

Optimization Techniques

Genetic Algorithms

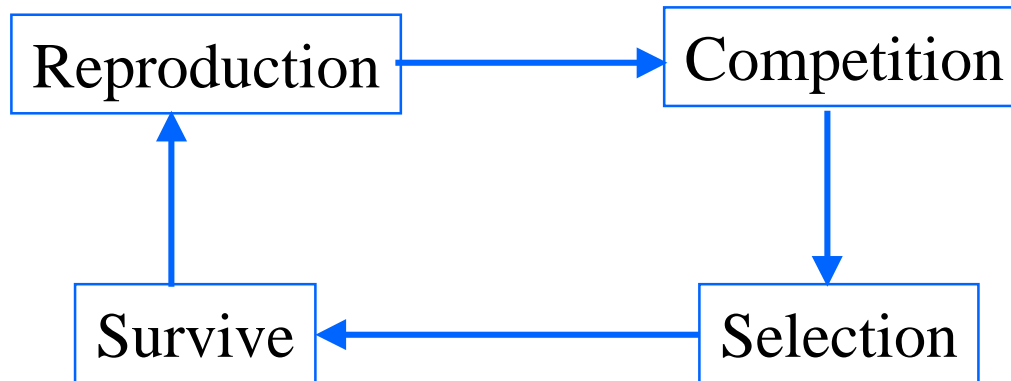
And other approaches for similar applications

Optimization Techniques

- Mathematical Programming
- Network Analysis
- Branch & Bound
- Genetic Algorithm
- Simulated Annealing
- Tabu Search

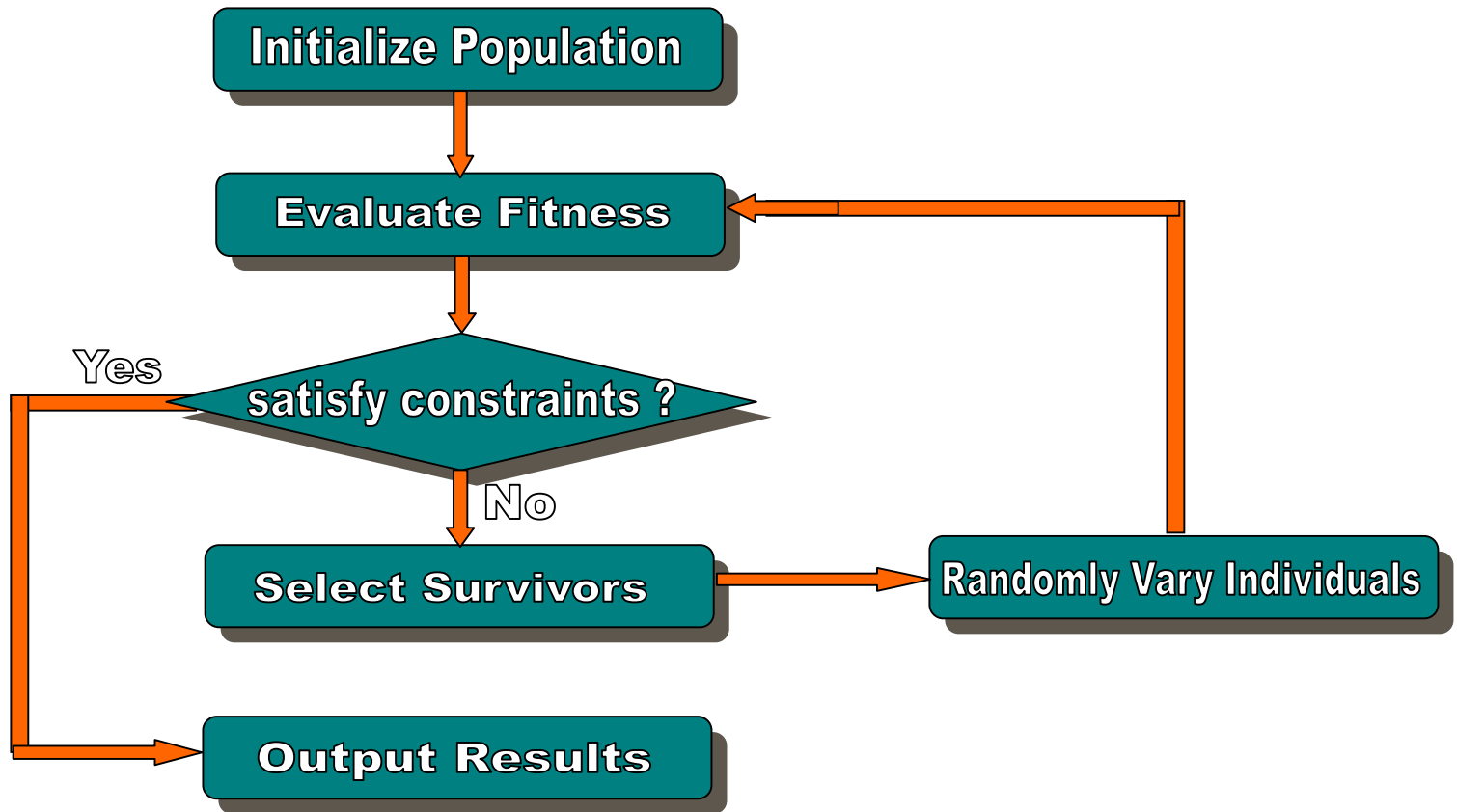
Genetic Algorithm

- Based on Darwinian Paradigm



- Intrinsicly a robust search and optimization mechanism

Conceptual Algorithm



Genetic Algorithm Introduction 1

- Inspired by **natural evolution**
- **Population** of individuals
 - Individual is feasible solution to problem
- Each individual is characterized by a **Fitness function**
 - Higher fitness is better solution
- Based on their fitness, parents are selected to reproduce **offspring** for a new **generation**
 - Fitter individuals have more chance to reproduce
 - New generation has same size as old generation; old generation dies
- Offspring has **combination** of properties of two parents
- If well designed, population will **converge** to optimal solution

Algorithm

BEGIN

Generate initial population;

Compute fitness of each individual;

REPEAT /* New generation */

FOR population_size / 2 DO

Select two parents from old generation;

/* biased to the fitter ones */

Recombine parents for two offspring;

Compute fitness of offspring;

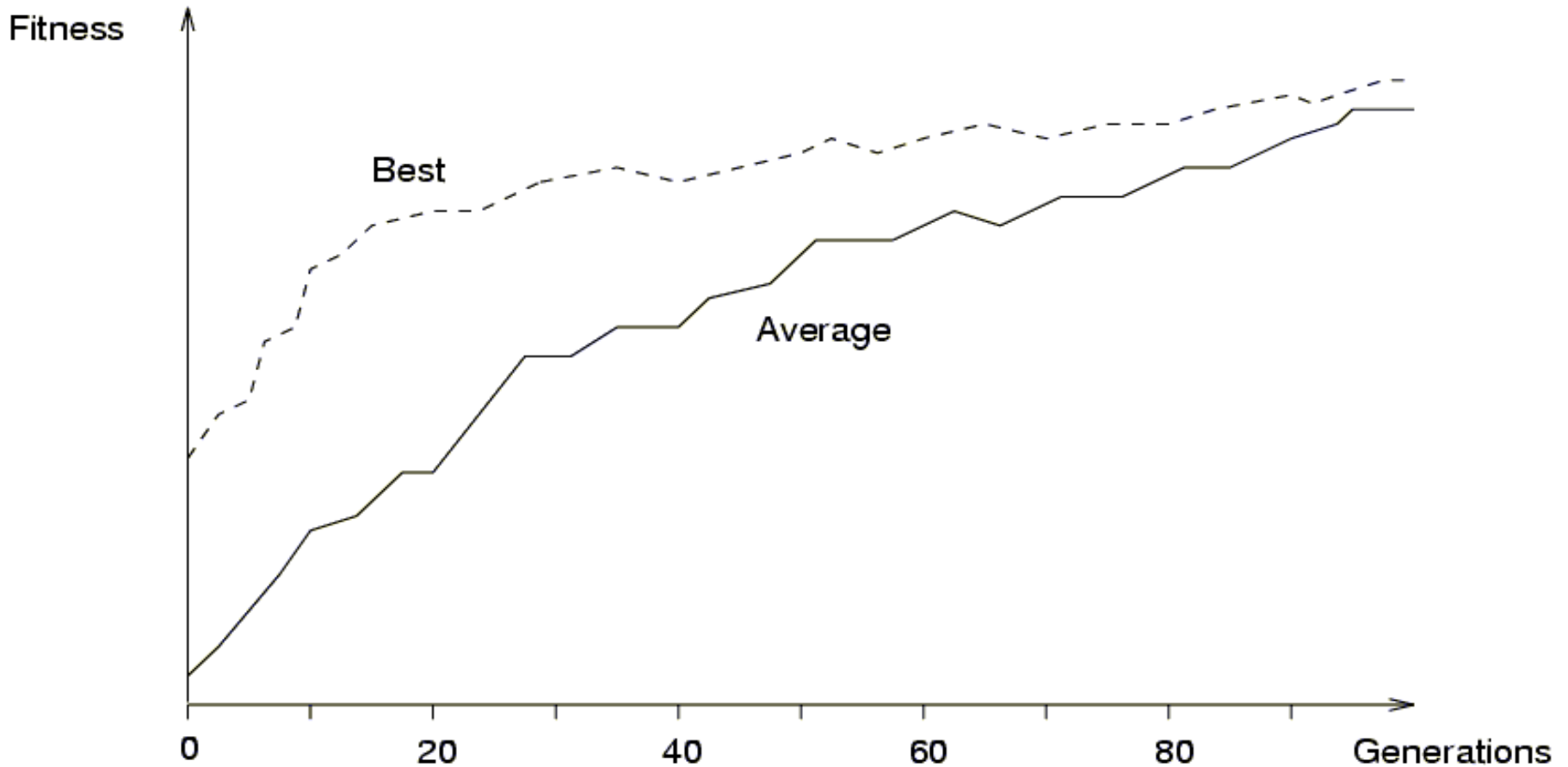
Insert offspring in new generation

END FOR

UNTIL population has converged

END

Example of convergence



Introduction 2

- Reproduction mechanisms have no knowledge of the problem to be solved
- Link between genetic algorithm and problem:
 - Coding
 - Fitness function

Basic principles 1

- Coding or Representation
 - String with all parameters
- Fitness function
 - Parent selection
- Reproduction
 - Crossover
 - Mutation
- Convergence
 - When to stop

Basic principles 2

- An individual is characterized by a set of parameters: **Genes**
 - The genes are joined into a string: **Chromosome**
-
- The chromosome forms the **genotype**
 - The genotype contains all information to construct an organism: the **phenotype**
-
- **Reproduction** is a “dumb” process on the chromosome of the **genotype**
 - **Fitness** is measured in the real world (‘struggle for life’) of the **phenotype**

Coding

- Parameters of the solution (**genes**) are concatenated to form a string (**chromosome**)
- All kind of **alphabets** can be used for a chromosome (numbers, characters), but generally a **binary alphabet** is used
- **Order** of genes on chromosome can be important
- Generally many **different codings** for the parameters of a solution are possible
- **Good coding is probably the most important factor for the performance of a GA**
- In many cases many possible chromosomes do not code for feasible solutions

Genetic Algorithm

- Encoding
- Fitness Evaluation
- Reproduction
- Survivor Selection

Encoding

- Design alternative → individual (chromosome)
- Single design choice → gene
- Design objectives → fitness

Example

- Problem
 - Schedule n jobs on m processors such that the maximum span is minimized.

Design alternative: job i ($i=1,2,\dots,n$) is assigned to processor j ($j=1,2,\dots,m$)

Individual: A n -vector \mathbf{x} such that $x_i = 1, \dots, \text{or } m$

Design objective: minimize the maximal span

Fitness: the maximal span for each processor

Reproduction

- Reproduction operators
 - Crossover
 - Mutation

Reproduction

- Crossover

- Two parents produce two offspring
- There is a chance that the chromosomes of the two parents are copied unmodified as offspring
- There is a chance that the chromosomes of the two parents are randomly recombined (crossover) to form offspring
- Generally the chance of crossover is between 0.6 and 1.0

- Mutation

- There is a chance that a gene of a child is changed randomly
- Generally the chance of mutation is low (e.g. 0.001)

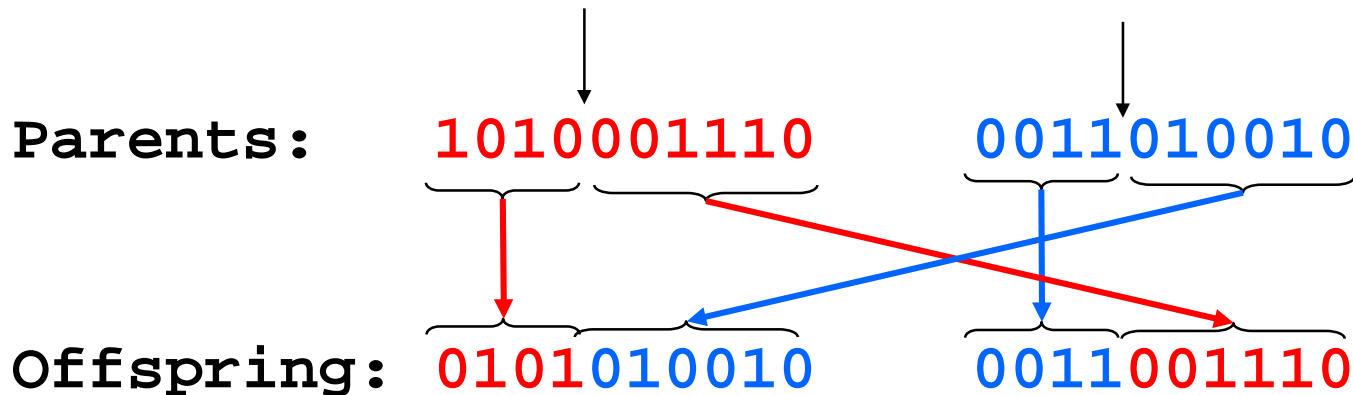
Reproduction Operators

- Crossover
 - Generating offspring from two selected parents
 - | Single point crossover
 - | Two point crossover (Multi point crossover)
 - | Uniform crossover

One-point crossover 1

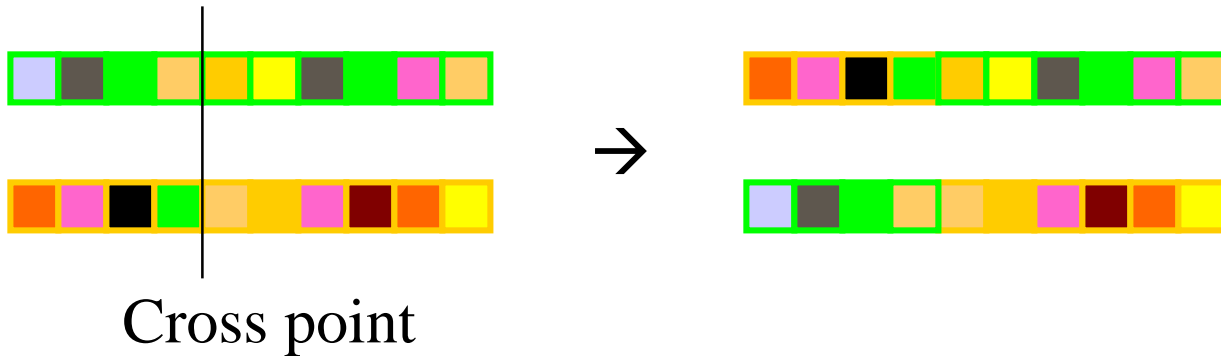
- Randomly one position in the chromosomes is chosen
- Child 1 is head of chromosome of parent 1 with tail of chromosome of parent 2
- Child 2 is head of 2 with tail of 1

Randomly chosen position

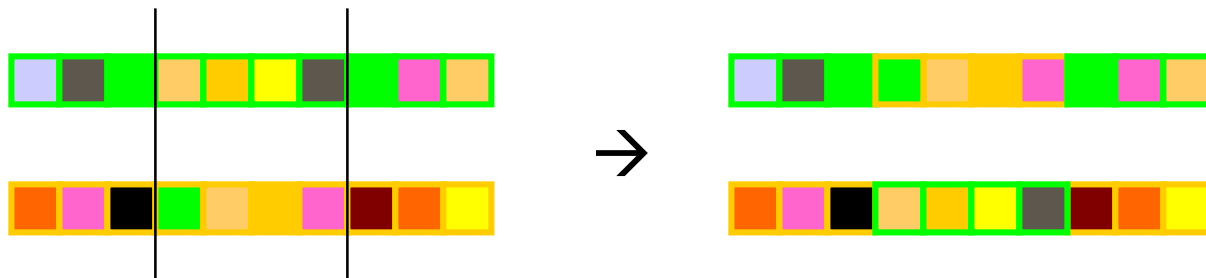


Reproduction Operators comparison

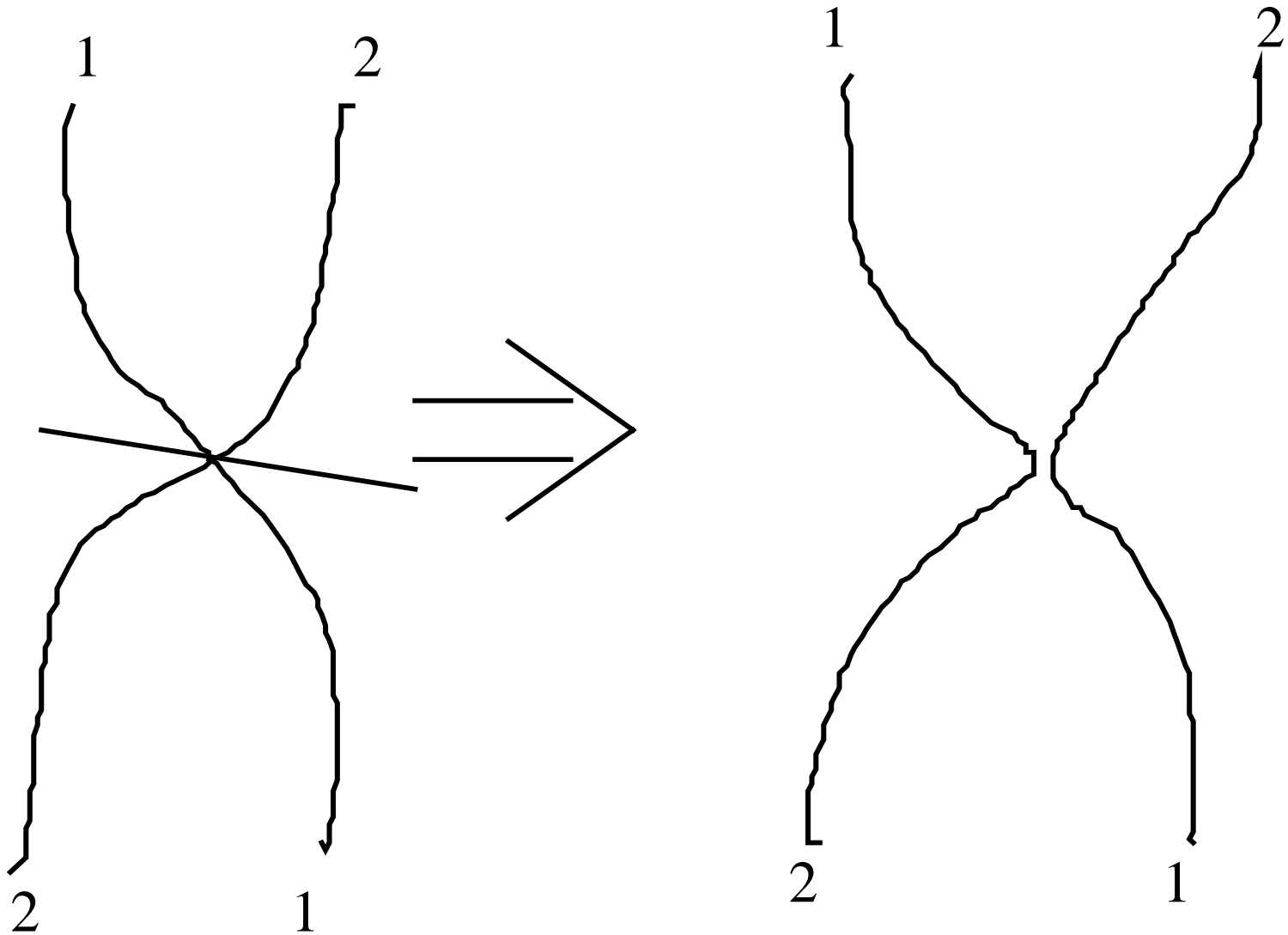
- Single point crossover



- Two point crossover (Multi point crossover)

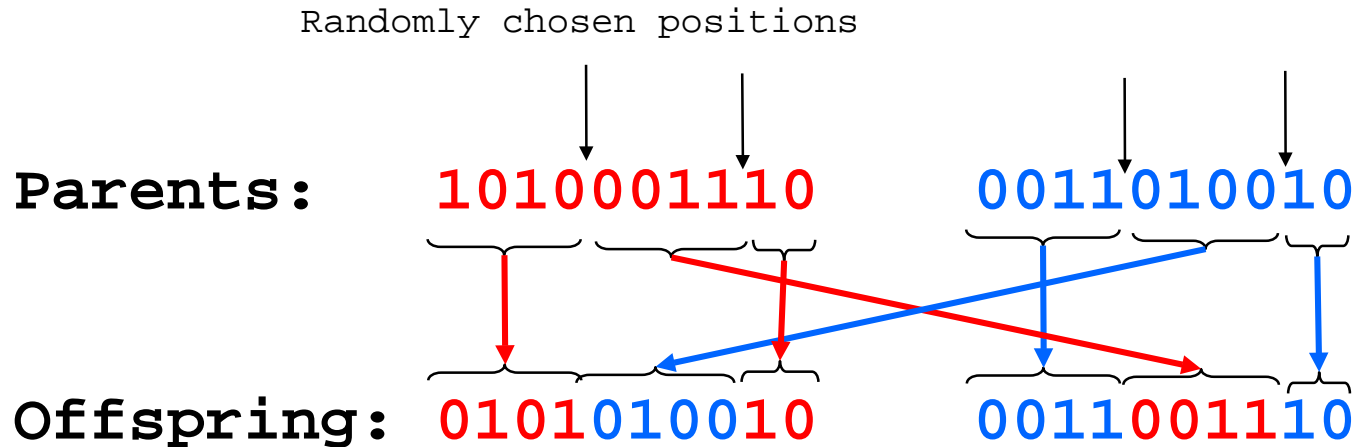


One-point crossover - Nature



Two-point crossover

- Randomly two positions in the chromosomes are chosen
- Avoids that genes at the head and genes at the tail of a chromosome are always split when recombined



Uniform crossover

- A random mask is generated
- The mask determines which bits are copied from one parent and which from the other parent
- Bit density in mask determines how much material is taken from the other parent (takeover parameter)

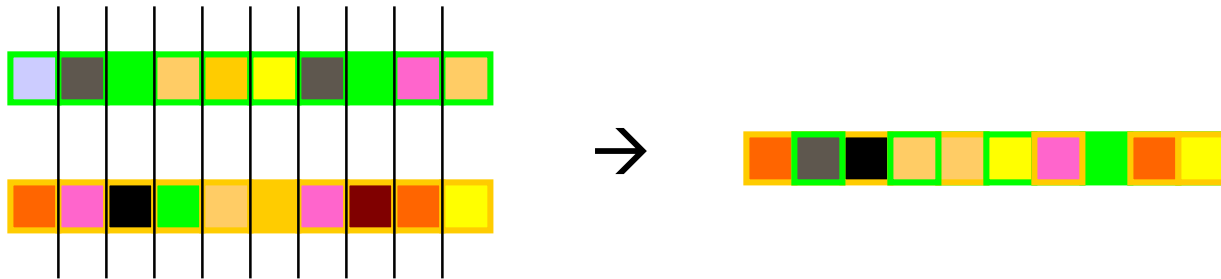
Mask: 0110011000 (Randomly generated)

Parents: 1010001110 0011010010

Offspring: 0011001010 1010010110

Reproduction Operators

- Uniform crossover



- Is uniform crossover better than single crossover point?
 - Trade off between
 - Exploration: introduction of new combination of features
 - Exploitation: keep the good features in the existing solution

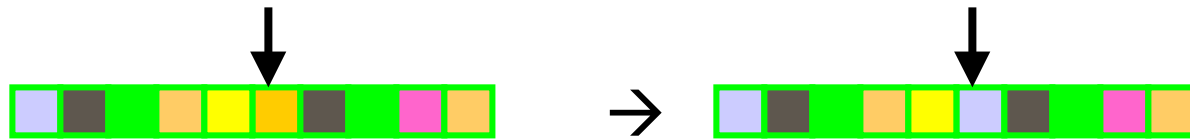
Problems with crossover

- Depending on coding, simple crossovers can have high chance to produce illegal offspring
 - E.g. in TSP with simple binary or path coding, most offspring will be illegal because not all cities will be in the offspring and some cities will be there more than once
- Uniform crossover can often be modified to avoid this problem
 - E.g. in TSP with simple path coding:
 - | Where mask is 1, copy cities from one parent
 - | Where mask is 0, choose the remaining cities in the order of the other parent

Reproduction Operators

- Mutation

- Generating new offspring from single parent



- Maintaining the diversity of the individuals
 - | Crossover can only explore the combinations of the current gene pool
 - | Mutation can “generate” new genes

Reproduction Operators

- Control parameters: **population size, crossover/mutation probability**
 - Problem specific
 - Increase population size
 - | Increase diversity and computation time for each generation
 - Increase crossover probability
 - | Increase the opportunity for recombination but also disruption of good combination
 - Increase mutation probability
 - | Closer to randomly search
 - | Help to introduce new gene or reintroduce the lost gene
- **Varies the population**
 - Usually using crossover operators to recombine the genes to generate the new population, then using mutation operators on the new population

Parent/Survivor Selection

- Strategies
 - Survivor selection
 - | Always keep the best one
 - | Elitist: deletion of the K worst
 - | Probability selection : inverse to their fitness
 - | Etc.

Parent/Survivor Selection

- Too strong fitness selection bias can lead to sub-optimal solution
- Too little fitness bias selection results in unfocused and meandering search

Parent selection

Chance to be selected as parent proportional to fitness

- Roulette wheel

To avoid problems with fitness function

- Tournament

Not a very important parameter

Parent/Survivor Selection

- Strategies
 - Parent selection
 - | Uniform randomly selection
 - | Probability selection : proportional to their fitness
 - | Tournament selection (Multiple Objectives)
 - Build a small comparison set
 - Randomly select a pair with the higher rank one beats the lower one
 - Non-dominated one beat the dominated one
 - Niche count**: the number of points in the population within certain distance, higher the niche count, lower the rank.
 - | Etc.

Others

- Global Optimal
- Parameter Tuning
- Parallelism
- Random number generators

Example of coding for TSP

Travelling Salesman Problem

- Binary
 - Cities are binary coded; chromosome is string of bits
 - | Most chromosomes code for illegal tour
 - | Several chromosomes code for the same tour
- Path
 - Cities are numbered; chromosome is string of integers
 - | Most chromosomes code for illegal tour
 - | Several chromosomes code for the same tour
- Ordinal
 - Cities are numbered, but code is complex
 - All possible chromosomes are legal and only one chromosome for each tour
- Several others

Roulette wheel

- Sum the fitness of all chromosomes, call it T
- Generate a random number N between 1 and T
- Return chromosome whose fitness added to the running total is equal to or larger than N
- Chance to be selected is exactly proportional to fitness

Chromosome:	1	2	3	4	5	6
Fitness:	8	2	17	7	4	11
Running total:	8	10	27	34	38	49
N (1 ≤ N ≤ 49):			23			
Selected:			3			

Tournament

- **Binary tournament**
 - Two individuals are randomly chosen; the fitter of the two is selected as a parent
- **Probabilistic binary tournament**
 - Two individuals are randomly chosen; with a chance p , $0.5 < p < 1$, the fitter of the two is selected as a parent
- **Larger tournaments**
 - n individuals are randomly chosen; the fittest one is selected as a parent
- By changing n and/or p , the GA can be adjusted dynamically

Problems with fitness range

- **Premature convergence**
 - Δ Fitness too large
 - Relatively superfit individuals dominate population
 - Population converges to a local maximum
 - Too much exploitation; too few exploration
- **Slow finishing**
 - Δ Fitness too small
 - No selection pressure
 - After many generations, average fitness has converged, but no global maximum is found; not sufficient difference between best and average fitness
 - Too few exploitation; too much exploration

Solutions for these problems

- Use tournament selection
 - Implicit fitness remapping
- Adjust fitness function for roulette wheel
 - Explicit fitness remapping
 - | Fitness **scaling**
 - | Fitness **windowing**
 - | Fitness **ranking**

Will be explained below

Fitness Function

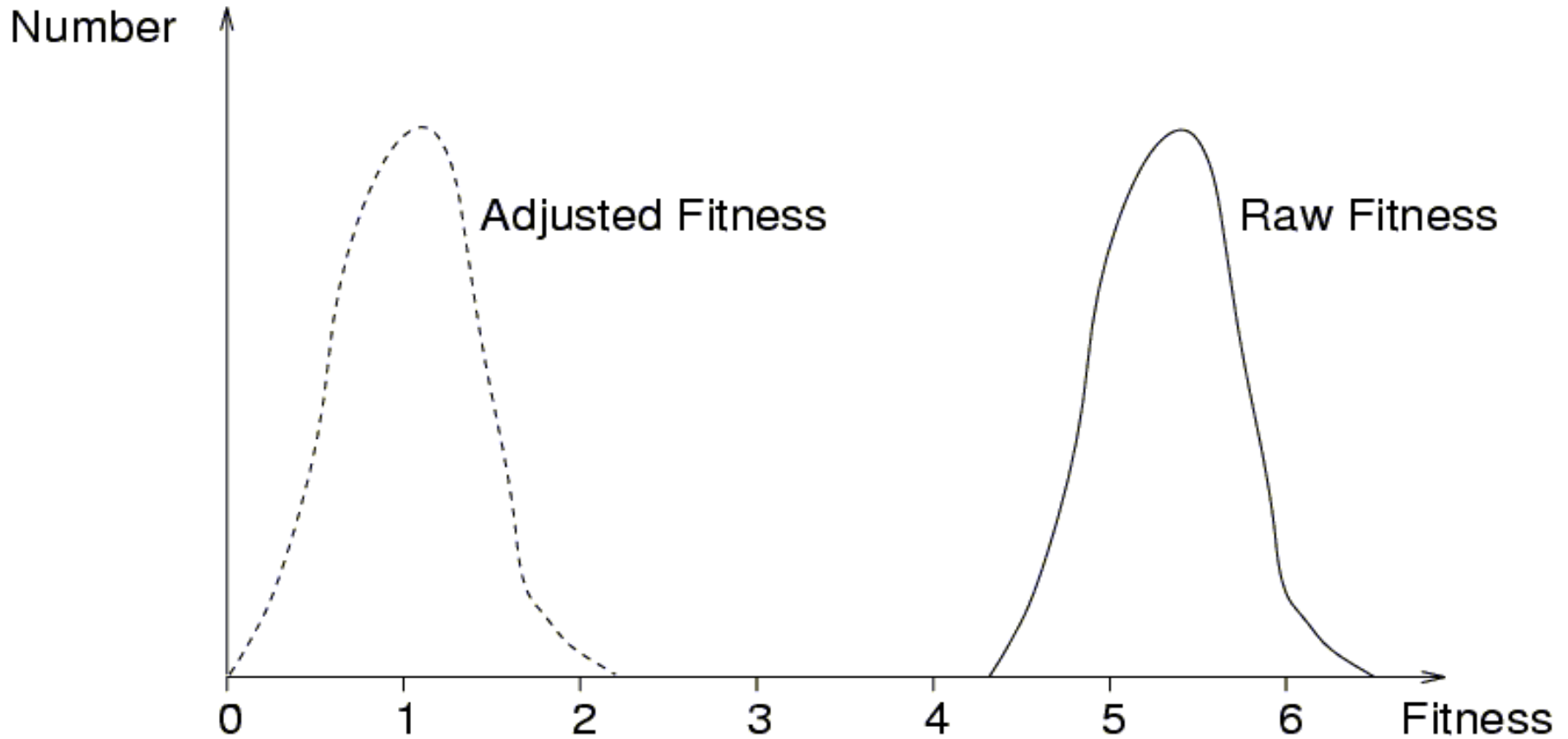
Purpose

- Parent selection
- Measure for convergence
- For Steady state: Selection of individuals to die
- Should reflect the value of the chromosome in some “real” way
- Next to coding the most critical part of a GA

Fitness scaling

- Fitness values are scaled by subtraction and division so that worst value is close to 0 and the best value is close to a certain value, typically 2
 - Chance for the most fit individual is 2 times the average
 - Chance for the least fit individual is close to 0
- Problems when the original maximum is very extreme (**super-fit**) or when the original minimum is very extreme (**super-unfit**)
 - Can be solved by defining a minimum and/or a maximum value for the fitness

Example of Fitness Scaling



Fitness windowing

- Same as window scaling, except the amount subtracted is the minimum observed in the n previous generations, with n e.g. 10
- Same problems as with scaling

Fitness ranking

- Individuals are numbered in order of increasing fitness
- The rank in this order is the adjusted fitness
- Starting number and increment can be chosen in several ways and influence the results

- No problems with super-fit or super-unfit
- Often superior to scaling and windowing

Fitness Evaluation

- A key component in GA
- Time/quality trade off
- Multi-criterion fitness

Multi-Criterion Fitness

- Dominance and indifference
 - For an optimization problem with more than one objective function ($f_i, i=1,2,\dots,n$)
 - given any two solution X_1 and X_2 , then
 - | X_1 dominates X_2 ($X_1 \succ X_2$), if
$$f_i(X_1) \geq f_i(X_2), \text{ for all } i = 1, \dots, n$$
 - | X_1 is indifferent with X_2 ($X_1 \sim X_2$), if X_1 does not dominate X_2 , and X_2 does not dominate X_1

Multi-Criterion Fitness

- Pareto Optimal Set
 - If there exists no solution in the search space which dominates any member in the set P , then the solutions belonging to the set P constitute a global Pareto-optimal set.
 - Pareto optimal front
- Dominance Check

Multi-Criterion Fitness

- Weighted sum

- $F(\mathbf{x}) = W_1f_1(x_1) + W_2f_2(x_2) + \dots + W_nf_n(x_n)$

- *Problems?*

- | *Convex and convex Pareto optimal front*

- Sensitive to the shape of the Pareto-optimal front*

- | *Selection of weights?*

- Need some pre-knowledge*

- Not reliable for problem involving uncertainties*

Multi-Criterion Fitness

- Optimizing single objective

- *Maximize:* $f_k(\mathbf{X})$

- Subject to:*

- $f_j(\mathbf{X}) \leq K_j, \quad j \neq k$

- \mathbf{X} in F where F is the solution space.

Multi-Criterion Fitness

- Weighted sum

- $F(\mathbf{x}) = W_1 f_1(x_1) + W_2 f_2(x_2) + \dots + W_n f_n(x_n)$

- *Problems?*

- | *Convex and convex Pareto optimal front*

- Sensitive to the shape of the Pareto-optimal front*

- | *Selection of weights?*

- Need some pre-knowledge*

- Not reliable for problem involving uncertainties*

Multi-Criterion Fitness

- Preference based weighted sum

(**ISMAUT** *Imprecisely Specific Multiple Attribute Utility Theory*)

- $F(\mathbf{x}) = W_1 f_1(x_1) + W_2 f_2(x_2) + \dots + W_n f_n(x_n)$

- Preference

| Given two known individuals \mathbf{X} and \mathbf{Y} , if we prefer \mathbf{X} than \mathbf{Y} , then

$$F(\mathbf{X}) > F(\mathbf{Y}),$$

that is

$$W_1(f_1(x_1) - f_1(y_1)) + \dots + W_n(f_n(x_n) - f_n(y_n)) > 0$$

Multi-Criterion Fitness

- | All the preferences constitute a linear space

$$W_n = \{W_1, W_2, \dots, W_n\}$$

$$W_1(f_1(x_1) - f_1(y_1)) + \dots + W_n(f_n(x_n) - f_n(y_n)) > 0$$

$$W_1(f_1(z_1) - f_1(p_1)) + \dots + W_n(f_n(z_n) - f_n(p_n)) > 0, \text{ etc}$$

- | For any two new individuals Y' and Y'' , how to determine which one is more preferable?

Multi-Criterion Fitness

$$\text{Min} : \mu = \sum_k w_k [f_k(\mathbf{Y}') - f_k(\mathbf{Y}'')]]$$

$$\text{s.t.} : \quad W_n$$

$$\text{Min} : \mu' = \sum_k w_k [f_k(\mathbf{Y}'') - f_k(\mathbf{Y}')]]$$

$$\text{s.t.} : \quad W_n$$

Multi-Criterion Fitness

Then,

$$\mu > 0 \Rightarrow Y' \succ Y''$$

$$\mu' > 0 \Rightarrow Y'' \succ Y'$$

Otherwise,

$$Y' \sim Y''$$

Construct the dominant relationship among some indifferent ones according to the preferences.

Other parameters of GA 1

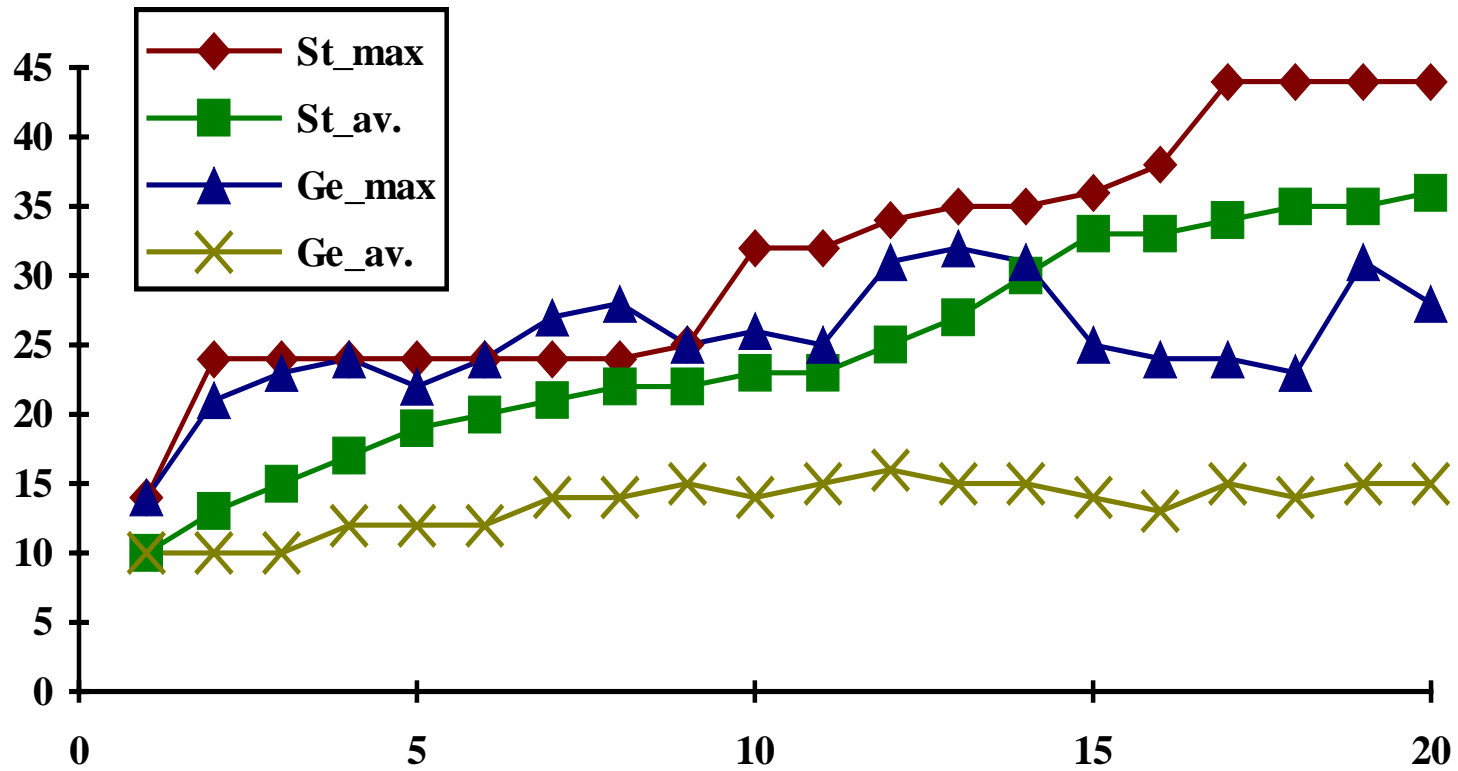
- **Initialization:**
 - Population size
 - Random
 - Dedicated greedy algorithm
- **Reproduction:**
 - Generational: as described before (insects)
 - Generational with elitism: fixed number of most fit individuals are copied unmodified into new generation
 - Steady state: two parents are selected to reproduce and two parents are selected to die; two offspring are immediately inserted in the pool (mammals)

Other parameters of GA 2

- **Stop criterion:**
 - Number of new chromosomes
 - Number of new and unique chromosomes
 - Number of generations
- **Measure:**
 - Best of population
 - Average of population
- **Duplicates**
 - Accept all duplicates
 - Avoid too many duplicates, because that degenerates the population (inteelt)
 - No duplicates at all

Example run

Maxima and Averages of steady state and generational replacement



Simulated Annealing

- What
 - Exploits an analogy between the annealing process and the search for the optimum in a more general system.

Annealing Process

- Annealing Process
 - Raising the temperature up to a very high level (melting temperature, for example), the atoms have a higher energy state and a high possibility to re-arrange the crystalline structure.
 - Cooling down slowly, the atoms have a lower and lower energy state and a smaller and smaller possibility to re-arrange the crystalline structure.

Simulated Annealing

- Analogy
 - Metal \leftrightarrow Problem
 - Energy State \leftrightarrow Cost Function
 - Temperature \leftrightarrow Control Parameter
 - A completely ordered crystalline structure \leftrightarrow the optimal solution for the problem

Global optimal solution can be achieved as long as the cooling process is slow enough.

Metropolis Loop

- The essential characteristic of simulated annealing
- Determining how to randomly explore new solution, reject or accept the new solution at a constant temperature T .
- Finished until equilibrium is achieved.

Metropolis Criterion

- Let
 - \mathbf{x} be the current solution and \mathbf{x}' be the new solution
 - $C(\mathbf{x})$ ($C(\mathbf{x}')$) be the energy state (cost) of \mathbf{x} (\mathbf{x}')
- Probability $P_{\text{accept}} = \exp [-(C(\mathbf{x}) - C(\mathbf{x}')) / T]$
- Let $N = \text{Random}(0, 1)$
- Unconditional accepted if
 - $C(\mathbf{x}') < C(\mathbf{x})$, *the new solution is better*
- Probably accepted if
 - $C(\mathbf{x}') \geq C(\mathbf{x})$, *the new solution is worse . Accepted only when $N < P_{\text{accept}}$*

Algorithm

Initialize initial solution \mathbf{x} , highest temperature T_h , and coolest temperature T_l

$T = T_h$

When the temperature is higher than T_l

While not in equilibrium

Search for the new solution \mathbf{X}'

Accept or reject \mathbf{X}' according to Metropolis Criterion

End

Decrease the temperature T

End

Simulated Annealing

- Definition of solution
- Search mechanism, i.e. the definition of a neighborhood
- Cost-function

Control Parameters

- Definition of equilibrium
 - Cannot yield any significant improvement after certain number of loops
 - A constant number of loops
- Annealing schedule (i.e. How to reduce the temperature)
 - A constant value, $T' = T - T_d$
 - A constant scale factor, $T' = T * R_d$
 - | A scale factor usually can achieve better performance

Control Parameters

- Temperature determination
 - Artificial, without physical significant
 - Initial temperature
 - | 80-90% acceptance rate
 - Final temperature
 - | A constant value, i.e., based on the total number of solutions searched
 - | No improvement during the entire Metropolis loop
 - | Acceptance rate falling below a given (small) value
- Problem specific and may need to be tuned

Example

- Traveling Salesman Problem (TSP)
 - Given 6 cities and the traveling cost between any two cities
 - A salesman need to start from city 1 and travel all other cities then back to city 1
 - Minimize the total traveling cost

Example

- Solution representation
 - An integer list, i.e., $(1, 4, 2, 3, 6, 5)$
- Search mechanism
 - Swap any two integers (except for the first one)
 - $(1, 4, 2, 3, 6, 5) \rightarrow (1, 4, 3, 2, 6, 5)$
- Cost function

Example

- Temperature
 - Initial temperature determination
 - | Around 80% acceptance rate for “bad move”
 - | Determine acceptable ($C_{\text{new}} - C_{\text{old}}$)
 - Final temperature determination
 - | Stop criteria
 - | Solution space coverage rate
 - Annealing schedule
 - | Constant number (90% for example)
 - | Depending on solution space coverage rate

Others

- Global optimal is possible, but near optimal is practical
- Parameter Tuning
 - Aarts, E. and Korst, J. (1989). *Simulated Annealing and Boltzmann Machines*. John Wiley & Sons.
- Not easy for parallel implementation
- Randomly generator

Optimization Techniques

- Mathematical Programming
- Network Analysis
- Branch & Bound
- Genetic Algorithm
- Simulated Annealing
- Tabu Search

Tabu Search

- What
 - Neighborhood search + memory
 - | Neighborhood search
 - | Memory
 - Record the search history
 - Forbid cycling search

Algorithm

- Choose an **initial** solution x
- Find a subset of $N(x)$ the neighbor of x which are not in the tabu list.
- Find the best one (x') in $N(x)$.
- If $F(x') > F(x)$ then set $x=x'$.
- Modify the tabu list.
- If a stopping condition is met **then stop**, else go to the **second step**.

Effective Tabu Search

- Effective Modeling
 - Neighborhood structure
 - Objective function (fitness or cost)
 - | **Example** Graph coloring problem: Find the minimum number of colors needed such that no two connected nodes share the same color.
- **Aspiration criteria**
 - The criteria for overruling the tabu constraints and differentiating the preference of among the neighbors

Effective Tabu Search

- Effective Computing
 - “Move” may be easier to be stored and computed than a completed solution
 - | move: the process of constructing of x' from x
 - Computing and storing the fitness difference may be easier than that of the fitness function.

Effective Tabu Search

- **Effective Memory Use**

- Variable tabu list size

- | For a constant size tabu list

- | Too long: deteriorate the search results

- | Too short: cannot effectively prevent from cycling

- Intensification of the search

- | Decrease the tabu list size

- Diversification of the search

- | Increase the tabu list size

- | Penalize the frequent move or unsatisfied constraints

Example

- A hybrid approach for graph coloring problem
 - R. Dorne and J.K. Hao, *A New Genetic Local Search Algorithm for Graph Coloring*, 1998

Problem

- Given an undirected graph $G=(V,E)$
 - $V=\{v_1, v_2, \dots, v_n\}$
 - $E=\{e_{ij}\}$
- Determine a partition of V in a minimum number of color classes C_1, C_2, \dots, C_k such that for each edge e_{ij} , v_i and v_j are not in the same color class.
- NP-hard

General Approach

- Transform an optimization problem into a decision problem
- Genetic Algorithm + Tabu Search
 - Meaningful crossover
 - Using Tabu search for efficient local search

Encoding

- Individual
 - $(C_{i1}, C_{i2}, \dots, C_{ik})$
- Cost function
 - Number of total conflicting nodes
 - | Conflicting node
 - having same color with at least one of its adjacent nodes
- Neighborhood (move) definition
 - Changing the color of a conflicting node
- Cost evaluation
 - Special data structures and techniques to improve the efficiency

Implementation

- Parent Selection
 - Random
- Reproduction/Survivor
- Crossover Operator
 - Unify independent set (UIS) crossover
 - | Independent set
 - Conflict-free nodes set with the same color
 - | Try to increase the size of the independent set to improve the performance of the solutions

UIS

Unify independent set

parent 1

0	2	1	0	1	2	1	2	0
---	---	---	---	---	---	---	---	---

conflict

⊗	○	⊗	○	○	○	⊗	○	○
---	---	---	---	---	---	---	---	---

parent 2

0	1	2	2	0	2	2	1	1
---	---	---	---	---	---	---	---	---

conflict

○	⊗	⊗	○	○	○	○	○	⊗
---	---	---	---	---	---	---	---	---

child 1

1	2	1	0	1	0	0	2	0
---	---	---	---	---	---	---	---	---

origin

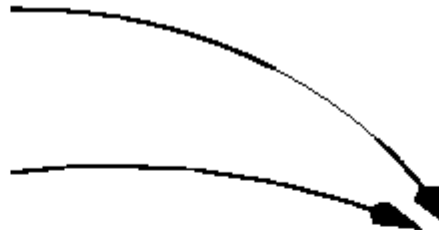
p2	p1		p1	p1	p2	p2	p1	p1
----	----	--	----	----	----	----	----	----

child 2

0	1	2	2	0	1	2	1	2
---	---	---	---	---	---	---	---	---

origin

p2	p1		p2	p2	p1	p2	p2	p1
----	----	--	----	----	----	----	----	----



unions obtained	
$l_{e1,0} = l_{p1,0} + l_{p2,2}$	$l_{e2,0} = l_{p2,0} + l_{p1,1}$
$l_{e1,1} = l_{p1,1} + l_{p2,0}$	$l_{e2,1} = l_{p2,1} + l_{p1,2}$
$l_{e1,2} = l_{p1,2} + l_{p2,1}$	$l_{e2,2} = l_{p2,2} + l_{p1,0}$



Implementation

- Mutation

- With Probability $P_{W'}$ randomly pick neighbor
- With Probability $1 - P_{W'}$ Tabu search

- | Tabu search

- Tabu list

- List of $\{V_i, C_j\}$

- Tabu tenure (the length of the tabu list)

- $$L = a * N_c + \text{Random}(g)$$

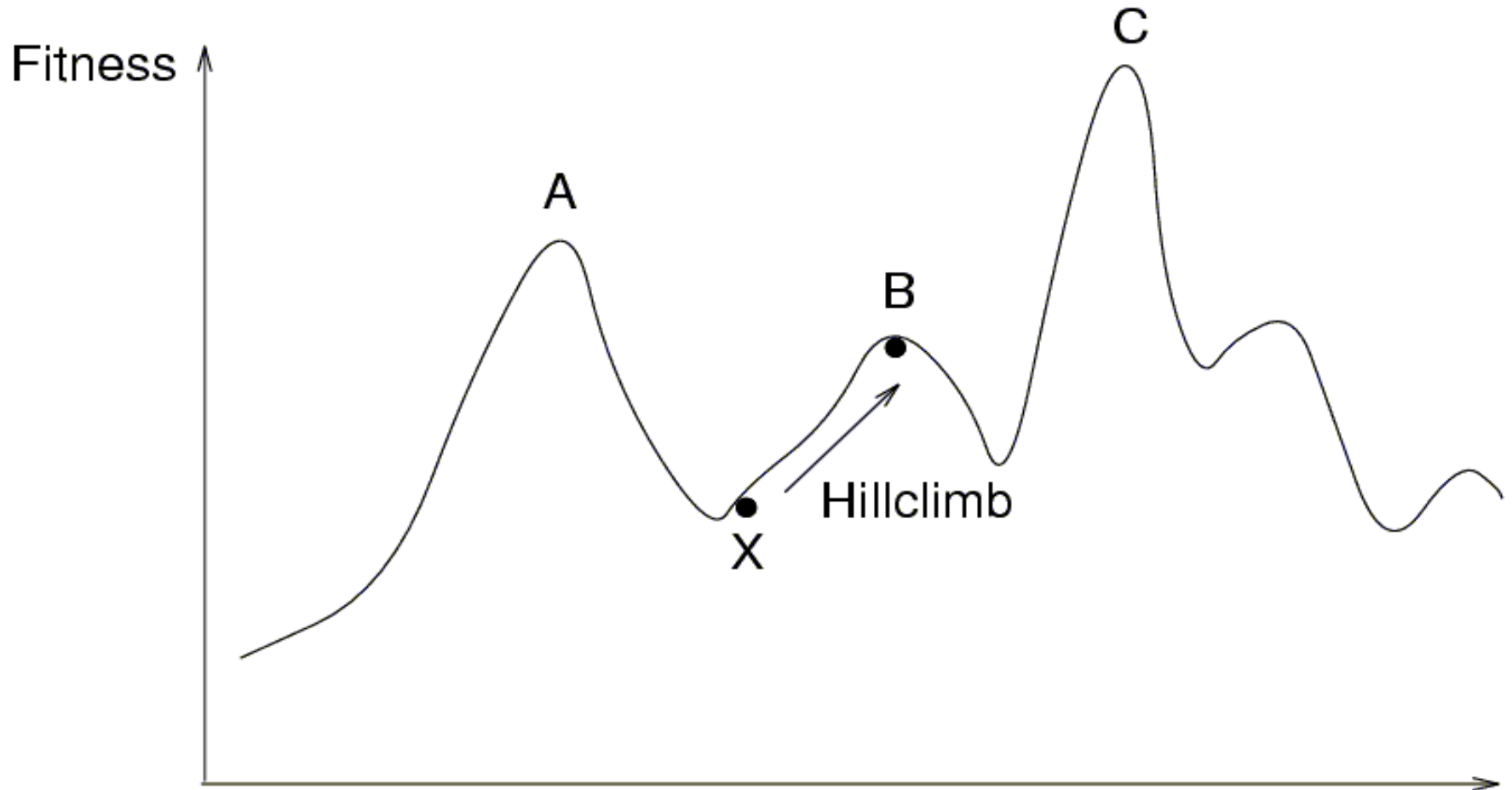
- N_c : Number of conflicted nodes

- a, g : empirical parameters

Summary

- Neighbor Search
- TS prevent being trapped in the local minimum with tabu list
- TS directs the selection of neighbor
- TS cannot guarantee the optimal result
- Sequential
- Adaptive

Hill climbing





Sources

Jaap Hofstede, Beasley, Bull, Martin
Version 2, October 2000

Department of Computer Science & Engineering
University of South Carolina
Spring, 2002