

GRASP: A Search Algorithm for Propositional Satisfiability

> EE878C Homework #2 2002/11/1 KAIST, EECS ICS Lab Lee, Dongsoo







Sat in a Nutshell

 Given a Boolean formula, find a variable assignment such that the formula evaluates to 1, or prove that no such assignment exists.

$$F = (a + b)(a' + b' + c)$$

 For n variable,s there are 2<sup>n</sup> possible truth assignments to be checked.



NP-Complete problem.









- Conjunctive Normal Form
  - F = (a+b)(a'+b'+c)
  - Simple representation (more efficient data structures)
- Logic circuit representation
  - Circuits have structural and direction information
- Circuit CNF conversion is straightforward







# DLL Algorithm

#### Davis, Logemann and Loveland

- M. Davis, G. Logemann and D. Loveland, "A Machine Program for Theorem-Proving", Communications of ACM, Vol. 5, No. 7, pp. 394-397, 1962
- Basic framework for many modern SAT solvers
- Also known as DPLL for historical reasons







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#### **Basic DLL Procedure - DFS**

(a' + b + c)(a + c + d) (a + c + d') (a + c' + d) (a + c' + d') (b' + c' + d) (a' + b + c') (a' + b' + c)







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#### **Basic DLL Procedure - DFS**

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(a' + b + c) (a + c + d) (a + c + d') (a + c' + d) (a + c' + d') (b' + c' + d) (a' + b + c') (a' + b' + c)







(a' + b + c)
(a + c + d)
(a + c + d')
(a + c' + d)
(a + c' + d')
(b' + c' + d)
(a' + b + c')
(a' + b' + c)











(a' + b + c)
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(a' + b' + c)













Implication Graph



Conflict!







(a' + b + c)		0	
(a + c + d)		$\checkmark$	
(a + c + d')	1	( b )	
(a + c' + d)	(		
(a + c' + d')	4		
(b' + c' + d)	( c	)	
(a' + b + c')			
(a' + b' + c)			
		(a + c + d)	<b>`</b>
			) Openfilianti
Implica	tion Graph		Conflict!
		(a + c + d') $(d=0)$	)

















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Conflict!







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Conflict!





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Implications and Boolean Constraint Propagation

- Implication
  - A variable is forced to be assigned to be True or False based on previous assignments.
- Unit clause rule (rule for elimination of one literal clauses)
  - An unsatisfied clause is a <u>unit clause if it has exactly one</u> unassigned literal.

(a +b'+ c)(b + c')(a' +c')

a = T, b = T, c is unassigned

Satisfied Literal Unsatisfied Literal Unassigned Literal

- The unassigned literal is implied because of the unit clause.
- Boolean Constraint Propagation (BCP)
  - Iteratively apply the unit clause rule until there is no unit clause available
- Workhorse of DLL based algorithms.







Marques-Silva and Sakallah [SS96,SS99]

- J.P.Marques-Silva and K.A.Sakallah, "GRASP A New Search Algorithm for Satisfiability," Proc. ICCAD 1996. (49 citations)
- Incorporates conflict driven learning and non-chronological backtracking
- Practical SAT instances can be solved in reasonable time
- Bayardo and Schrag's ReISAT also proposed conflict driven learning [BS97]





















x1 + x4

x7 + x8 + x10'x7 + x10 + x12'

















































































x1 + x3' + x8' x1 + x8 + x12 x2 + x11 x7' + x3' + x9 x7' + x8 + x9' x7 + x8 + x10' x7 + x10 + x12'

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Conflict Clause

# Significantly prune the search speace Learned clause is useful forever

Useful in generating future conflict clauses



