



Fast Fourier  
Transform  
S. Schüppel

Introduction  
FFT  
Altera's FFT  
IP-Core  
Literature

# The Fast Fourier Transform

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# Introduction

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Twiddle Factor

FFT

Altera's FFT  
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## Use of FFT algorithm

- originally discovered by Gauß in 1804
- 'rediscovery' by Danielson and Lanczos in 1942
- generally known through Cooley / Tukey in mid 60s
- DFT:  $O(n^2)$  vs. FFT:  $O(n \log_2(n))$

$$X = \sum_{k=0}^{N-1} x_k e^{-i\left(\frac{2\pi}{N}\right)kn} \quad (1)$$

- use  $N = 2^p$  number of samples
- divide and conquer algorithm



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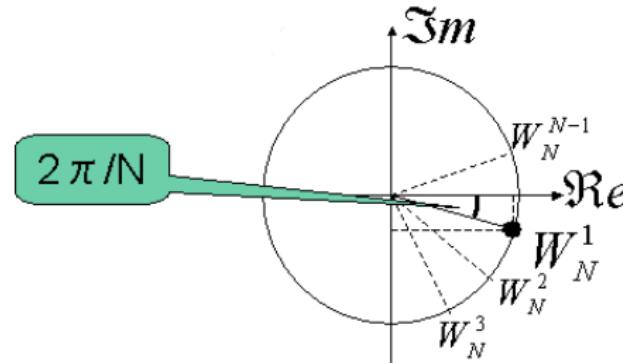
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## Twiddle Factor

$$e^{-i\frac{2\pi}{N}} = W_N \quad (2)$$

$$X(n) = \sum_{k=0}^{N-1} x_k e^{-i(\frac{2\pi}{N})kn} = \sum_{k=0}^{N-1} x_k W_N^{kn} \quad (3)$$

- Euler's Law:  $e^{i\theta} = \sin(\theta) + i \cos(\theta) = \sin(\frac{2\pi}{N}) - i \cos(\frac{2\pi}{N})$





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## Divide and Conquer

$$\begin{aligned} X(n) &= \sum_{k=0}^{(N/2)-1} x_{2k} e^{-i(\frac{2\pi}{N})(2k)n} + \sum_{k=0}^{(N/2)-1} x_{2k+1} e^{-i(\frac{2\pi}{N})(2k+1)n} \\ &= \sum_{k=0}^{(N/2)-1} x_{2k} e^{-i(\frac{2\pi}{N})(2k)n} + e^{-i(\frac{2\pi}{N})n} \sum_{k=0}^{(N/2)-1} x_{2k+1} e^{-i(\frac{2\pi}{N})(2k)n} \\ &= \sum_{k=0}^M x_{2k} W_M^{kn} + W_N^n \sum_{k=0}^M x_{2k+1} W_M^{kn} \quad \left(M = \frac{N}{2}\right) \end{aligned} \quad (4)$$

## Solution through recursion

$$N = 1 \implies X(0) = x_0 W_1^0 = x_0$$



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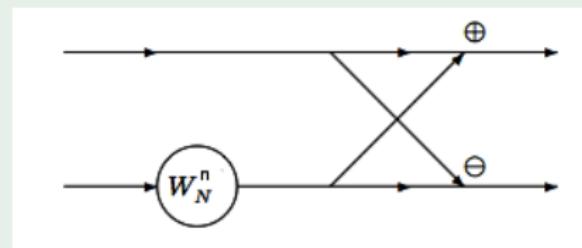
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## Basic Building Block - Butterfly

$$X(n) = X_{2k} + W_N^n X_{2k+1} \quad \left( M = \frac{N}{2} \right) \quad (5)$$

$$X(n + M) = X_{2k} - W_N^n X_{2k+1} \quad (6)$$



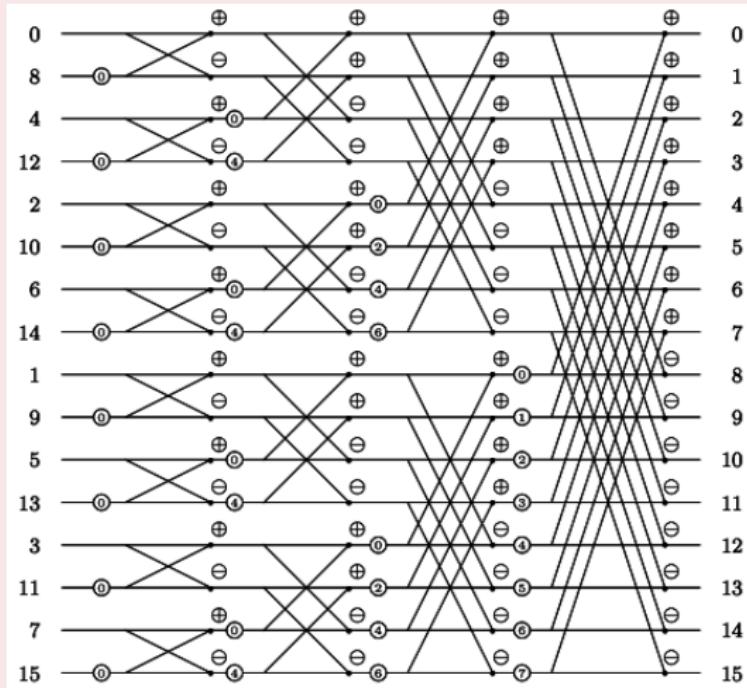


# Fast Fourier Transform - N = 16

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bit-reversed order !





# FFT IP-Cores

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## Why using IP-Cores?

- Easy to Implement
- Free (Altera)
- Optimized in Speed and Area

## Altera Features

- Matlab Model
- radix-4 and mixed radix-4/2
- different Input and Output orders
- transformation length  $2^m \quad 6 \leq m \leq 14$

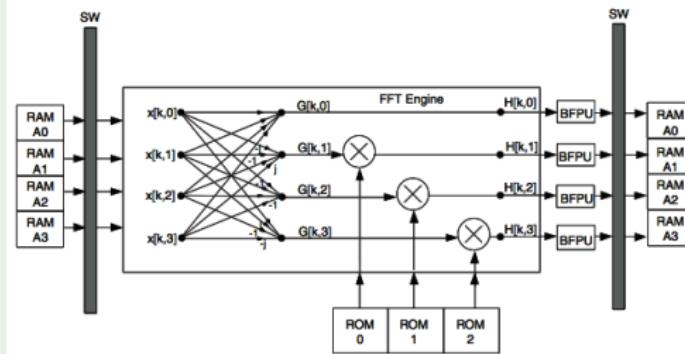
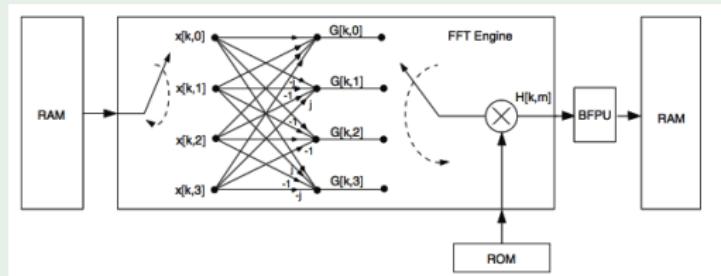


# Inside of the Altera Core

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## Single and Quad radix-4 Output





# Recommended Literature

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- "Fourier Transform for Pedestrians"  
published by Springer Verlag