

# Quantum Logic

**Marek Perkowski**

# Sources

Origin of slides: John Hayes, Peter Shor, Martin Lukac, Mikhail Pivtoraiko, Alan Mishchenko, Pawel Kerntopf, Mosca, Ekert

**Lee Spector**

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# Introduction

- Short-Term Objectives

Introduce Quantum Computing Basics to interested students at KAIST.  
Especially non-physics students

- Long-Term Objectives

Engage into AI/CS/Math Research projects benefiting from Quantum Computing. Continue our previous projects in quantum computing

- Prerequisite

- No linear algebra or quantum mechanics assumed
- A ECE, math, physics or CS background would be beneficial, practically-oriented class.

# Introduction

- MainTextbook

## Quantum Computation & Quantum Information

**Michael A. Nielsen**  
**Isaac L. Chuang**

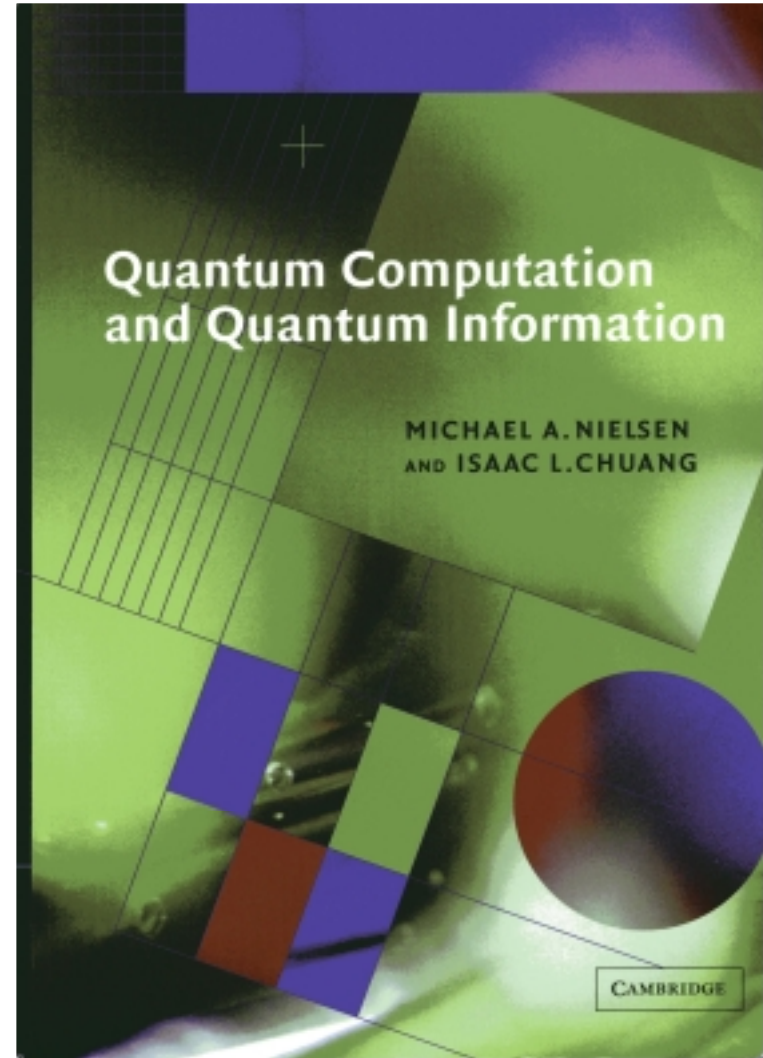
ISBN: 0 521 63503 9

Paperback

ISBN: 0 521 63235 8 Hardback

Cost: \$48.00 New Paperback  
\$35.45 Used Paperback

(<http://www.amazon.com>)  
also in KAIST bookstore



# Presentation Overview

## Qubits

1 Qubit -> Bloch Sphere,  
2 Qubits -> Bell States,  
n Qubits

## Quantum Computation

Gates: Single Qubit, Arbitrary Single Qubit -> Universal Quantum Gates, Multiple Qubit Gates -> CNOT  
Other Computational Bases

## Quantum Circuits

Qubit Swap Circuit  
Qubit Copying Circuit  
Bell State Circuit -> Quantum Teleportation

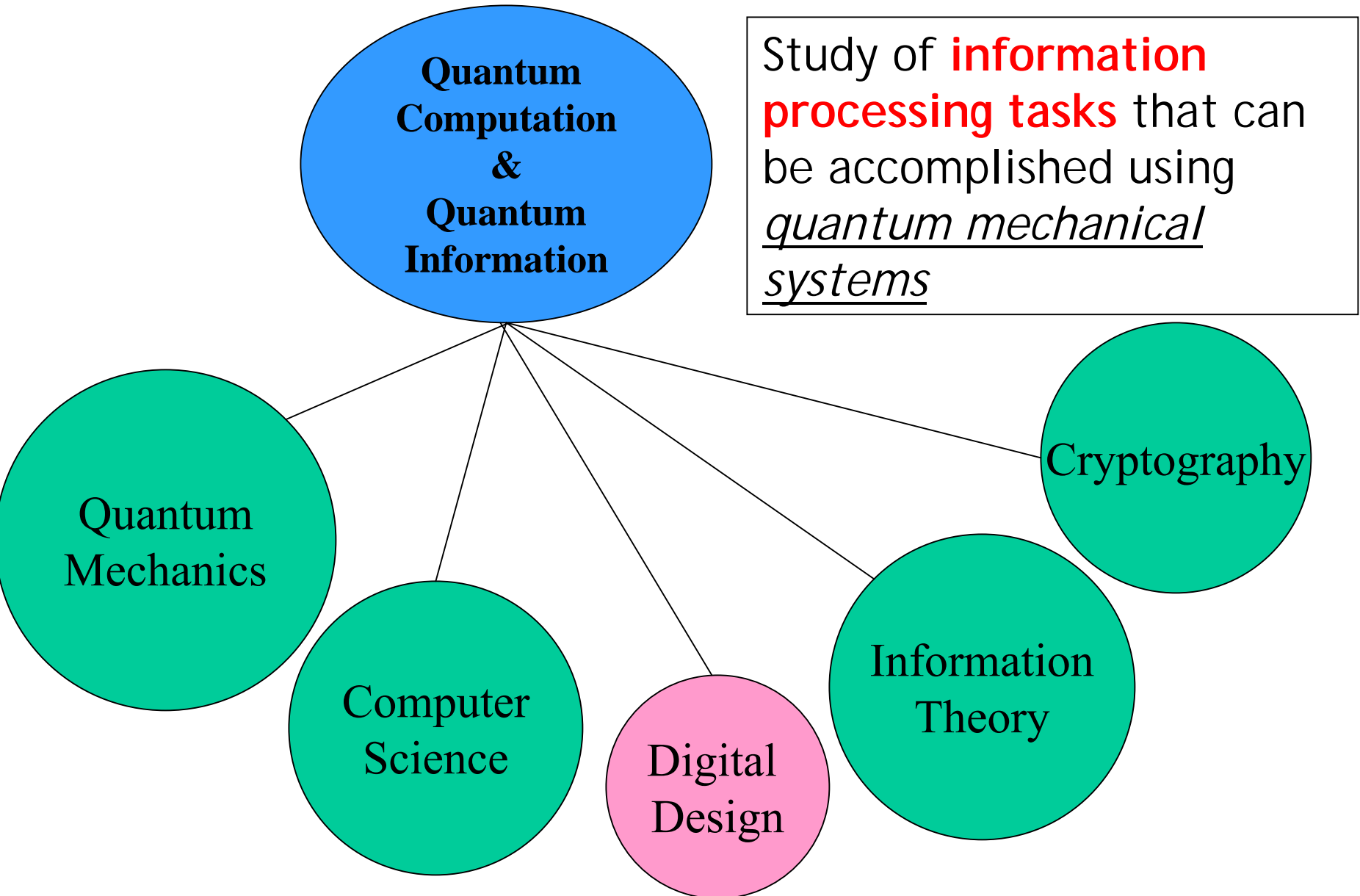
## Quantum Algorithms

Toffoli Gate -> Quantum Parallelism -> Hadamard Transform  
Deutsch's Algorithm, Deutsch-Josa Algorithm  
Other Algorithms  
- Fourier Transform, Quantum Search, Quantum Simulation

## Quantum Information Processing

Stern-Gerlach, Optical Techniques, Traps, NMR, Quantum Dots

# Historical Background and Links



# What will be discussed?

- **Background**
- **Quantum circuits synthesis and algorithms**
- **Quantum circuits simulation**
- **Quantum Computation**
- **AI for quantum computation**
  
- **Quantum computation for AI**
- **Quantum logic emulation and evolvable hardware**
- **Quantum circuits verification**
- **Quantum-based robot control**

# What is quantum computation?

- Computation with **coherent atomic-scale dynamics**.
- The behavior of a quantum computer is governed by the laws of quantum mechanics.



# Why bother with quantum computation?

- **Moore's Law:** We hit the quantum level 2010~2020.
- Quantum computation is more powerful than classical computation.
- More can be computed in less time—the complexity classes are different!

# The power of quantum computation

- In quantum systems possibilities count, **even if they never happen!**
- Each of **exponentially many possibilities** can be used to **perform a part of a computation** at the same time.

# **Nobody understands quantum mechanics**

**“No, you’re not going to be able to understand it. . . . You see, my physics students don’t understand it either. That is because I don’t understand it. Nobody does. ... The theory of quantum electrodynamics describes Nature as absurd from the point of view of common sense. And it agrees fully with an experiment. So I hope that you can accept Nature as She is -- absurd.**

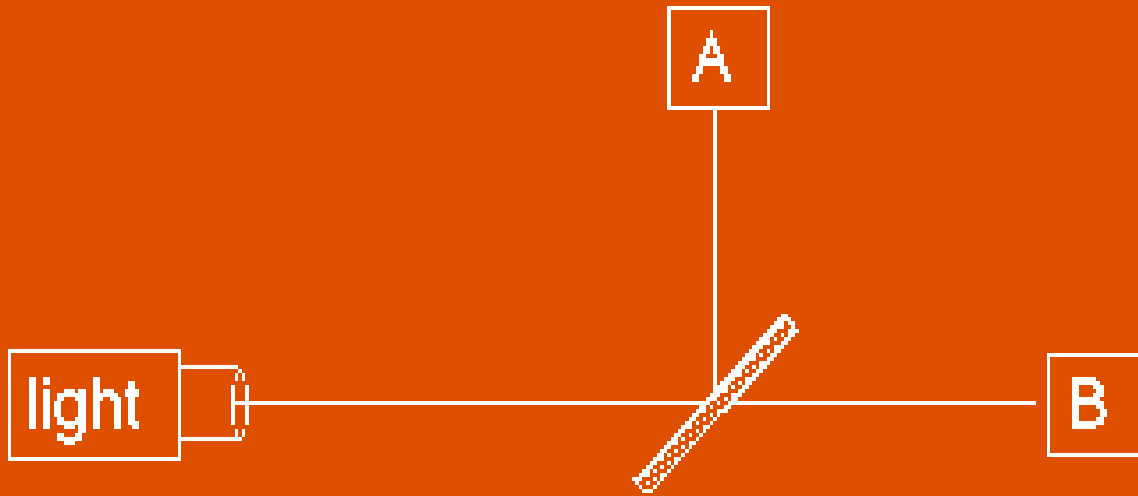
**Richard Feynman**

# **Absurd but taken seriously (not just quantum mechanics but also quantum computation)**

- **Under active investigation by many of the top physics labs around the world (including CalTech, MIT, AT&T, Stanford, Los Alamos, UCLA, Oxford, l'Université de Montréal, University of Innsbruck, IBM Research . . .)**
- **In the mass media (including The New York Times, The Economist, American Scientist, Scientific American, . . .)**
- **Here.**

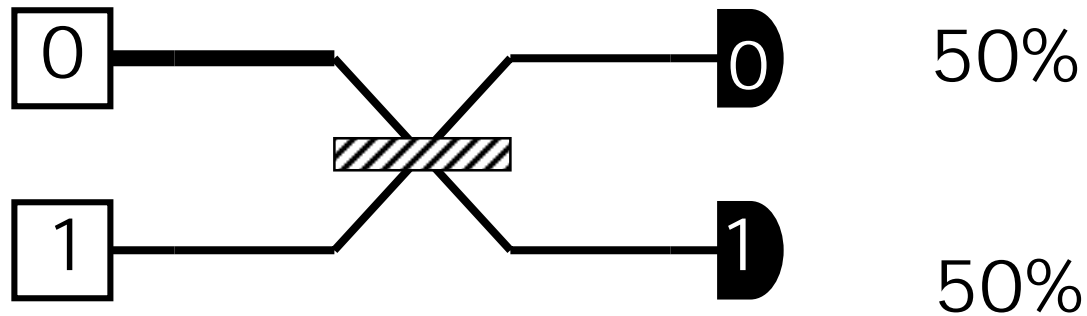
# Quantum Logic Circuits

# A beam splitter



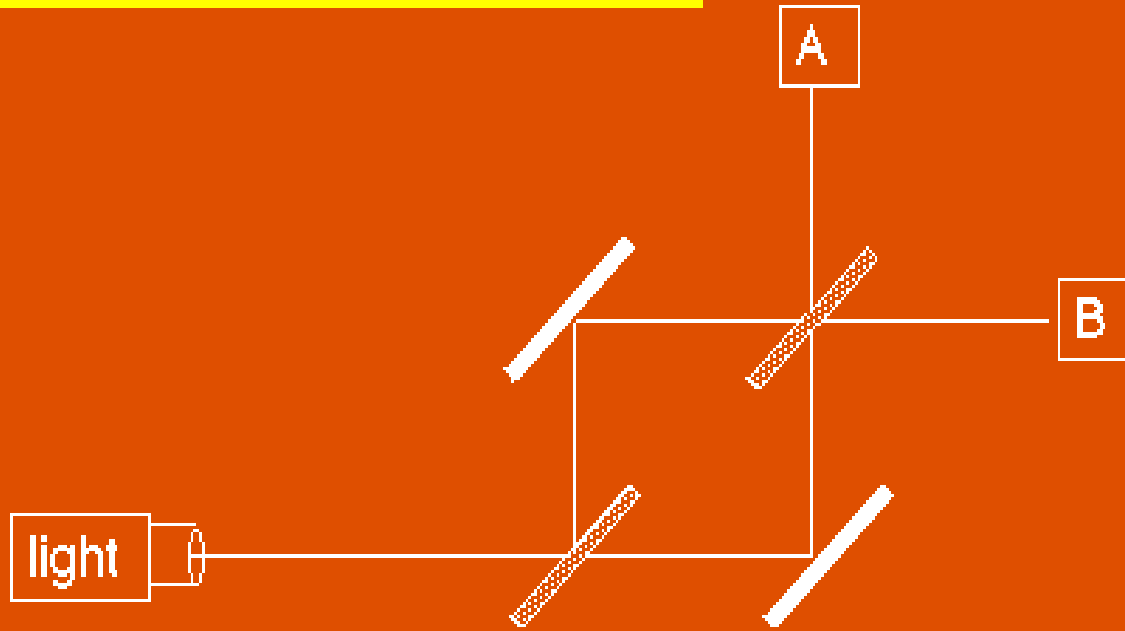
**Half of the photons leaving the light source arrive at detector A;  
the other half arrive at detector B.**

# A beam-splitter



The simplest explanation is that the beam-splitter acts as a classical coin-flip, randomly sending each photon one way or the other.

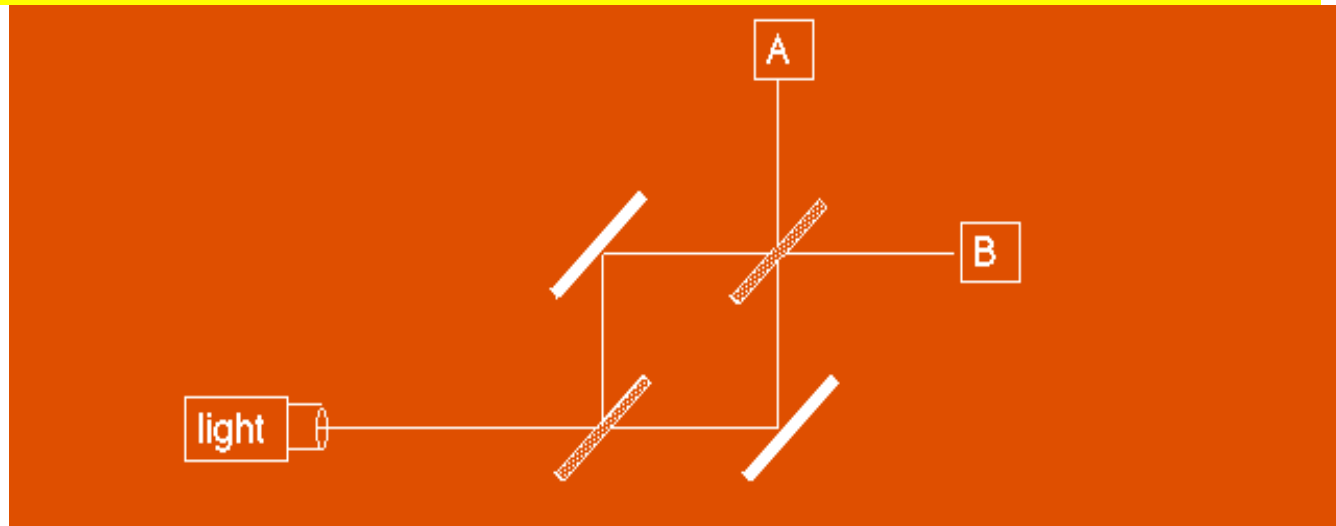
# An interferometer



- **Equal path lengths, rigid mirrors.**
- **Only one photon in the apparatus at a time.**
- **All photons leaving the source arrive at B.**
- **WHY?**



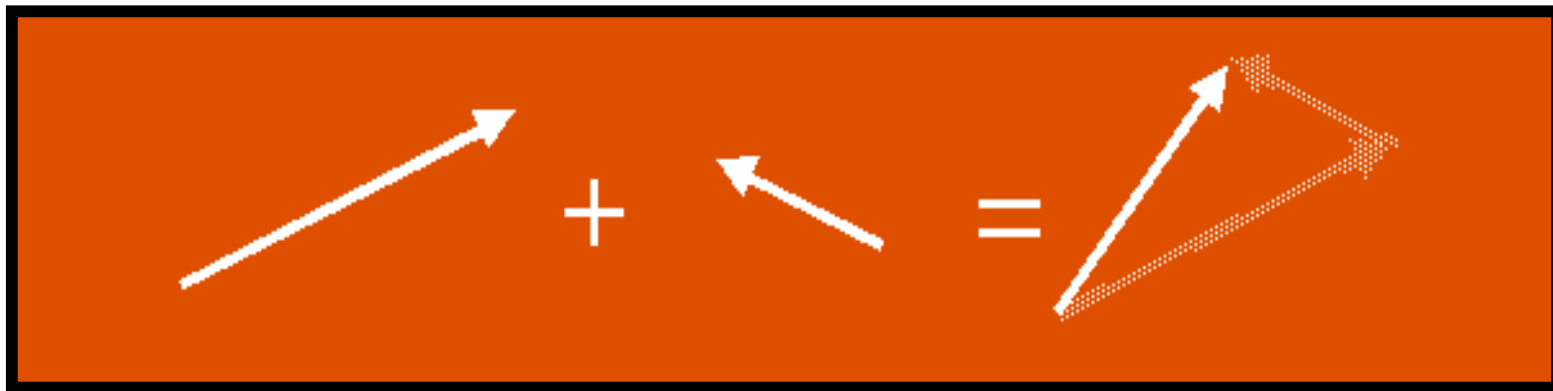
# Possibilities count



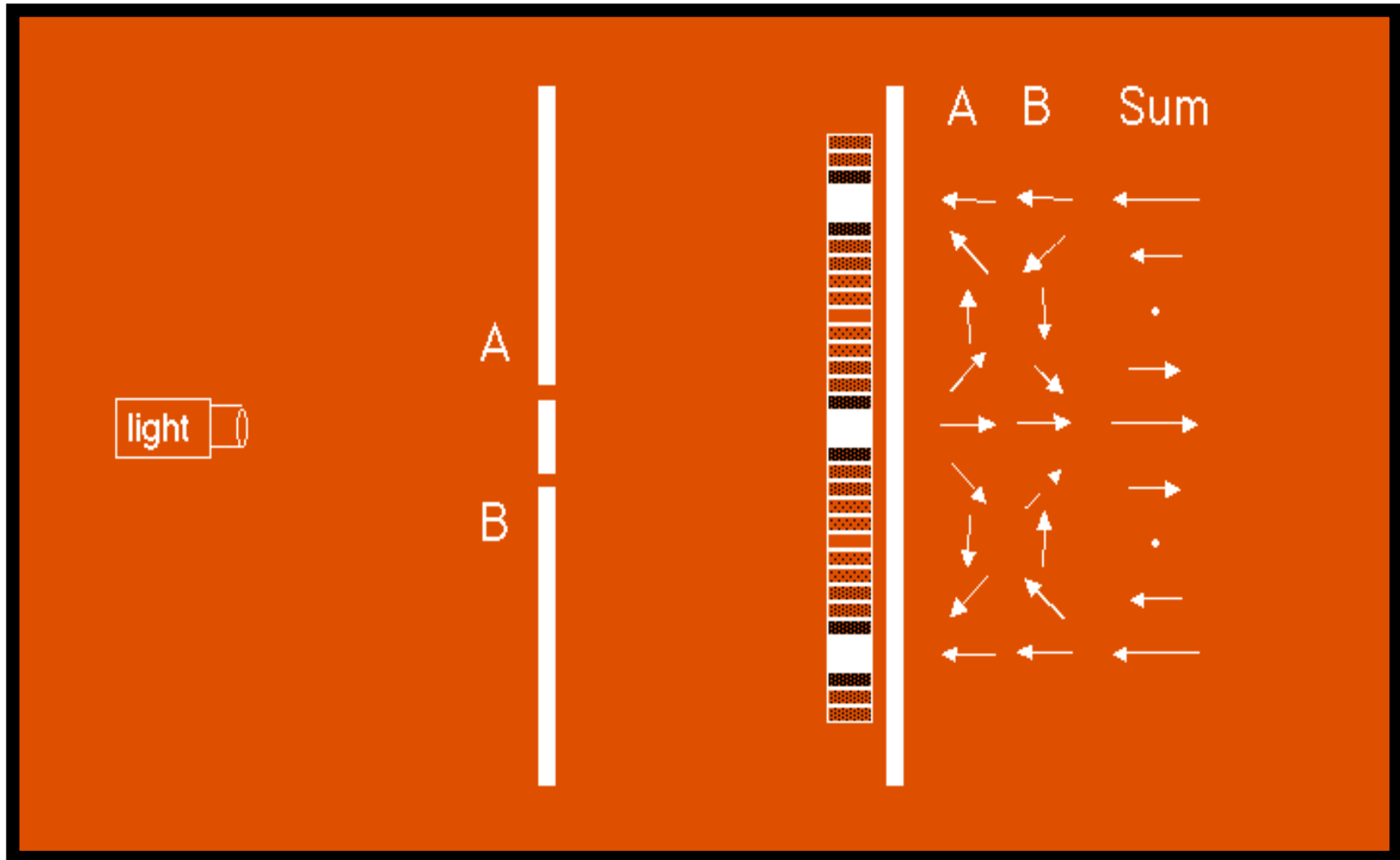
- There is a quantity that we'll call the **“amplitude”** for each possible path that a photon can take.
- The **amplitudes can interfere constructively and destructively**, even though each photon takes only one path.
- The amplitudes at detector A interfere destructively; those at detector B interfere constructively.

# Calculating interference

- Arrows for each possibility.
- Arrows rotate; speed depends on frequency.
- Arrows flip  $180^\circ$  at mirrors, rotate  $90^\circ$  counter-clockwise when reflected from beam splitters.
- Add arrows and square the length of the result to determine the probability for any possibility.



# Double slit interference



# Quantum Interference : Amplitudes are added and not intensities !

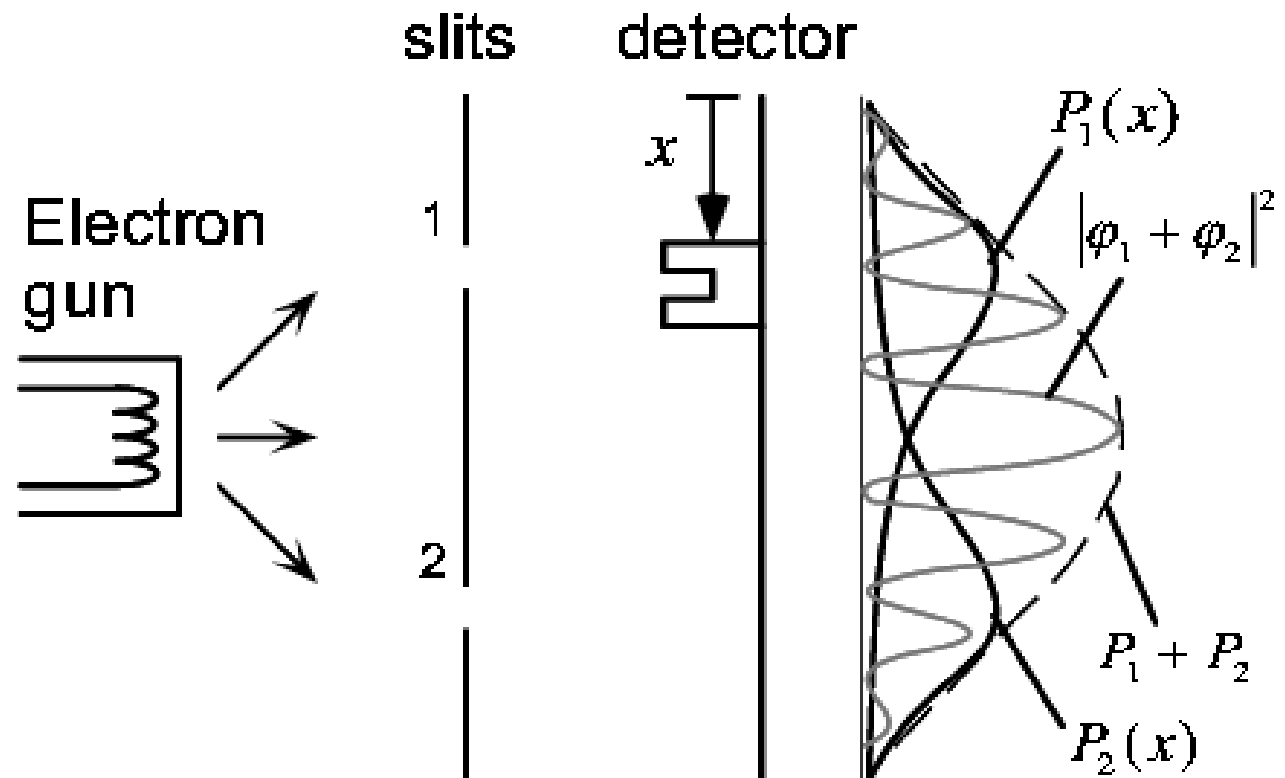
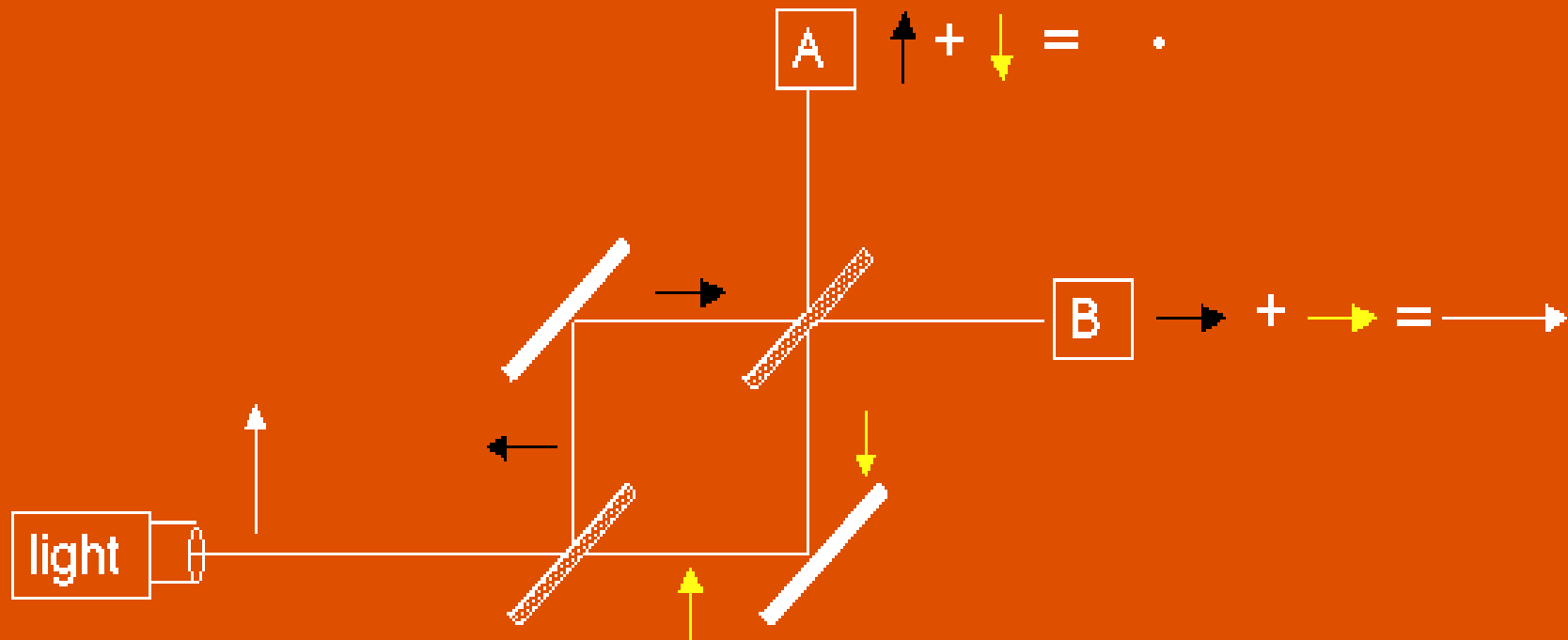
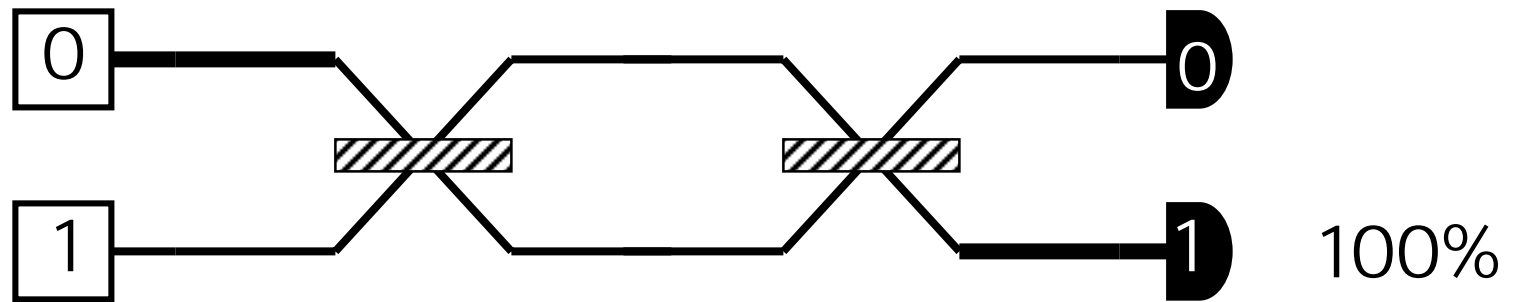


Figure 1: Two-slit experiment.

# Interference in the interferometer

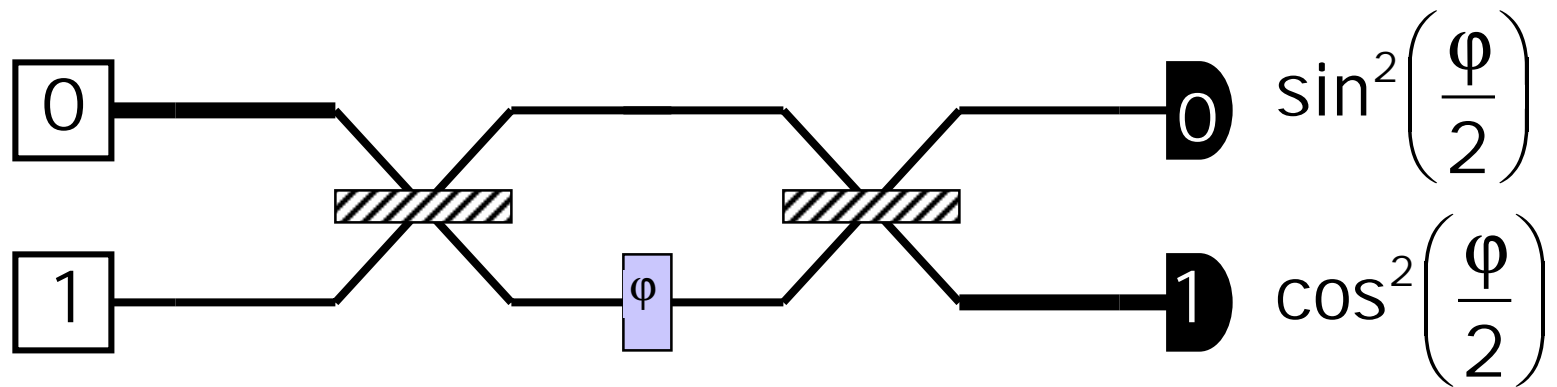


# Quantum Interference



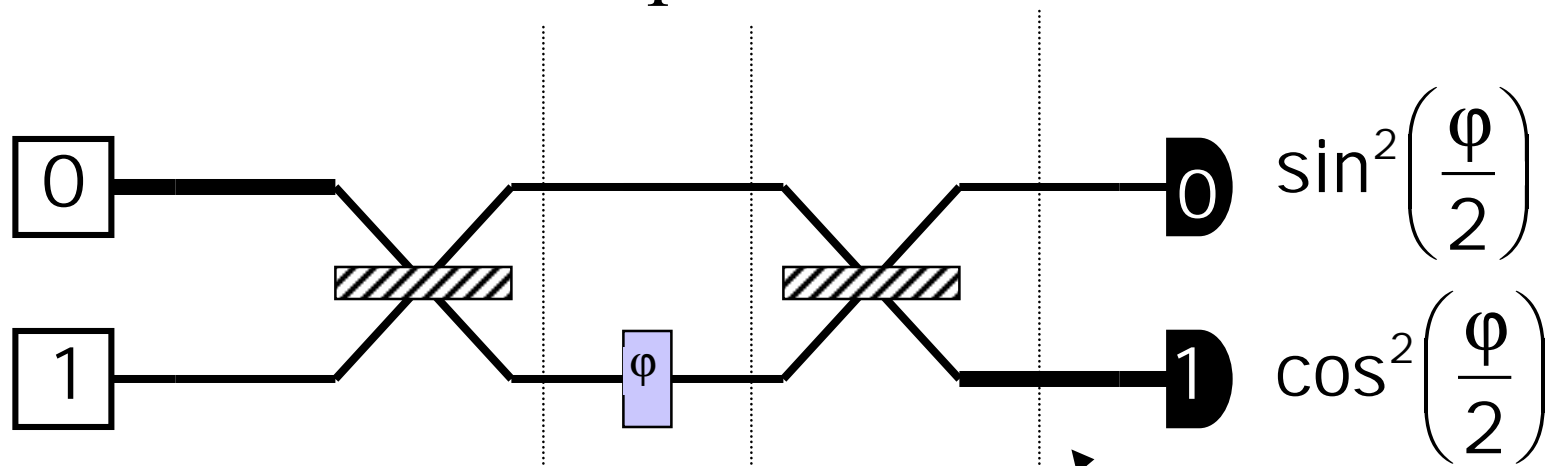
The simplest explanation must be wrong, since it would predict a 50-50 distribution.

# More experimental data



# A new theory

The particle can exist in a linear combination or *superposition* of the two paths



$$\frac{i}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$$

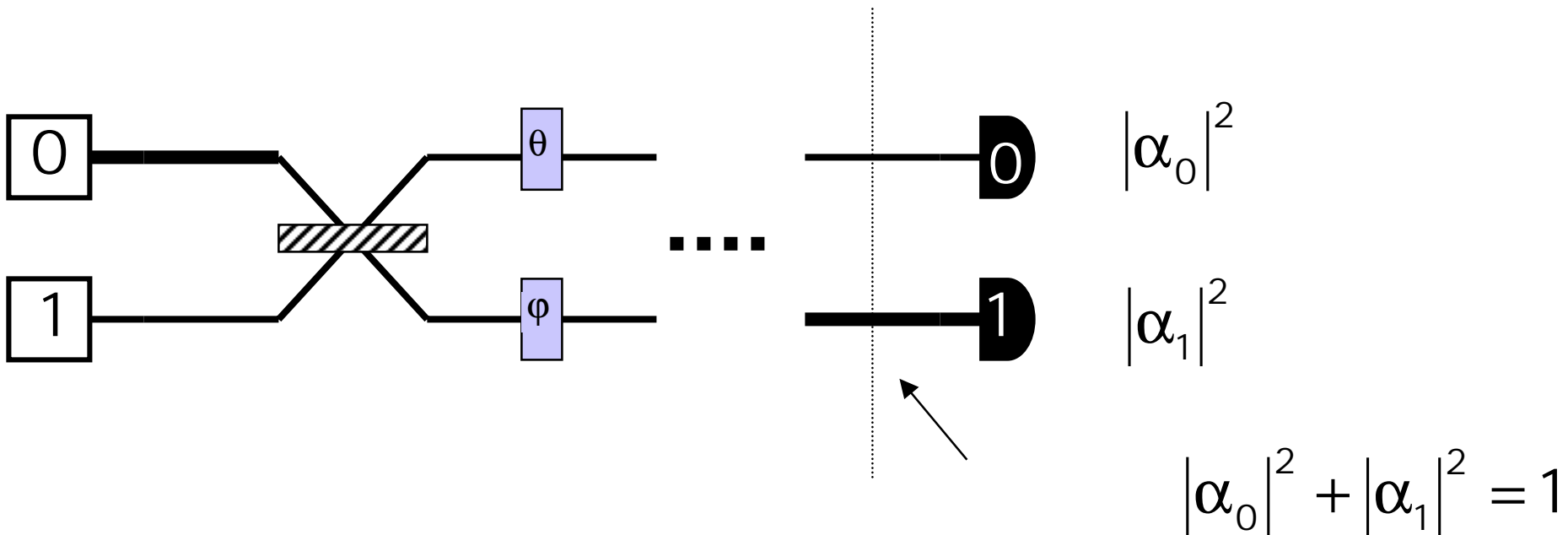
$$\frac{i}{\sqrt{2}}|0\rangle + \frac{e^{i\phi}}{\sqrt{2}}|1\rangle$$

$$\frac{e^{i\phi} - 1}{2}|0\rangle + \frac{i(e^{i\phi} + 1)}{\sqrt{2}}|1\rangle$$



# Probability Amplitude and Measurement

If the photon is measured when it is in the state  $\alpha_0|0\rangle + \alpha_1|1\rangle$  then we get  $|0\rangle$  with probability  $|\alpha_0|^2$  and  $|1\rangle$  with probability of  $|\alpha_1|^2$



# Quantum Operations

The operations are induced by the apparatus *linearly*, that is, **if**

$$|0\rangle \rightarrow \frac{i}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$$

**and**

$$|1\rangle \rightarrow \frac{1}{\sqrt{2}}|0\rangle + \frac{i}{\sqrt{2}}|1\rangle$$

**then**

$$\begin{aligned}\alpha_0|0\rangle + \alpha_1|1\rangle &\rightarrow \alpha_0\left(\frac{i}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle\right) + \alpha_1\left(\frac{1}{\sqrt{2}}|0\rangle + \frac{i}{\sqrt{2}}|1\rangle\right) \\ &= \left(\alpha_0 \frac{i}{\sqrt{2}} + \alpha_1 \frac{1}{\sqrt{2}}\right)|0\rangle + \left(\alpha_0 \frac{1}{\sqrt{2}} + \alpha_1 \frac{i}{\sqrt{2}}\right)|1\rangle\end{aligned}$$

# Quantum Operations

Any linear operation that takes states

$$\alpha_0|0\rangle + \alpha_1|1\rangle \quad \text{satisfying} \quad |\alpha_0|^2 + |\alpha_1|^2 = 1$$

and maps them to states

$$\alpha'_0|0\rangle + \alpha'_1|1\rangle \quad \text{satisfying} \quad |\alpha'_0|^2 + |\alpha'_1|^2 = 1$$

must be **UNITARY**

# Linear Algebra

$$U = \begin{bmatrix} u_{00} & u_{01} \\ u_{10} & u_{11} \end{bmatrix}$$

is unitary **if and only if**

$$UU^t = \begin{bmatrix} u_{00} & u_{01} \\ u_{10} & u_{11} \end{bmatrix} \begin{bmatrix} u_{00}^* & u_{10}^* \\ u_{01}^* & u_{11}^* \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = I$$

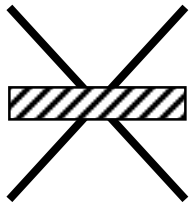
# Linear Algebra

$|0\rangle$  corresponds to  $\begin{pmatrix} 1 \\ 0 \end{pmatrix}$

$|1\rangle$  corresponds to  $\begin{pmatrix} 0 \\ 1 \end{pmatrix}$

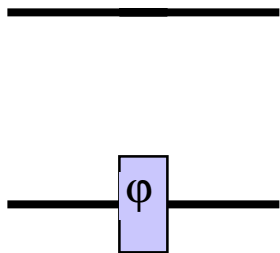
$\alpha_0|0\rangle + \alpha_1|1\rangle$  corresponds to  $\alpha_0\begin{pmatrix} 1 \\ 0 \end{pmatrix} + \alpha_1\begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} \alpha_0 \\ \alpha_1 \end{pmatrix}$

# Linear Algebra



corresponds to

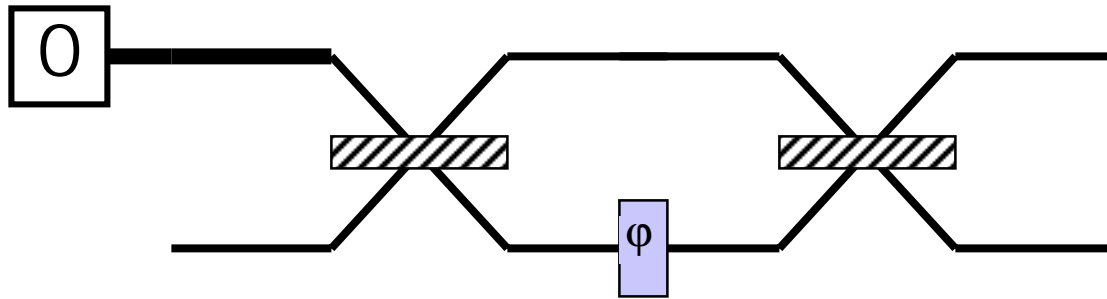
$$\begin{pmatrix} \frac{i}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{i}{\sqrt{2}} \end{pmatrix}$$



corresponds to

$$\begin{pmatrix} 1 & 0 \\ 0 & e^{i\varphi} \end{pmatrix}$$

# Linear Algebra



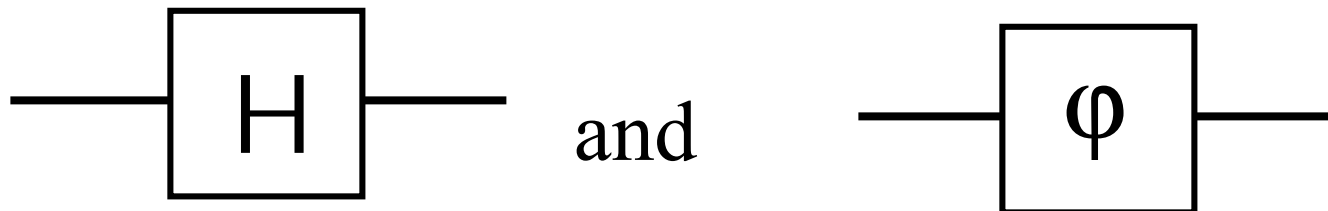
corresponds to

$$\begin{pmatrix} \frac{i}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{i}{\sqrt{2}} \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & e^{i\phi} \end{pmatrix} \begin{pmatrix} \frac{i}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{i}{\sqrt{2}} \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

# Abstraction

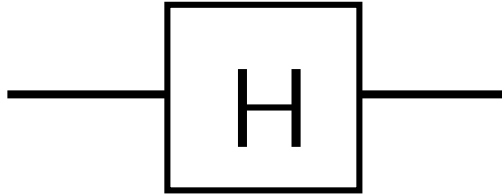
The **two position states** of a photon in a **Mach-Zehnder apparatus** is just one example of a quantum bit or ***qubit***

Except when addressing a particular physical implementation, we will simply talk about “**basis**” states  $|0\rangle$  and  $|1\rangle$  and **unitary operations** like





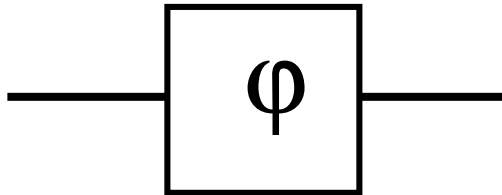
where



corresponds to

$$\begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{pmatrix}$$

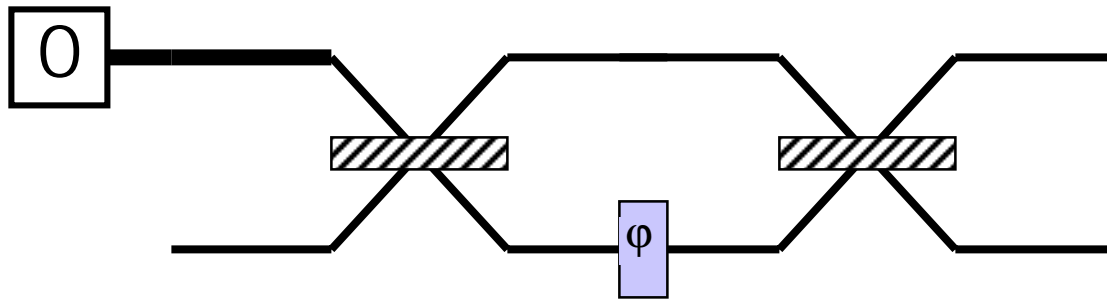
and



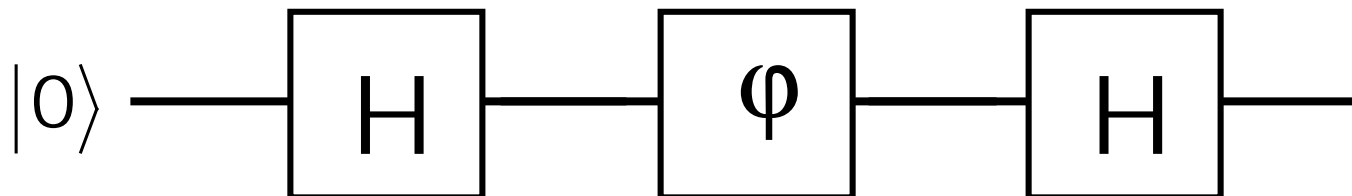
corresponds to

$$\begin{pmatrix} 1 & 0 \\ 0 & e^{i\phi} \end{pmatrix}$$

An arrangement like



is represented with a network like



# More than one qubit

If we concatenate two qubits

$$(\alpha_0|0\rangle + \alpha_1|1\rangle) (\beta_0|0\rangle + \beta_1|1\rangle)$$

we have a 2-qubit system with **4 basis states**

$$|0\rangle|0\rangle = |00\rangle \quad |0\rangle|1\rangle = |01\rangle \quad |1\rangle|0\rangle = |10\rangle \quad |1\rangle|1\rangle = |11\rangle$$

and we can also describe the state as

$$\alpha_0\beta_0|00\rangle + \alpha_0\beta_1|01\rangle + \alpha_1\beta_0|10\rangle + \alpha_1\beta_1|11\rangle$$

or by the vector

$$\begin{pmatrix} \alpha_0\beta_0 \\ \alpha_0\beta_1 \\ \alpha_1\beta_0 \\ \alpha_1\beta_1 \end{pmatrix} = \begin{pmatrix} \alpha_0 \\ \alpha_1 \end{pmatrix} \otimes \begin{pmatrix} \beta_0 \\ \beta_1 \end{pmatrix}$$

# More than one qubit

In general we can have arbitrary superpositions

$$\alpha_{00}|0\rangle|0\rangle + \alpha_{01}|0\rangle|1\rangle + \alpha_{10}|1\rangle|0\rangle + \alpha_{11}|1\rangle|1\rangle$$

$$|\alpha_{00}|^2 + |\alpha_{01}|^2 + |\alpha_{10}|^2 + |\alpha_{11}|^2 = 1$$

where there is no factorization into the tensor product of two independent qubits.

These states are called *entangled*.

# Entanglement

- Qubits in a multi-qubit system are not independent—they can become “entangled.”
- To represent the state of  $n$  qubits we use  $2^n$  complex number amplitudes.

# Measuring multi-qubit systems

If we measure both bits of

$$\alpha_{00}|0\rangle|0\rangle + \alpha_{01}|0\rangle|1\rangle + \alpha_{10}|1\rangle|0\rangle + \alpha_{11}|1\rangle|1\rangle$$

we get  $|x\rangle|y\rangle$  with probability  $|\alpha_{xy}|^2$

# Measurement

- $\sum |\alpha|^2$ , for amplitudes of all states matching an output bit-pattern, gives the probability that it will be read.
- Example:  
 $0.316|00\rangle + 0.447|01\rangle + 0.548|10\rangle + 0.632|11\rangle$   
–The probability to read the rightmost bit as 0 is  $|0.316|^2 + |0.548|^2 = 0.4$
- Measurement **during** a computation changes the state of the system but can be used in some cases to increase efficiency (measure and halt or continue).

Classical  
Versus  
Quantum



# Classical vs. Quantum Circuits

- Goal: Fast, low-cost implementation of useful algorithms using standard components (gates) and design techniques

- Classical Logic Circuits

- Circuit behavior is governed implicitly by classical physics
- Signal states are simple bit vectors, e.g.  $X = 01010111$
- Operations are defined by Boolean Algebra
- No restrictions exist on copying or measuring signals
- Small well-defined sets of universal gate types, e.g. {NAND}, {AND,OR,NOT}, {AND,NOT}, etc.
- Well developed CAD methodologies exist
- Circuits are easily implemented in fast, scalable and macroscopic technologies such as CMOS

# Classical vs. Quantum Circuits

- Quantum Logic Circuits

- Circuit behavior is governed explicitly by quantum mechanics
- Signal states are vectors interpreted as a **superposition** of binary “qubit” vectors with complex-number coefficients

$$|\Psi\rangle = \sum_{i=0}^{2^n-1} c_i |i_{n-1}i_{n-1}\dots i_0\rangle$$

- Operations are defined by linear algebra over Hilbert Space and can be represented by unitary matrices with complex elements
- Severe restrictions exist on **copying** and **measuring** signals
- Many universal gate sets exist but the best types are not obvious
- Circuits must use microscopic technologies that are slow, fragile, and not yet scalable, e.g., NMR

# Quantum Circuit Characteristics

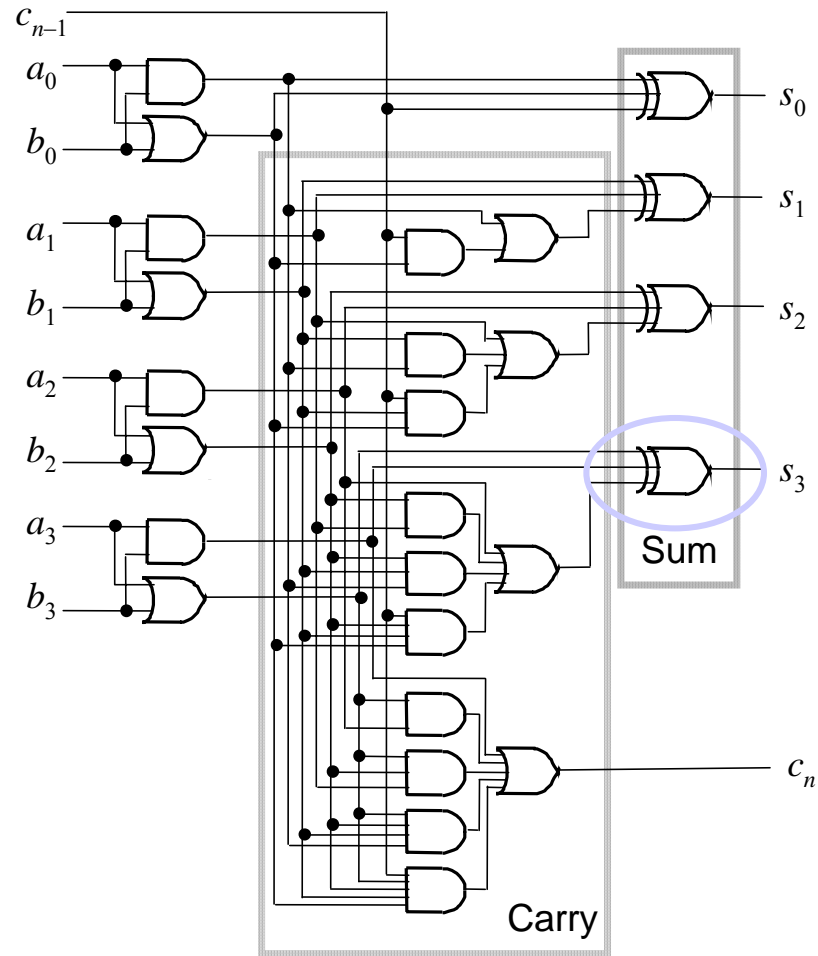
- Unitary Operations
  - Gates and circuits must be reversible (information-lossless)
    - Number of output signal lines = Number of input signal lines
    - The circuit function must be a bijection, implying that output vectors are a permutation of the input vectors
  - **Classical** logic behavior can be represented by permutation matrices
  - **Non-classical** logic behavior can be represented including **state sign** (phase) and **entanglement**

# Quantum Circuit Characteristics

- Quantum Measurement
  - Measurement yields only one state  $X$  of the superposed states
  - Measurement also makes  $X$  the new state and so *interferes with computational processes*
  - $X$  is determined with some **probability**, implying uncertainty in the result
  - States cannot be copied (“cloned”), implying that **signal fanout is not permitted**
  - Environmental interference can cause a **measurement-like state collapse (decoherence)**

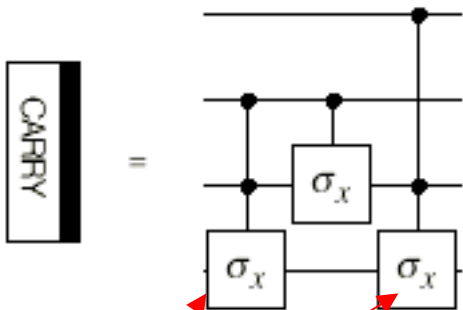
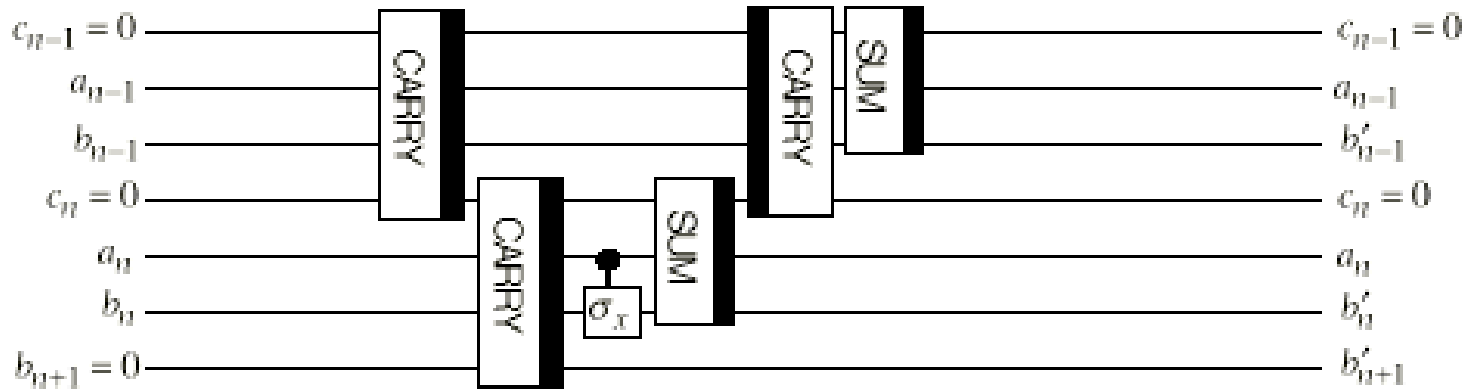
# Classical vs. Quantum Circuits

Classical adder

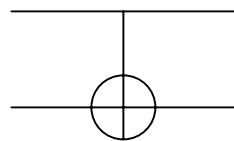
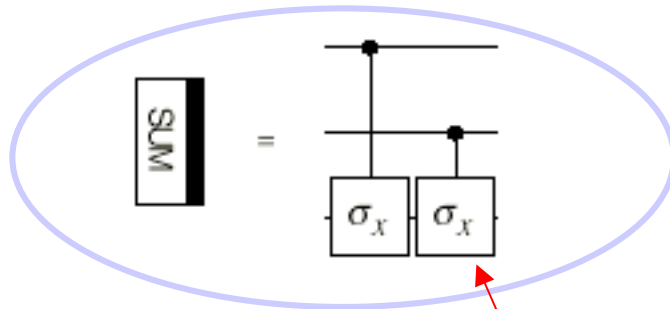


# Classical vs. Quantum Circuits

Quantum adder



Controlled-controlled  $\sigma_x$  is the same as Toffoli



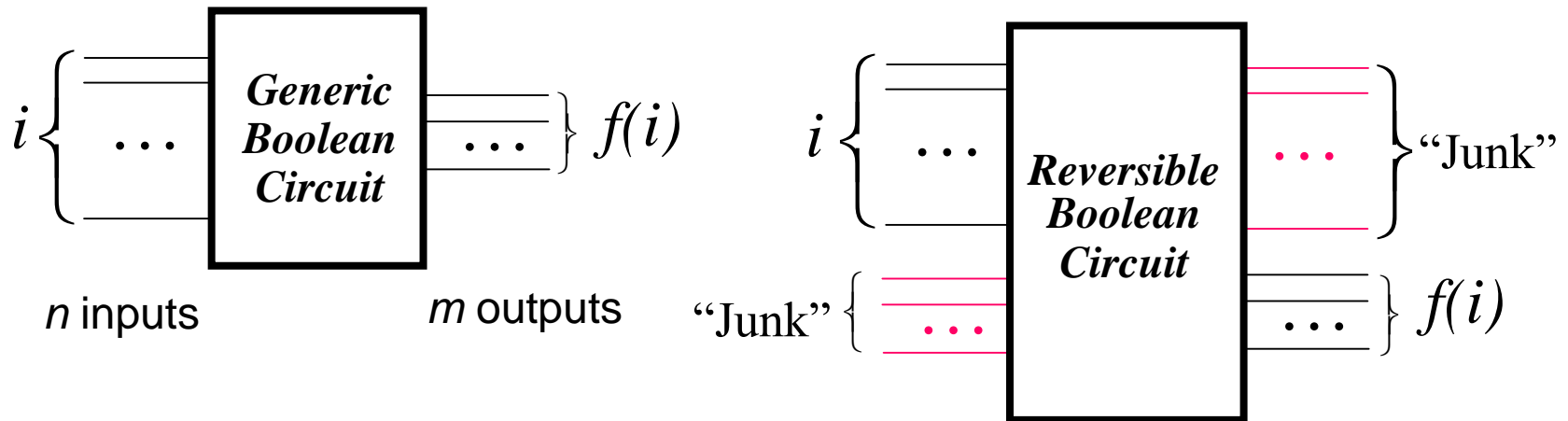
Controlled  $\sigma_x$  is the same as Feynman

- Here we use Pauli rotations notation.
- Controlled  $\sigma_x$  is the same as controlled NOT

# Reversible Circuits

# Reversible Circuits

- Reversibility was studied around 1980 motivated by power minimization considerations
- Bennett, Toffoli et al. showed that any classical logic circuit  $C$  can be made reversible with modest overhead

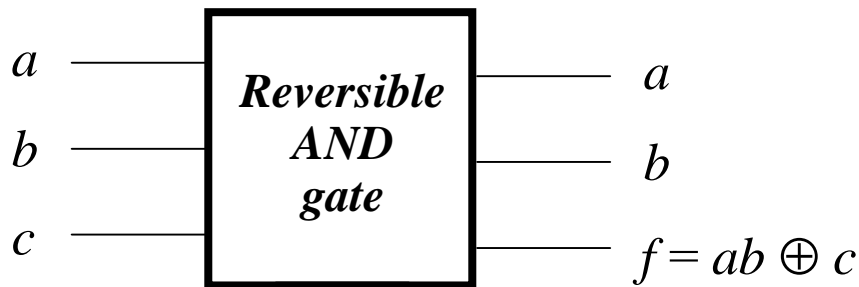




# Reversible Circuits

- How to make a given  $f$  reversible
  - Suppose  $f: i \rightarrow f(i)$  has  $n$  inputs  $m$  outputs
  - Introduce  $n$  extra outputs and  $m$  extra inputs
  - Replace  $f$  by  $f_{\text{rev}}: i, j \rightarrow i, f(i) \oplus j$  where  $\oplus$  is XOR

- Example 1:  $f(a,b) = \text{AND}(a,b)$

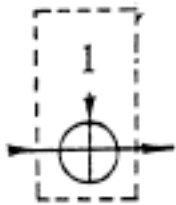


$a$	$b$	$c$	$a$	$b$	$f$
0	0	0	0	0	0
0	0	1	0	0	1
0	1	0	0	1	0
0	1	1	0	1	1
1	0	0	1	0	0
1	0	1	1	0	1
1	1	0	1	1	1
1	1	1	1	1	0

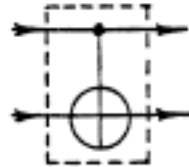
- This is the well-known Toffoli gate, which realizes AND when  $c = 0$ , and NAND when  $c = 1$ .

# Reversible Circuits

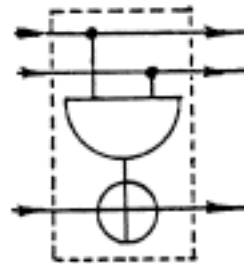
- Reversible gate family [Toffoli 1980]



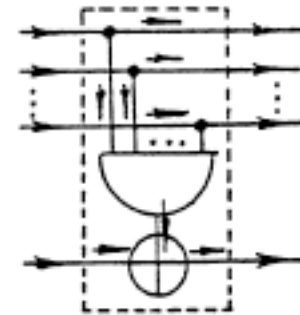
NOT



XOR/FAN-OUT



AND/NAND  
(Toffoli gate)



generalized AND/NAND

- Every Boolean function has a reversible implementation using Toffoli gates.
- There is no universal reversible gate with fewer than three inputs

Quantum

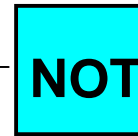
Gates

# Quantum Gates

- **One-Input gate: NOT**

- Input state:  $c_0|0\rangle + c_1|1\rangle$

- Output state:  $c_1|0\rangle + c_0|1\rangle$

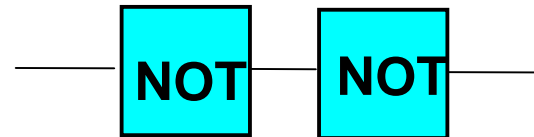


- Pure states are mapped thus:  $|0\rangle \rightarrow |1\rangle$  and  $|1\rangle \rightarrow |0\rangle$

- Gate operator (matrix) is  $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$   $|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$   $|1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$

- As expected:

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$



# Quantum Gates

- **One-Input gate: “Square root of NOT”**

- Some matrix elements are imaginary

- Gate operator (matrix):

$$\begin{pmatrix} i/\sqrt{1/2} & 1/\sqrt{1/2} \\ 1/\sqrt{1/2} & i/\sqrt{1/2} \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} i & 1 \\ 1 & i \end{pmatrix}$$

- We find:

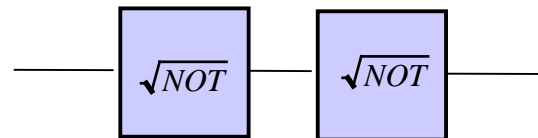
$$\frac{1}{\sqrt{2}} \begin{pmatrix} i & 1 \\ 1 & i \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} i \\ 1 \end{pmatrix} \quad \text{so } |0\rangle \rightarrow |0\rangle \text{ with probability } |i/\sqrt{2}|^2 = 1/2$$

$$\text{and } |0\rangle \rightarrow |1\rangle \text{ with probability } |1/\sqrt{2}|^2 = 1/2$$

Similarly, this gate randomizes input  $|1\rangle$

- But concatenation of two gates eliminates the randomness!

$$\frac{1}{2} \begin{pmatrix} i & 1 \\ 1 & i \end{pmatrix} \begin{pmatrix} i & 1 \\ 1 & i \end{pmatrix} = \begin{pmatrix} 0 & i \\ i & 0 \end{pmatrix}$$



**Other variant of square root of not - we do not use complex numbers  
- only real numbers**

## A square-root-of-NOT (SRN) gate

$$\begin{bmatrix} \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{bmatrix}$$

- Applied once to a classical state, this ~randomizes the value of the qubit.
- Applied twice in a row, this is ~equivalent to NOT:

$$\begin{bmatrix} \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{bmatrix} * \begin{bmatrix} \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix}$$

# Quantum Gates

- **One-Input gate: Hadamard**

$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad \text{---} \boxed{\text{H}} \text{---}$$

- Maps  $|0\rangle \rightarrow 1/\sqrt{2} |0\rangle + 1/\sqrt{2} |1\rangle$  and  $|1\rangle \rightarrow 1/\sqrt{2} |0\rangle - 1/\sqrt{2} |1\rangle$ .
- Ignoring the normalization factor  $1/\sqrt{2}$ , we can write  
 $|x\rangle \rightarrow (-1)^x |x\rangle - |1-x\rangle$

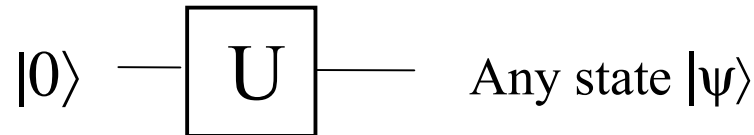
- **One-Input gate: Phase shift**

$$\begin{pmatrix} 1 & 0 \\ 0 & e^{i\phi} \end{pmatrix} \quad \text{---} \boxed{\phi} \text{---}$$

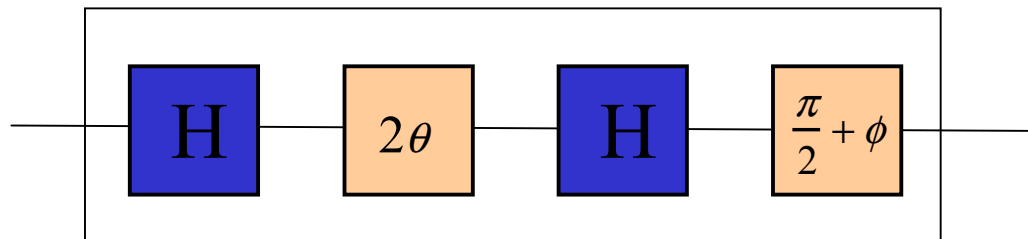
# Quantum Gates

## Universal One-Input Gate Sets

- Requirement:



- **Hadamard** and **phase-shift** gates form a universal gate set of 1-qubit gates, every 1-qubit gate can be built from them.
- *Example:* The following circuit generates  $|\psi\rangle = \cos \theta |0\rangle + e^{i\phi} \sin \theta |1\rangle$  up to a global factor





# Other Quantum Gates

$$\text{Rotation } (U_\theta): \begin{bmatrix} \cos(\theta) & \sin(\theta) \\ -\sin(\theta) & \cos(\theta) \end{bmatrix}$$

$$\text{Hadamard } (H): \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

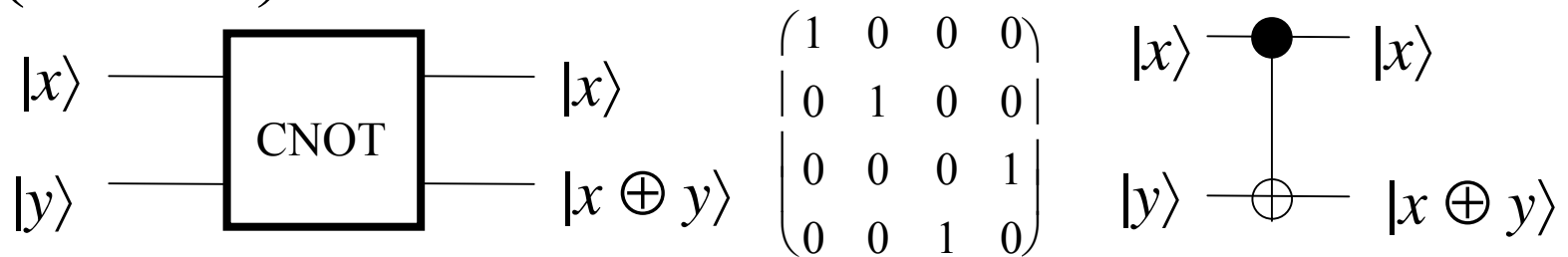
$$\text{CNOT: } \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \quad \text{CPHASE: } \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{i\alpha} \end{bmatrix}$$

There are many small “universal” sets of gates.

Gates must be unitary:  $U^\dagger U = U U^\dagger = I$ , where  $U^\dagger$  is the Hermitean adjoint of  $U$ .

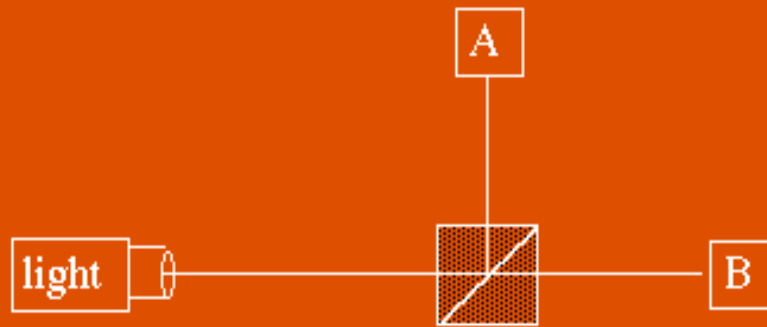
# Quantum Gates

- **Two-Input Gate: Controlled NOT (CNOT)**



- CNOT maps  $|x\rangle|0\rangle \rightarrow |x\rangle|x\rangle$  and  $|x\rangle|1\rangle \rightarrow |x\rangle|\text{NOT } x\rangle$
- $|x\rangle|0\rangle \rightarrow |x\rangle|x\rangle$  *looks like cloning*, but it's not. These mappings are valid only for the pure states  $|0\rangle$  and  $|1\rangle$
- Serves as a “non-demolition” measurement gate

# Polarizing Beam-Splitter CNOT gate from [Cerf, Adami, Kwiat]

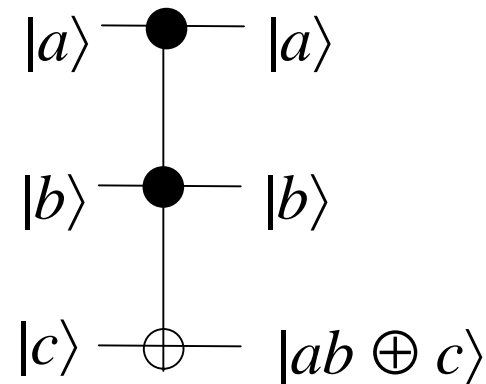


- Two qubits encoded in one photon, one in momentum (direction) and one in polarization.
- Polarization controls change in momentum.
- Cannot be scaled up directly, but demonstrates an implementation of a 2-qubit gate.

# Quantum Gates

- **3-Input gate: Controlled CNOT** (C<sup>2</sup>NOT or Toffoli gate)

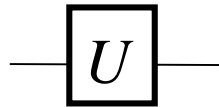
$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$



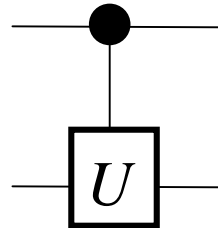
# Quantum Gates

- General controlled gates that control some 1-qubit unitary operation  $U$  are useful

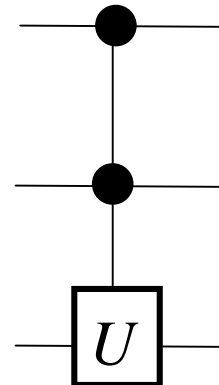
$$\begin{pmatrix} u_{00} & u_{01} \\ u_{10} & u_{11} \end{pmatrix}$$



$U$



$C(U)$



$C^2(U)$

etc.

# Quantum Gates

## Universal Gate Sets

- To implement any unitary operation on  $n$  qubits **exactly** requires an **infinite** number of gate types
- The (infinite) set of all 2-input gates is universal
  - Any  $n$ -qubit unitary operation can be implemented using  $\Theta(n^3 4^n)$  gates [Reck et al. 1994]
- CNOT and the (infinite) set of all 1-qubit gates is universal

# Quantum Gates

## Discrete Universal Gate Sets

- The **error** on implementing  $U$  by  $V$  is defined as

$$E(U, V) = \max_{|\Psi\rangle} \|(U - V)|\Psi\rangle\|$$

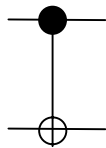
- If  $U$  can be implemented by  $K$  gates, we can simulate  $U$  with a total error less than  $\epsilon$  with a gate overhead that is polynomial in  $\log(K/\epsilon)$
- A discrete set of gate types  $G$  is **universal**, if we can approximate any  $U$  to within any  $\epsilon > 0$  using a sequence of gates from  $G$

# Quantum Gates

## Discrete Universal Gate Set

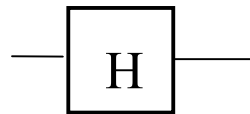
- **Example 1:** Four-member “standard” gate set

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$



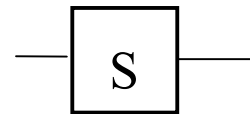
CNOT

$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$



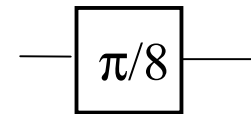
Hadamard

$$\begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}$$



Phase

$$\begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{pmatrix}$$



$\pi/8$  (T) gate

- **Example 2:** {CNOT, Hadamard, Phase, Toffoli}



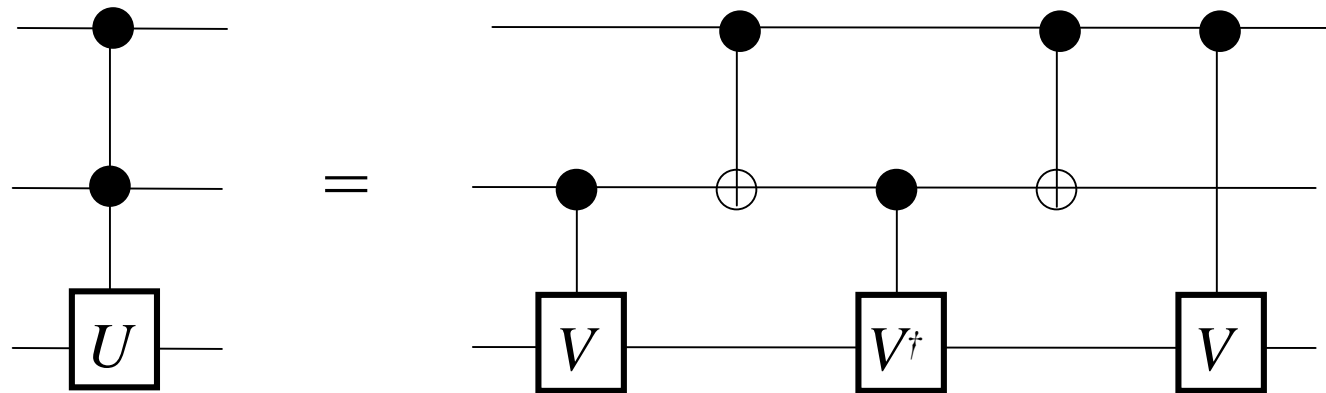
# Quantum Circuits

# Quantum Circuits

- A quantum (combinational) circuit is a sequence of quantum gates, linked by “wires”
- The circuit has fixed “width” corresponding to the number of qubits being processed
- Logic design (classical and quantum) attempts to find circuit structures for needed operations that are
  - Functionally correct
  - Independent of physical technology
  - Low-cost, e.g., use the minimum number of qubits or gates
- Quantum logic design is not well developed!

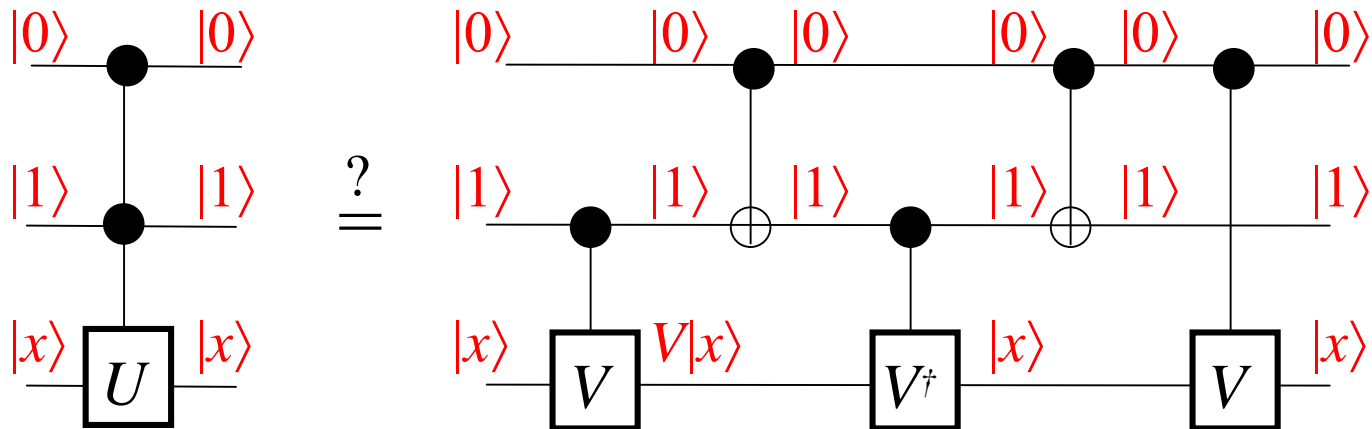
# Quantum Circuits

- Ad hoc designs known for many specific functions and gates
- **Example 1** illustrating a theorem by [Barenco et al. 1995]: Any  $C^2(U)$  gate can be built from CNOTs,  $C(V)$ , and  $C(V^\dagger)$  gates, where  $V^2 = U$



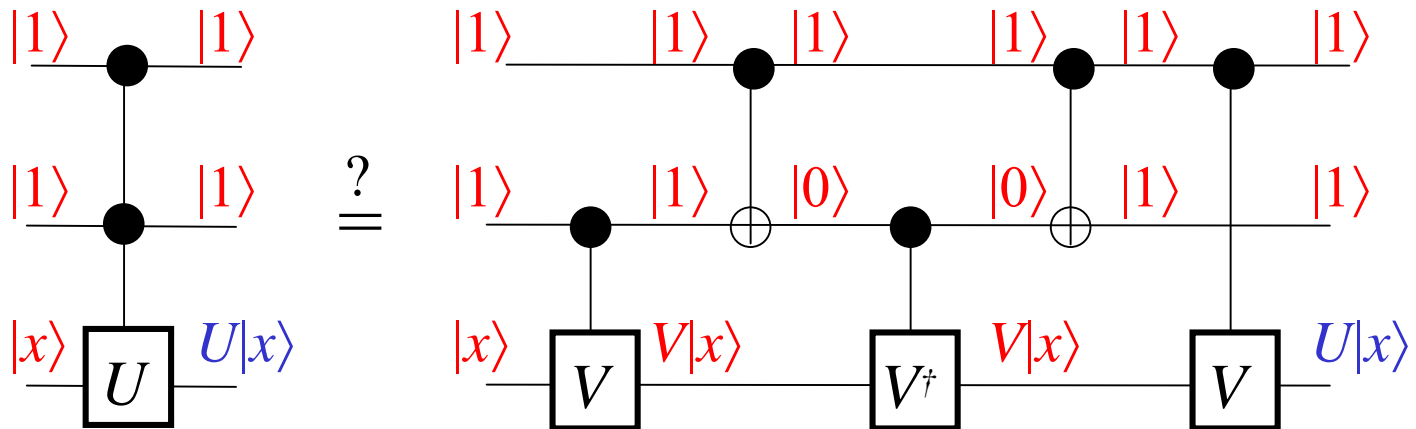
# Quantum Circuits

## Example 1: Simulation



# Quantum Circuits

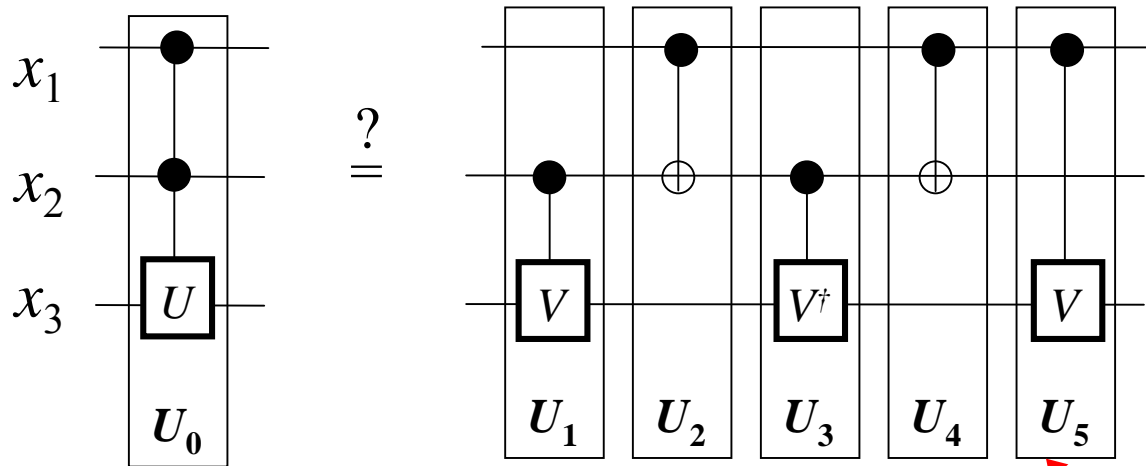
## Example 1: Simulation (contd.)



- *Exercise:* Simulate the two remaining cases

# Quantum Circuits

## Example 1: Algebraic analysis



- Is  $U_0(x_1, x_2, x_3) = U_5 U_4 U_3 U_2 U_1(x_1, x_2, x_3)$   
 $= (x_1, x_2, x_1 x_2 \oplus U(x_3))$  ?

We will verify unitary matrix of Toffoli gate

Observe that the order of matrices  $U_i$  is inverted.

# Quantum Circuits

**Example 1 (contd);**

We calculate the Unitary Matrix  $U_1$  of the first block from left.

$$U_1 = I_1 \otimes C(V)$$

$$= \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \otimes \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & v_{00} & v_{01} \\ 0 & 0 & v_{10} & v_{11} \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & v_{00} & v_{01} & 0 & 0 & 0 & 0 \\ 0 & 0 & v_{10} & v_{11} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{00} & v_{01} \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{10} & v_{11} \end{pmatrix}$$

Unitary matrix of a wire

Kronecker since this is a parallel connection

Unitary matrix of a controlled V gate (from definition)

# Quantum Circuits

**Example 1 (contd);**

We calculate the Unitary Matrix  $U_2$  of the second block from left.

$$U_2 = U_4 = CNOT(x_1, x_2) \otimes I_1$$

$$= \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \otimes \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}$$

Unitary matrix of  
CNOT or Feynman  
gate with EXOR down

As we can check in the schematics, the Unitary Matrices  $U_2$  and  $U_4$  are the same



# Quantum Circuits

**Example 1 (contd);**

$$U_2 = U_4 = CNOT(x_1, x_2) \otimes I_1$$

$$= \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \otimes \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}$$

# Quantum Circuits

## Example 1 (contd);

- $U_5$  is the same as  $U_1$  but has  $x_1$  and  $x_2$  permuted (tricky!)
- It remains to evaluate the product of five  $8 \times 8$  matrices  $U_5 U_4 U_3 U_2 U_1$  using the fact that  $VV^\dagger = I$  and  $VV = U$

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & v_{00} & v_{01} & 0 & 0 \\ 0 & 0 & 0 & 0 & v_{10} & v_{11} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{00} & v_{01} \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}
 \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}
 \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & v_{00} & v_{10} & 0 & 0 & 0 & 0 \\ 0 & 0 & v_{01} & v_{11} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{00} & v_{10} \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{01} & v_{11} \end{pmatrix}
 \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}
 \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & v_{00} & v_{01} & 0 & 0 & 0 & 0 \\ 0 & 0 & v_{10} & v_{11} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{00} & v_{01} \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{10} & v_{11} \end{pmatrix}$$

$$= \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{00}v_{00} + v_{10}v_{10} & v_{00}v_{01} + v_{10}v_{11} \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{01}v_{00} + v_{11}v_{10} & v_{01}v_{01} + v_{11}v_{11} \end{pmatrix} = U_0$$

# Quantum Circuits

**Example 1** (contd);

– We calculate matrix  $U_3$

$$\begin{vmatrix} 1 & 0 \\ 0 & 1 \end{vmatrix} \otimes \begin{vmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & \mathbf{v}_{00} & \mathbf{v}_{10} \\ 0 & 0 & \mathbf{v}_{01} & \mathbf{v}_{11} \end{vmatrix}$$

**=**

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & \mathbf{v}_{00} & \mathbf{v}_{10} & 0 & 0 & 0 & 0 \\ 0 & 0 & \mathbf{v}_{01} & \mathbf{v}_{11} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & \mathbf{v}_{00} & \mathbf{v}_{10} \\ 0 & 0 & 0 & 0 & 0 & 0 & \mathbf{v}_{01} & \mathbf{v}_{11} \end{pmatrix}$$

This is a hermitian matrix, so we transpose and next calculate complex conjugates, we denote complex conjugates by bold symbols

# Quantum Circuits

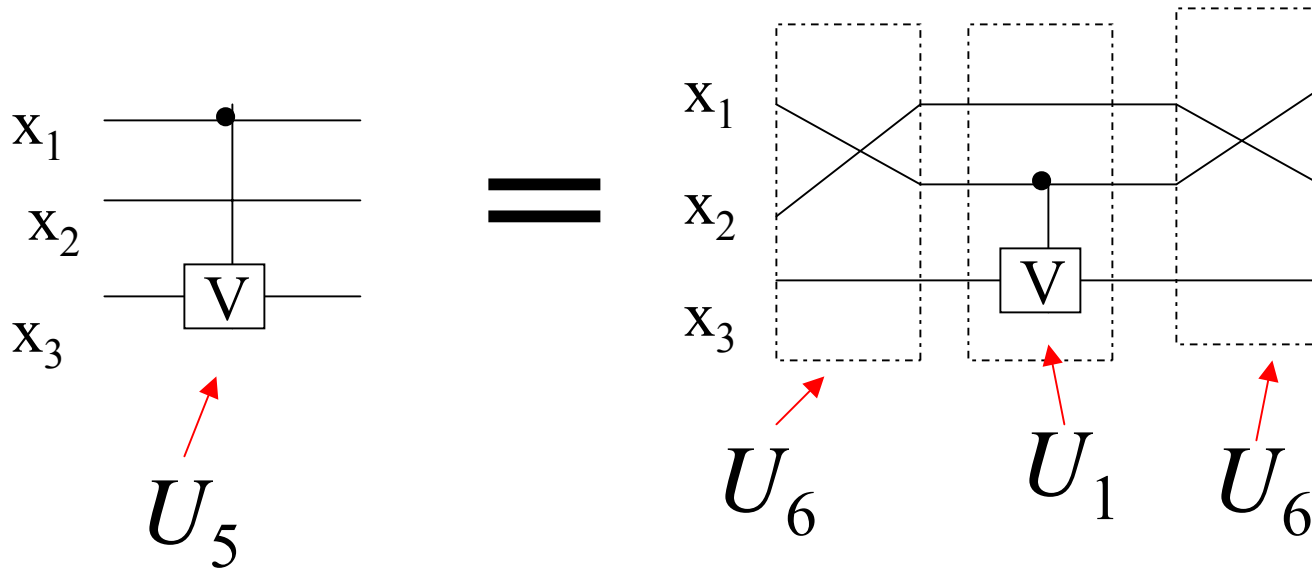
## Example 1 (contd);

- $U_5$  is the same as  $U_1$  but has  $x_1$  and  $x_2$  permuted because in  $U_1$  black dot is in variable  $x_2$  and in  $U_5$  black dot is in variable  $x_1$
- This can be also checked by definition, see next slide.

$$U_5 = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & v_{00} & v_{01} & 0 & 0 \\ 0 & 0 & 0 & 0 & v_{10} & v_{11} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{00} & v_{01} \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{10} & v_{11} \end{pmatrix}$$

# Quantum Circuits

**Example 1** (here we explain in detail how to calculate  $U_5$ )



$U_6$  is calculated as a Kronecker product of  $U_7$  and  $I_1$

$U_7$  is a unitary matrix of a swap gate

$$U_5 = U_6 U_1 U_6$$

# Quantum Circuits

## Example 1 (contd);

- It remains to evaluate the product of five 8 x 8 matrices  $U_5 U_4 U_3 U_2 U_1$  using the fact that  $VV^\dagger = I$  and  $VV = U$

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & v_{00} & v_{01} & 0 & 0 \\ 0 & 0 & 0 & 0 & v_{10} & v_{11} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{00} & v_{01} \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{10} & v_{11} \end{pmatrix}
 \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}
 \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & \mathbf{v}_{00} & \mathbf{v}_{10} & 0 & 0 & 0 & 0 \\ 0 & 0 & \mathbf{v}_{01} & \mathbf{v}_{11} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & \mathbf{v}_{00} & \mathbf{v}_{10} \\ 0 & 0 & 0 & 0 & 0 & 0 & \mathbf{v}_{01} & \mathbf{v}_{11} \end{pmatrix}
 \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}
 \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & v_{00} & v_{01} & 0 & 0 & 0 & 0 \\ 0 & 0 & v_{10} & v_{11} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{00} & v_{01} \\ 0 & 0 & 0 & 0 & 0 & 0 & v_{10} & v_{11} \end{pmatrix}$$



$U_1$

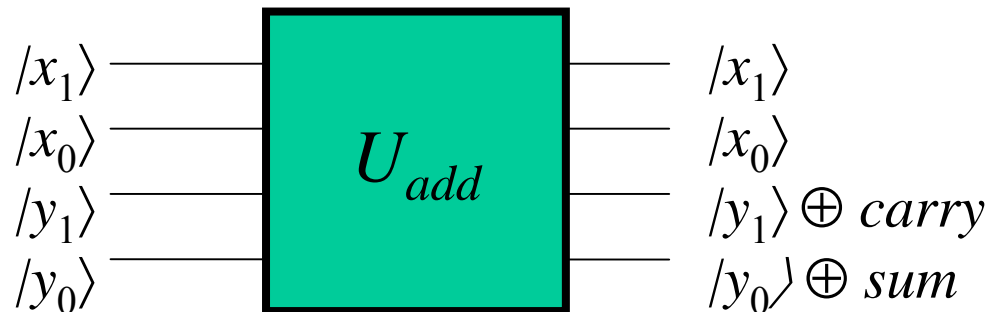
$$= \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & \mathbf{v}_{00}v_{00} + \mathbf{v}_{10}v_{10} & \mathbf{v}_{00}v_{01} + \mathbf{v}_{10}v_{11} \\ 0 & 0 & 0 & 0 & 0 & 0 & \mathbf{v}_{01}v_{00} + \mathbf{v}_{11}v_{10} & \mathbf{v}_{01}v_{01} + \mathbf{v}_{11}v_{11} \end{pmatrix} = U_0$$

# Quantum Circuits

- **Implementing a Half Adder**

- *Problem:* Implement the classical functions  $sum = x_1 \oplus x_0$  and  $carry = x_1 x_0$

- **Generic design:**







# Quantum Circuits

- **Half Adder.** Specific (reduced) design

