

CSCE790 Topics in Information Technology

# Computational Models (Lecture 5)

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# Review

- Data Flow Graph
  - data dependency
- Control/Data Flow Graph
  - control dependency
- How about a reactive system?

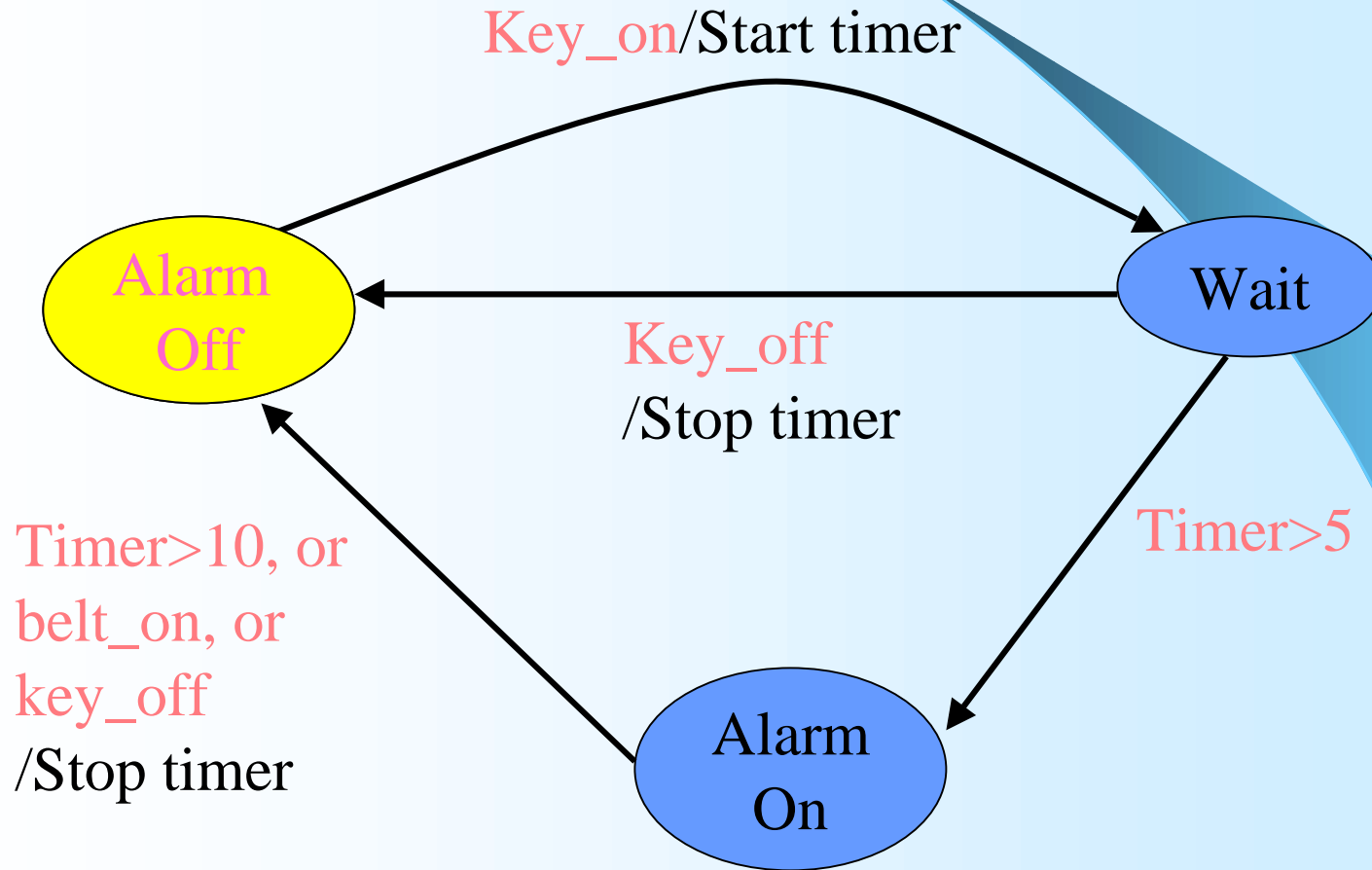
# Finite State Machine

- What ?

*If the driver turns on the key, and does not  
fasten the seat belt within 5 seconds  
then*

*an alarm beeps for 5 seconds,  
or until the driver fastens the seat belt,  
or until the driver turns off the key*

# An FSM



# An FSM (Cont'd)

- **States**
  - Alarm off, Alarm on, Wait
- **Initial State**
  - Alarm off
- **Inputs**
  - Turn on/off the key, fasten the seat belt, timer reads
- **Outputs**
  - Start/stop the timer
- **Start transitions**
  - Alarm off + Turn on the key → Wait
- **Output**
  - Alarm off + Turn on the key → start the timer

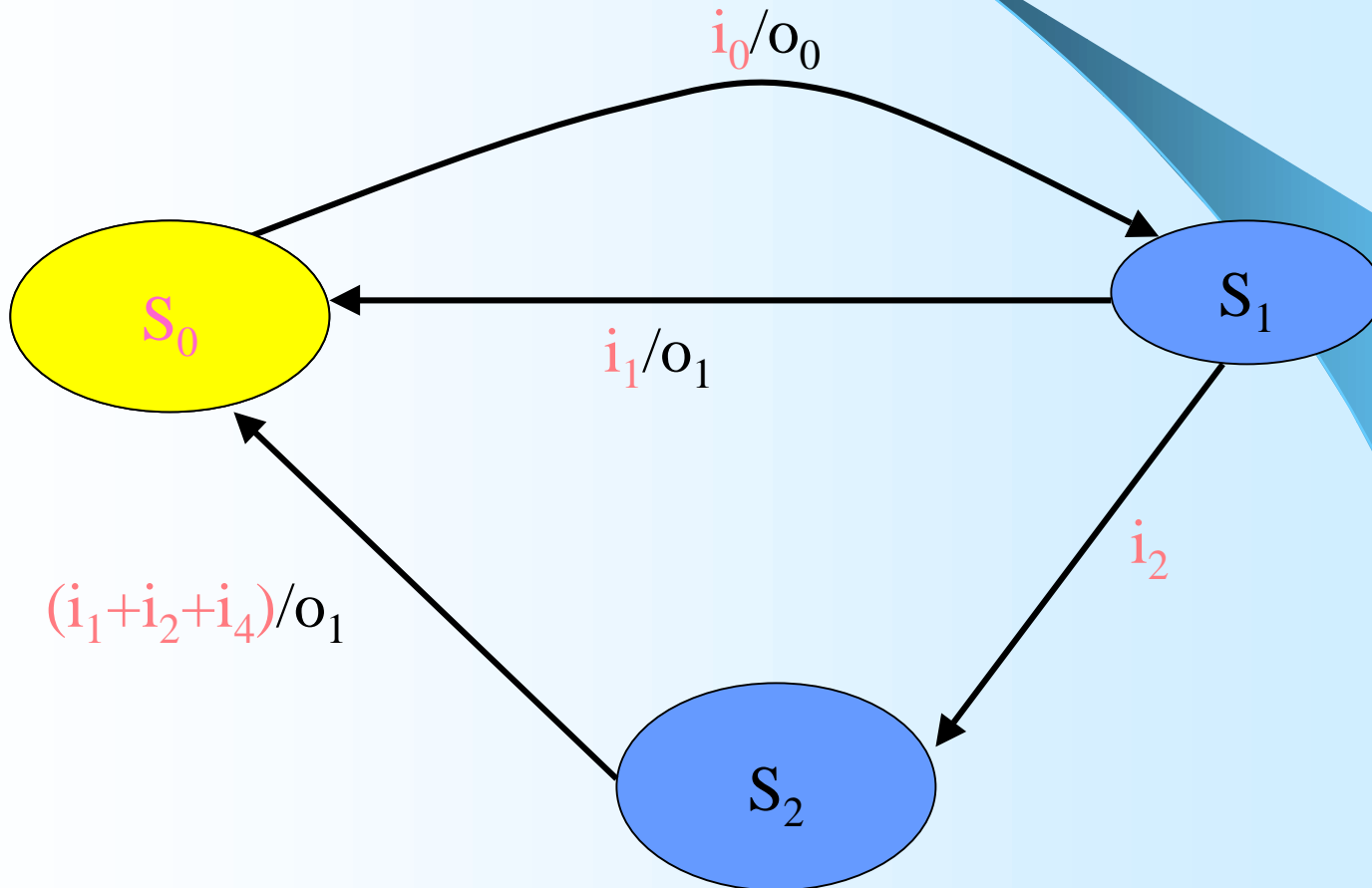
# Finite State Machine

- **FSM = ( S, I, O, s<sub>0</sub>,  $\delta$ ,  $\lambda$  )**
  - **S = {s<sub>0</sub>, s<sub>1</sub>, ..., s<sub>k</sub>}**
  - **I = {i<sub>1</sub>, i<sub>2</sub>, ..., i<sub>m</sub>}**
  - **O = {o<sub>1</sub>, o<sub>2</sub>, ..., o<sub>n</sub>}**
  - **$\delta: S \times I \rightarrow S$  (Transition function)**
  - **$\lambda: S \times I \rightarrow O$  (Output function)**
- **Given an input sequence, an output sequence is produced which is depended on s<sub>0</sub>,  $\delta$ , and  $\lambda$  .**

# Representation

- Given
  - States
    - Alarm off ( $s_0$ ), Alarm on ( $s_1$ ), Wait ( $s_2$ )
  - Initial State
    - Alarm off ( $s_0$ )
  - Inputs
    - Turn on/off the key ( $i_0/i_1$ ), fasten the seat belt ( $i_2$ ), timer > 5 ( $i_3$ ), time > 10 ( $i_4$ )
  - Outputs
    - Start/stop the timer ( $o_0/o_1$ )

# Transition Graph





# Transition Function

- **Transition Function**

$$s_1 = s_0 * i_0 \quad s_0 = s_1 * i_1$$

$$s_2 = s_1 * i_3 \quad s_0 = s_2 * (i_1 + i_2 + i_4)$$

- **Output Function**

$$O_0 = s_0 * i_0 \quad O_1 = s_1 * i_1$$

$$O_1 = s_2 * (i_1 + i_2 + i_4)$$

# Transition Table

State

|       | $S_0$ | $S_1$ | $S_2$ |
|-------|-------|-------|-------|
| $i_0$ | $S_1$ | X     | X     |
| $i_1$ | X     | $S_0$ | $S_0$ |
| $i_2$ | X     | $S_2$ | $S_0$ |
| $i_3$ | X     | X     | X     |
| $i_4$ | X     | X     | $S_0$ |

Output

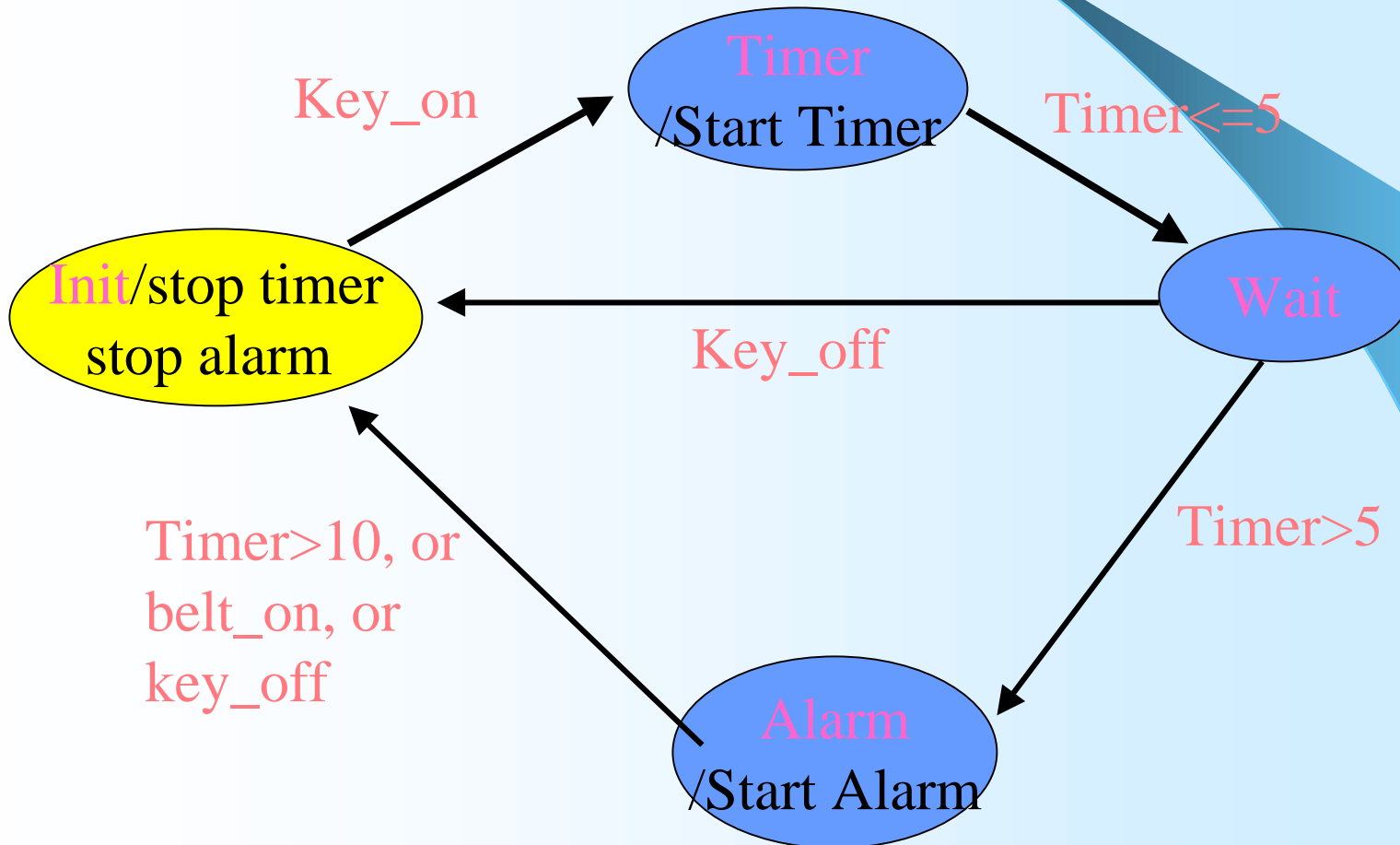
|       | $S_0$ | $S_1$ | $S_2$ |
|-------|-------|-------|-------|
| $i_0$ | $O_0$ | -     | -     |
| $i_1$ | -     | $O_1$ | $O_1$ |
| $i_2$ | -     | -     | $O_1$ |
| $i_3$ | -     | -     | -     |
| $i_4$ | -     | -     | $O_1$ |

X: don't care

# Mealy Machine and Moore Machine

- Mealy Machine
  - The output is a function of both the current state and the input
- Moore Machine
  - The output is only a function of the current state

# Transition Graph For Moore Machine



# Mealy/Moore Machine

- An FSM can be realized either by Mealy or Moore machine
- Mealy machine may use less flip-flops and output signals are immediately after the transition
- Moore machine may use more flip-flops and output signals valid except during the transition

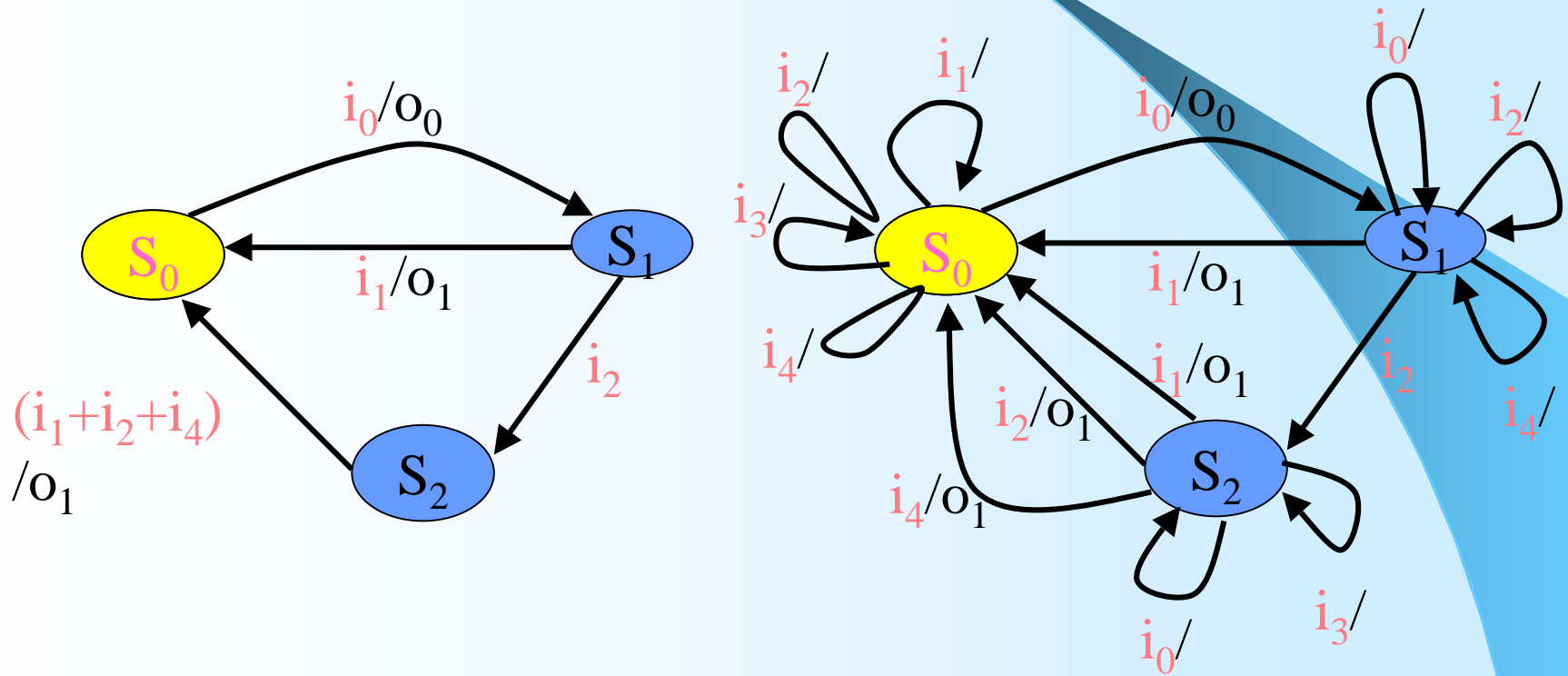
# Nondeterministic FSM

- Deterministic FSM
  - Given a state and input, there is exactly one next state
- Nondeterministic FSM (NFSM)
  - Given a state and input, there maybe more than one next state, or a state can transform from one state to anther without any input, or for some given input there no next state at all
- For any NFSM, there is always one equivalent FSM

# Nondeterministic FSM

- For unknown/unspecified behavior
- Less states, more compact
- Useful for
  - Optimization
  - Verification
- Can be refined
- For any NFSM, there is always one equivalent DFSM

# NFSM and FSM

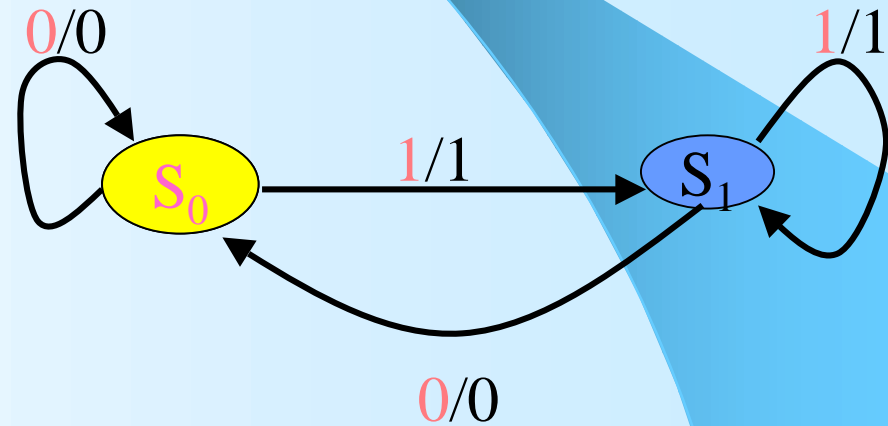
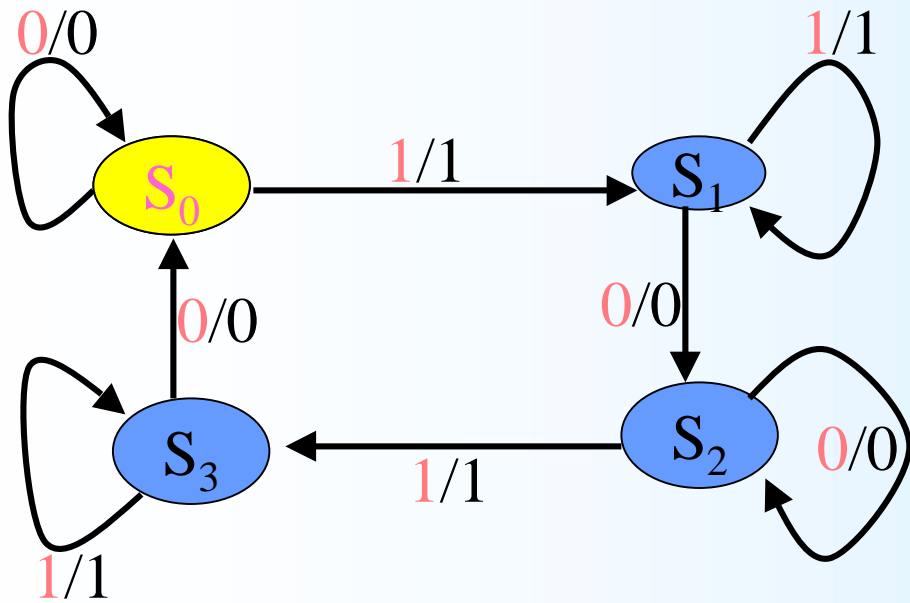




# Equivalence

- Two FSMs are equivalent iff for any given input sequence, identical output sequences are produced

# Equivalence



# Minimization

- What

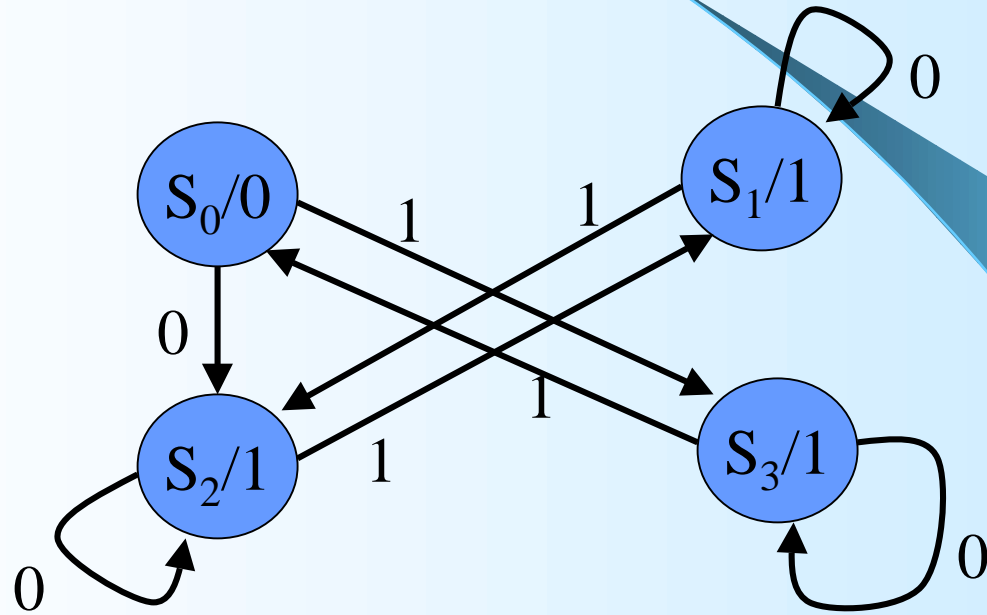
- Given an FSM, find the equivalent FSM with a minimum number of states
  - Two states  $s_1$  and  $s_2$  in an FSM are equivalent iff each input sequence beginning from  $s_1$  yields an output sequence identical to that obtained by starting from  $s_2$

- How

# Minimization(Moore Machine)

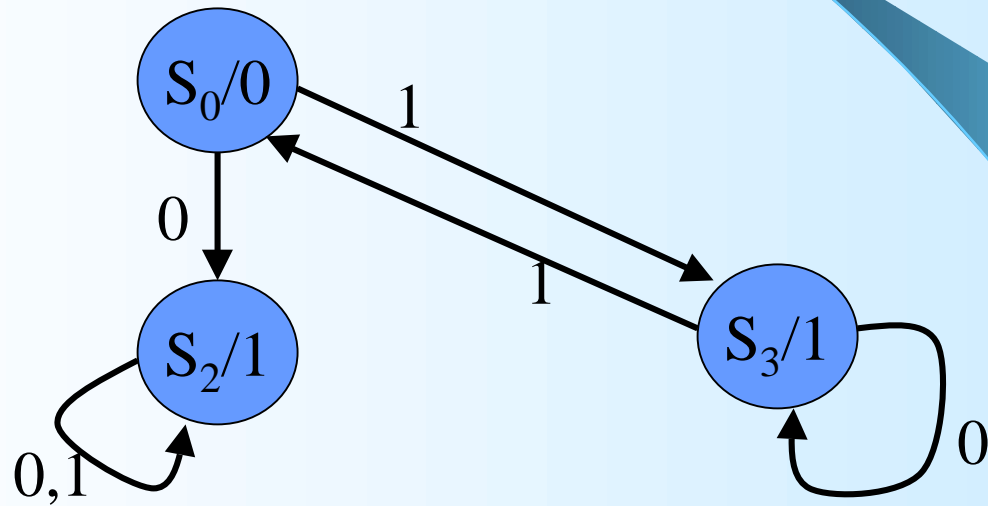
*For each pair of the states (si,sj)*  
*If si and sj have different output*  
*Mark si and si as not equivalent*  
*End for*  
*Do*  
*for each unmarked pair*  
*for each input, si and sj are transferred to states which*  
*are not equivalent*  
*Mark si and sj as not equivalent*  
*end for*  
*end for*  
*until no mark is possible*  
*Unmarked pairs are equivalent*

# Minimization



$(s_0, s_1)$   $(s_0, s_2)$   $(s_0, s_3)$   $(s_1, s_2)$   $(s_1, s_3)$   $(s_2, s_3)$

# Minimization



(s0, s1) (s0,s2) (s0,s3) (s1,s2) (s1,s3) (s2,s3)