D	D	D	D	D
D	D	D	D	D	D
D	D	D	D	D	D

D	D	D	D	D	D
D	D	D	D	D	D
D	D	D	D	D	D
D	D	D	D	D	D
D	D	D	D	D	D
D	D	D	D	D	D

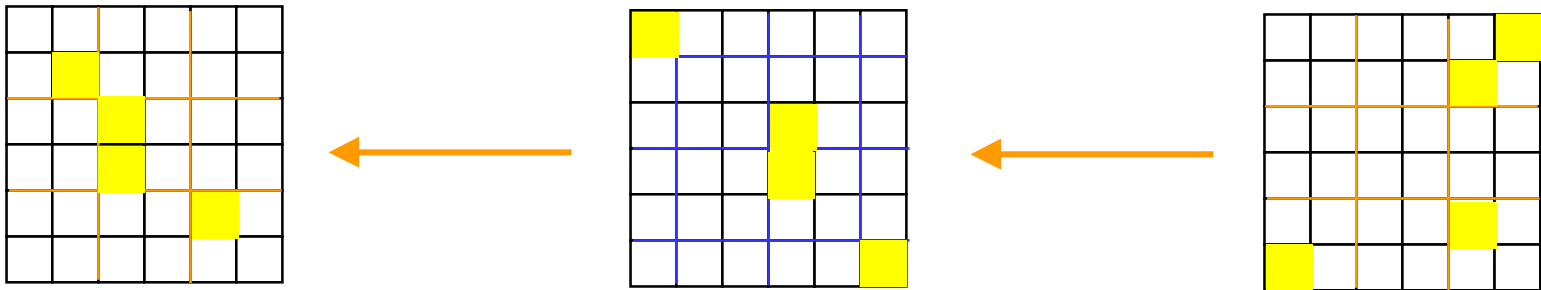
Margolus
Neighborhood

Example



Reverse

⇒ Given the **beginning and ending positions** of the prior slide show that it run backwards

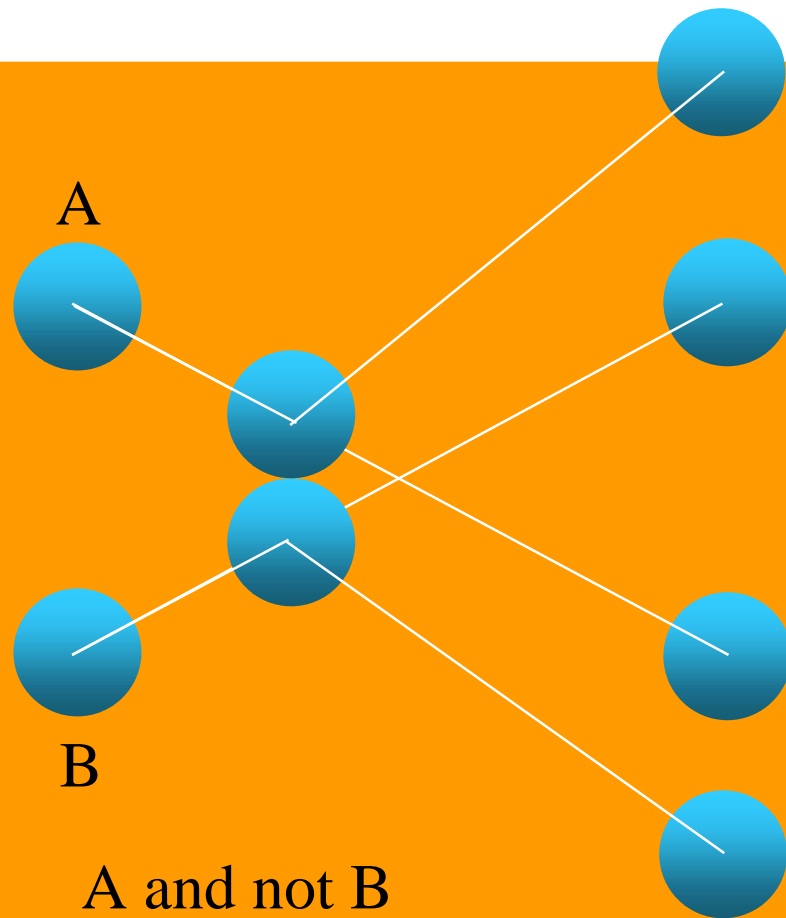


Reminder: A Billiard Ball Model of Computation

- ⇒ Create a model of computation based on the motion of particles
 - ⇒ Billiard Ball Model (BBM)
 - ⇒ Run on a reversible cellular automata
 - ⇒ The position of a billiard model is indicated by a 1 in a cell
 - ⇒ Every place where a collision of finite-diameter hard spheres might occur can be viewed as a logic gate

Example

→ We can put billiard balls at either points A, B, or both



A

B

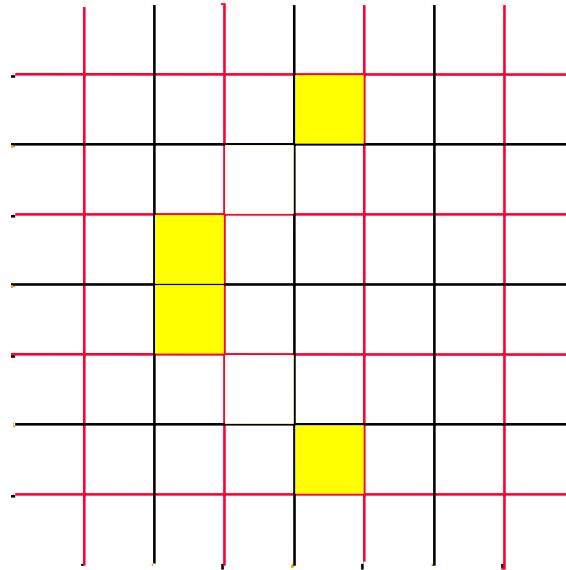
A and not B

B and not A

A and B

Cellular Automata Behavior

⇒ A collision looks like:



This type of system is used to model the classical behavior of many particle systems in physics

Possible Homeworks with CA Focus

- ⇒ Explore the nature of reversible computing
 - ⇒ Why it is necessary
 - ⇒ How is it implemented
- ⇒ Build a BBM and explore the different behavior patterns