

# CS137: Electronic Design Automation

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Systolic Algorithms



# Today

- Systolic Style
- Systolic Algorithms and Structures
  - Filters
  - Matching
  - Matrix Multiply and Transitive Closure
  - Dynamic Programming
  - Priority Queue
- Decomposition

# Systolic

- **Systole** a rhythmically recurrent contraction; *esp*: the contraction of the heart by which the blood is forced onward and the circulation kept up --- *Merriam Webster*
- **Systolic architectures** are characterized by regular, rhythmic, data flow through numerous processing elements

# Motivation

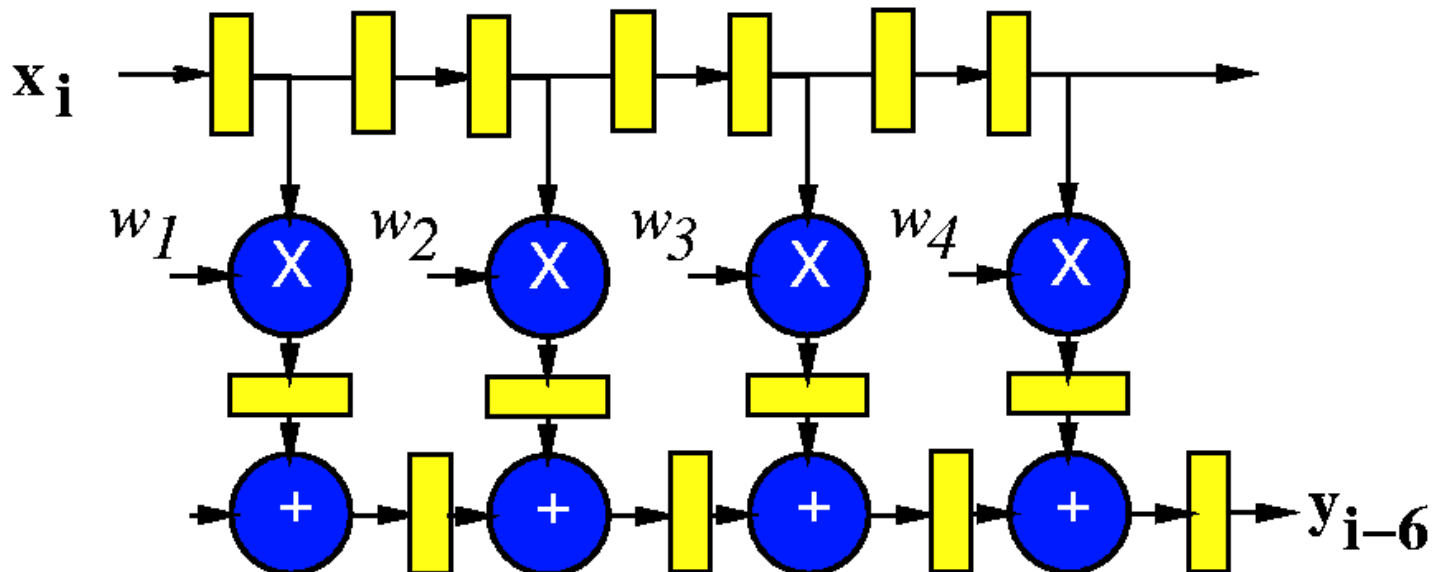
- Accelerate computationally bound problems
- For each data item read from memory
  - Use it multiple times
  - Perform many operations
  - Compute on intermediate results before discarding/writing back to memory

# Philosophy

- Match VLSI Costs
  - Exploit concurrency
  - Wires take up space
  - Long wires take time
  - Minimize design effort
    - Exploit regularity
    - Scale with problem size and/or perf. requirements

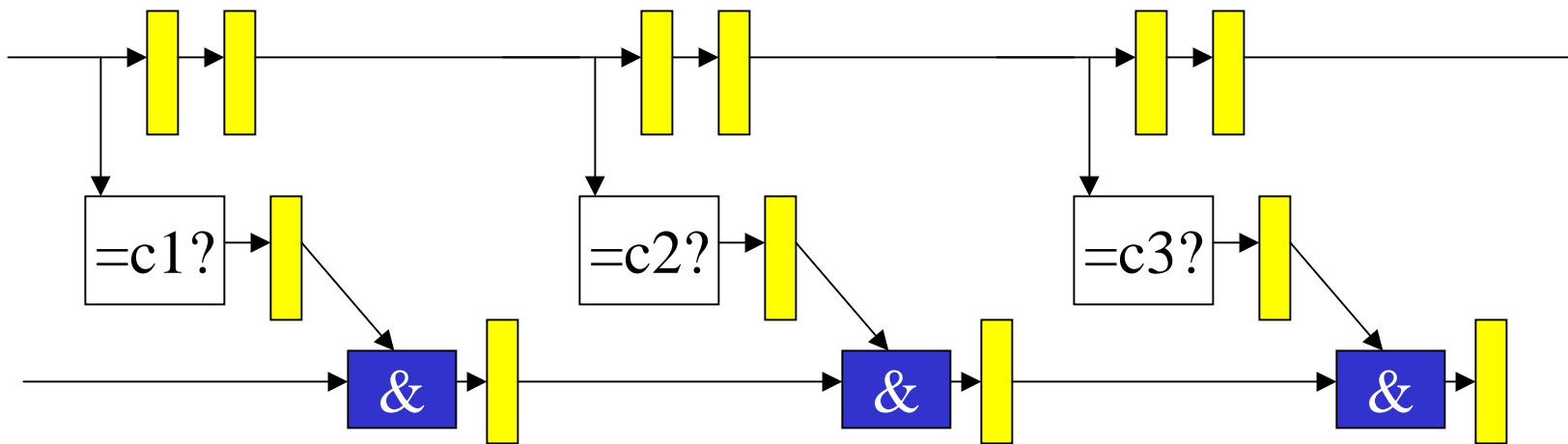
# Systolic Filter

- Finite-Impulse Response Filter
  - $y[n]=w_1 * x[n]+w_2 * x[n-1]+w_3 * x[n-3]...$
  - weighted sum of  $k$  previous samples



# Simple String Matching

- Stream text past set of matches



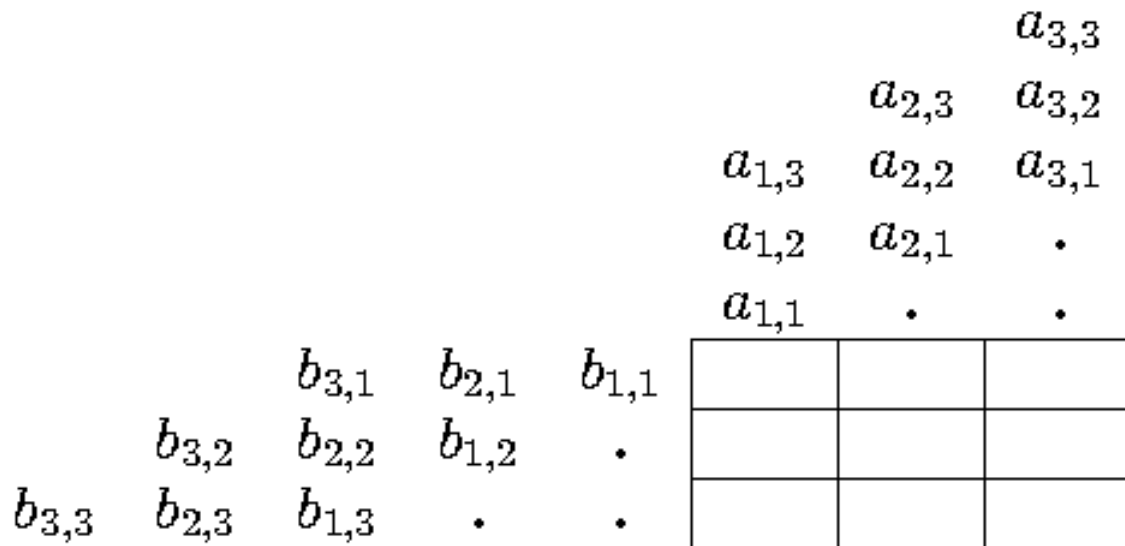
# Matrix Multiply

- $C = B * A$
- $C[i,j] = \sum_k A[i,k] * B[k,j]$
- NxN matrix
- Takes  $N^3$  multiplications



# Systolic Matrix Multiply

- Time  $N$  with  $N^2$  cells





# Transitive Closure

- Same basic shape
- $C[i,j]=a[i,j]+OR_k (a[i,k]\&a[k,j])$
- Replace
  - Sum with OR
  - Multiply with AND

# Dynamic Programming

- Two Examples
  - SPLASH sequence matching/edit distances
    - $O(N^2)$  operation in  $O(N)$  time with  $O(N)$  hardware
  - CMU parenthesis matching
    - $O(N^3)$  operation in  $O(N)$  time with  $O(N^2)$  hardware

# Sequence Matching

- Find *edit distance* between two strings
  - E.g.
    - Insert cost 1
    - Delete cost 1
    - Replace cost 2
    - Match 0

# Edit Example

- SHMOO
  - Add E (cost 1)
- SHMOOE
  - Remove M (cost  $1 + 1=2$ )
- SHOOE
  - Replace O with R (cost  $2+2=4$ )
- SHORE

# Dynamic Programming

- Build a table representing string prefixes
- Fill in costs

		S	H	O	R	E
S						
H						
M						
O						
O						

# Local Move Costs

- $D(0,0) = 0$
- $D(i,0) = D(i-1,0) + \text{Delete}(S_i)$
- $D(0,j) = D(0,j-1) + \text{Insert}(T_j)$
- $D(i,j) = \min$ 
  - $D(i-1,j) + \text{Delete}(S_i)$
  - $D(i,j-1) + \text{Insert}(T_j)$
  - $D(i-1,j-1) + \text{Replace}(S_i, T_j)$

		S	H	O	R	E
S						
H						
M						
O						
O						



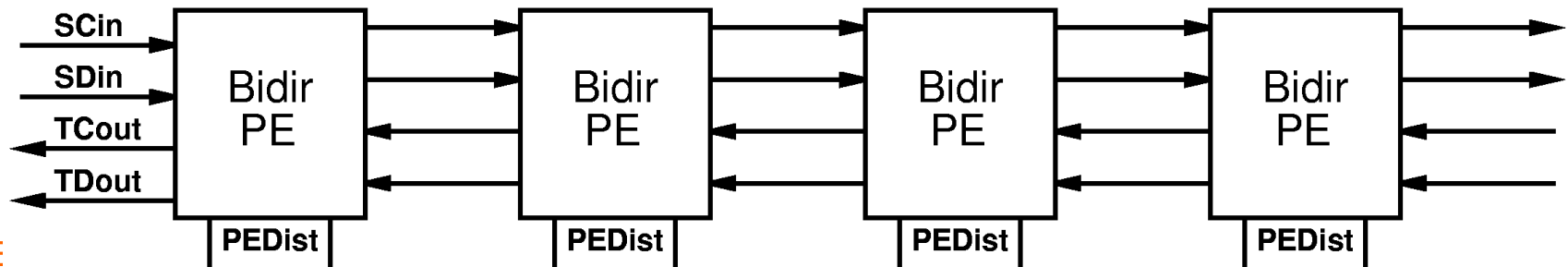
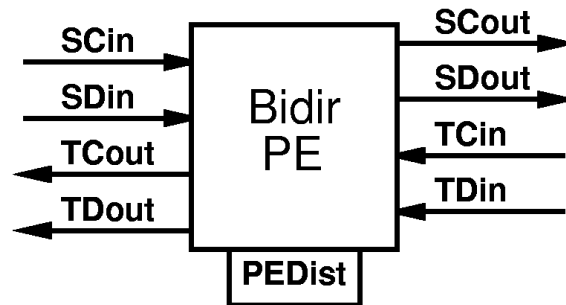
# Edit Distance Table

		S	H	O	R	E
	0	1	2	3	4	5
S	1	0	1	2	3	4
H	2	1	0	1	2	3
M	3	2	1	2	3	4
O	4	3	2	1	2	3
O	5	4	3	2	3	4

# Systolic Array

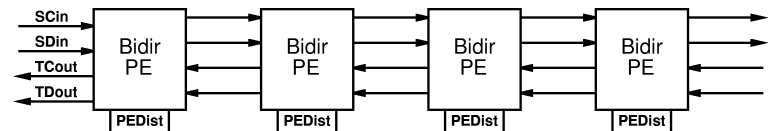
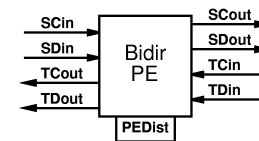
- Feed Strings from opposite ends
- Compute along diagonals

		S	H	O	R	E
S						
H						
M						
O						
O						



# Systolic Array

- $PEDist = \min$ 
  - $PEDist + Cost(Substitute(TCin, SCin))$
  - $TDin + Cost(Delete(SCin))$
  - $SDin + Cost(Insert(TCin))$
- $SDout = TDout = PEDist$
- ....plus details for edge cases



# In Operation

		<del>S</del>	<del>H</del>	O	R	E
	<del>0</del>	<del>1</del>	<del>2</del>	3	4	5
<del>S</del>	<del>1</del>	<del>0</del>	1	2	3	4
<del>H</del>	<del>2</del>	1	<del>0</del>	1	2	3
M	3	2	1	2	3	4
O	4	3	2	1	2	3
O	5	4	3	2	3	4

TC	O		O		M		H		S										
TD	5		4		3		2		1										
TDout	5		4		3		2		1										
SC											S		H		O		R		E
SD											1		2		3		4		5
SDout											1		2		3		4		5
PEDist	5	4	4	3	3	2	2	1	1	0	1	1	2	2	3	3	4	4	5
TC		O		O		M		H		S									
TD		5		4		3		2		1									
TDout		5		4		3		2		0									
SC										S		H		O		R		E	
SD										1		2		3		4		5	
SDout										0		2		3		4		5	
PEDist	5	5	4	4	3	3	2	2	1	0	1	2	2	3	3	4	4	5	5
TC			O		O		M		H		S								
TD			5		4		3		2		0								
TDout			5		4		3		1		1								
SC									S		H		O		R		E		
SD									0		2		3		4		5		
SDout									1		1		3		4		5		
PEDist	5	5	5	4	4	3	3	2	1	0	1	2	3	3	4	4	5	5	5

# In Operation (Cont.)

		S	H	<del>O</del>	R	E
	0	1	2	3	4	5
S	1	0	1	2	3	4
H	2	1	0	1	2	3
M	3	2	1	2	3	4
O	4	3	2	1	2	3
O	5	4	3	2	3	4

TC				O		O		M		H		S							
TD				5		4		3		1		1							
TDout				5		4		2		0		2							
SC								S		H		O		R		E			
SD								1		1		3		4		5			
SDout								2		0		2		4		5			
PEDist	5	5	5	5	4	4	3	2	1	0	1	2	3	4	4	5	5	5	5
TC					O		O		M		H		S						
TD					5		4		2		0		2						
TDout					5		3		1		1		3						
SC							S		H		O		R		E				
SD							2		0		2		4		5				
SDout							3		1		1		3		5				
PEDist	5	5	5	5	5	4	3	2	1	0	1	2	3	4	5	5	5	5	5
TC						O		O		M		H		S					
TD						5		3		1		1		3					
TDout						4		2		2		2		4					
SC						S		H		O		R		E					
SD						3		1		1		3		5					
SDout						4		2		2		2		4					
PEDist	5	5	5	5	5	4	3	2	1	2	1	2	3	4	5	5	5	5	5

		S	H	O	R	E
	0	1	2	3	4	5
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H	2	1	0	1	2	3
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O	4	3	2	1	2	3
O	5	4	3	2	3	4

# In Operation (Cont.)

TC							O		O		M		H		S				
TD							4		2		2		2		4				
TDout							3		2		3		3		4				
SC				S			H		O		R		E						
SD				4			2		2		2		4						
SDout				4			3		2		3		3						
PEDist	5	5	5	5	4	4	3	2	2	2	3	2	3	4	4	5	5	5	5
TC									O		O		M		H		S		
TD									3		2		3		3		4		
TDout									2		3		4		3		4		
SC				S			H		O		R		E						
SD				4			3		2		3		3						
SDout				4			3		2		3		4						
PEDist	5	5	5	4	4	3	3	2	2	3	3	4	3	3	4	4	5	5	5
TC										O		O		M		H		S	
TD										2		3		4		3		4	
TDout										3		4		4		3		4	
SC			S		H		O		R		E								
SD			4		3		2		3		4		4						
SDout			4		3		2		3		4		4						
PEDist	5	5	4	4	3	3	2	2	3	3	4	4	4	3	3	4	4	5	5

# In Operation (Cont.)

		S	H	O	R	E
	0	1	2	3	4	5
S	1	0	1	2	3	4
H	2	1	0	1	2	3
M	3	2	1	2	3	4
O	4	3	2	1	2	3
O	5	4	3	2	3	4

TC										O		O		M		H		S	
TD										3		4		4		3		4	
TDout										4		4		4		3		4	
SC		S		H		O		R		E									
SD		4		3		2		3		4									
SDout		4		3		2		3		4									
PEDist	5	4	4	3	3	2	2	3	3	4	4	4	4	4	3	3	4	4	5

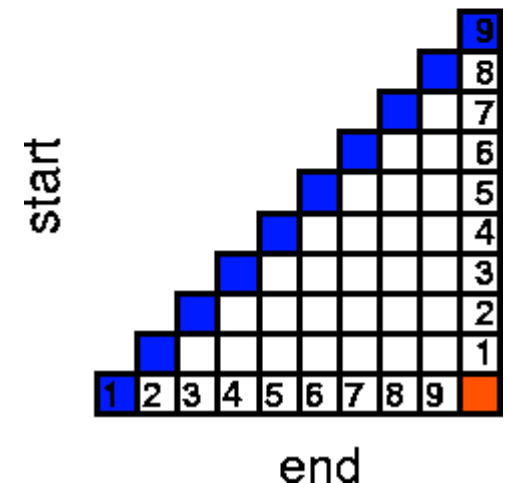
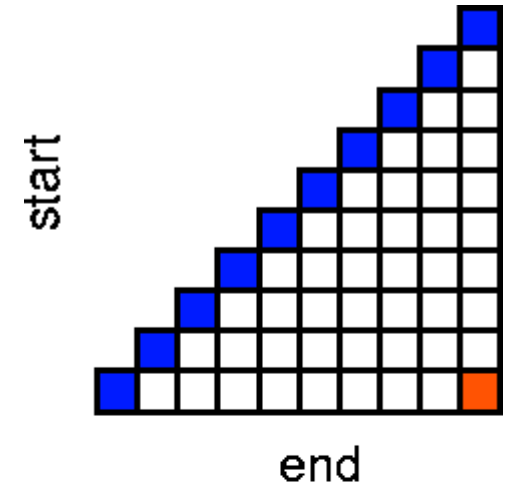
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O	4	3	2	1	2	3
O	5	4	3	2	3	4



# Parenthesis Matching

- Similar
- But compute from all breaks across a diagonal
  - Not just nearest neighbor
- Hence extra  $O(N)$

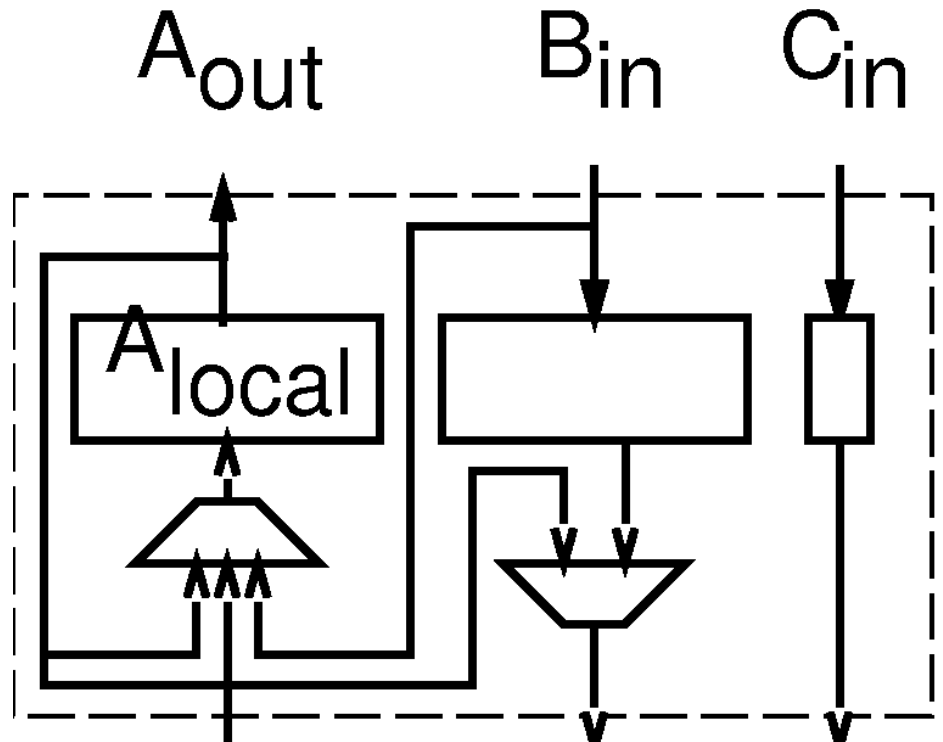


# Priority Queue

- Insert top
- Extract Largest
- W/  $O(N)$  cells
- $O(1)$  Extract

# Priority Queue Cell

- If ( $C_{in} = \text{insert}$ )  
     $A_{local} \leftarrow \text{largest}$   
     $B_{out} \leftarrow \text{smallest}$
- If ( $C_{in} = \text{extract}$ )  
     $A_{local} \leftarrow A_{in}$   
     $B_{out} \leftarrow B_{in}$
- $C_{out} \leftarrow C_{in}$



# Decomposition