

Transformation-Based Optimization

Basis for a series of optimizer frameworks:
Exodus, Volcano, Cascades, Columbia,
Colombia

Optimize an expression: look for all equivalent plans

- requires optimizing sub-expressions
- "memoize" results for later reuse

Produce logically equivalent expressions via rewrite

(then map them to physical plans and estimate their costs)

Based on Algebraic Equivalences

Example:

$$\sigma_p(u \cap v) \equiv \sigma_p(u) \cap \sigma_p(v)$$

Need to know such equivalences are correct

Usually easiest to work from set-theoretic definition

$$\sigma_p(r) = \{t \mid t \in r \text{ and } P(t)\}$$

$$r \cap s = \{t \mid t \in r \text{ and } t \in s\}$$

Expand Left Side

$$\begin{aligned} \sigma_P(u \cap v) &= \\ \sigma_P(\{t1 \mid t1 \in u \text{ and } t1 \in v\}) &= \\ \{t2 \mid t2 \in \{t1 \mid t1 \in u \text{ and } t1 \in v\} \text{ and } & \\ P(t2)\} &= \\ \{t2 \mid (t2 \in u \text{ and } t2 \in v) \text{ and } P(t2)\} & \end{aligned}$$

↳ ist true

Expand Right Side

$$\begin{aligned} \sigma_P(u) \cap \sigma_P(v) &= \\ \{t1 \mid t1 \in u \text{ and } P(t1)\} \cap \{t2 \mid t2 \in v & \\ \text{and } P(t2)\} &= \\ \{t3 \mid t3 \in \{t1 \mid t1 \in u \text{ and } P(t1)\} \text{ and } & \\ t3 \in \{t2 \mid t2 \in v \text{ and } P(t2)\}\} &= \\ \{t3 \mid (t3 \in u \text{ and } P(t3)) \text{ and } & \\ (t3 \in v \text{ and } P(t3))\} &= \end{aligned}$$

Continuing ...

$$\{t3 \mid (t3 \in u \text{ and } P(t3)) \text{ and } (t3 \in v \text{ and } P(t3))\} = \quad [\text{simplify}]$$

$$\{t3 \mid t3 \in u \text{ and } P(t3) \text{ and } t3 \in v\} = \quad [\text{rearrange}]$$

$$\{t3 \mid (t3 \in u \text{ and } t3 \in v) \text{ and } P(t3)\} = \quad [\text{change variable}]$$

$$\{t2 \mid (t2 \in u \text{ and } t2 \in v) \text{ and } P(t2)\}$$

et voila!

Equivalences Can Have Side Conditions

An equivalence may hold only under certain conditions

$$\pi_X(\sigma_P(r)) \equiv \sigma_P(\pi_X(r))$$

if mentioned(P) \subseteq X

mentioned(P) = set of attributes used in P

$$\pi_{ABC}(\sigma_{A>B}(r)) \equiv \sigma_{A>B}(\pi_{ABC}(r)) \quad \text{OK}$$

$$\pi_{BC}(\sigma_{A>B}(r)) \equiv \sigma_{A>B}(\pi_{BC}(r)) \quad \text{not OK}$$

Rewriting Group-by

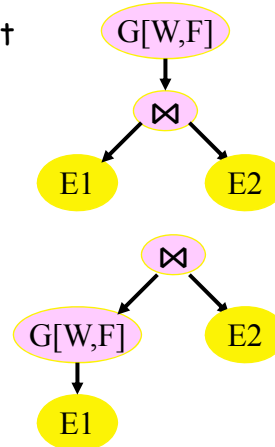
Initial translation usually puts group-by above join

May be useful to move past other operators

Move down: makes one of the join inputs smaller

Choose based on cost

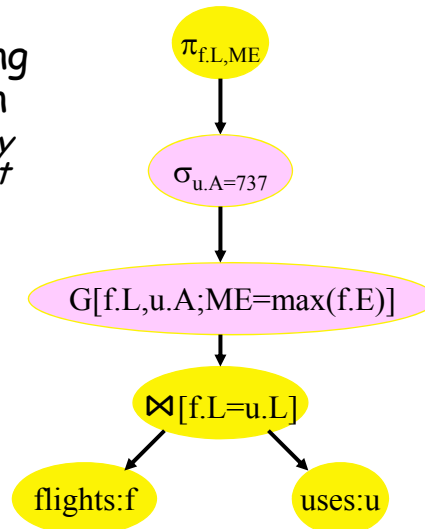
Move up: might make input to group-by smaller



Select and Group-by

Consider the following starting expression

Find max # of economy seats for flights that use a 737



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Do Selection Early

Is this okay?
 Yes - the select applies to the grouping attributes, which are constant for a group
 Want groups to pass in an all-or-nothing manner

The diagram shows a query plan starting with a projection operator $\pi_{f.L,ME}$ on top. Below it is a grouping operator $G[f.L,u.A;ME=\max(f.E)]$. This is followed by a selection operator $\sigma_{u.A=737}$. Below the selection is a join operator $\bowtie[f.L=u.L]$, which connects to two tables: 'flights:f' and 'uses:u'. A blue arrow points from the selection operator to the join operator with the text 'could push further'.

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I'm Confused!

Paper* says you can move a select if the attributes it uses are functionally determined by the grouping attributes
 Assume $L \rightarrow A$ on uses
 Doesn't seem right to me: $u.A$ isn't visible above the group-by

The diagram shows a query plan starting with a grouping operator $G[u.L;ME=\max(f.E)]$. Below it is a selection operator $\sigma_{u.A=737}$. A double-headed blue arrow indicates a potential swap between the selection and the grouping operator. Below the arrow is another selection operator $\sigma_{u.A=737}$, followed by the grouping operator $G[u.L;ME=\max(f.E)]$. A blue arrow points from the text 'Doesn't seem right to me...' to the double-headed arrow.

* Orthogonal Optimization of Subqueries and Aggregation

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Add an Attribute?

If we lift a selection above a Group-by, add the attribute to the grouping condition?

- Won't create more groups if $L \rightarrow A$.
- Will need to project away A at some point.

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Group-by and Join

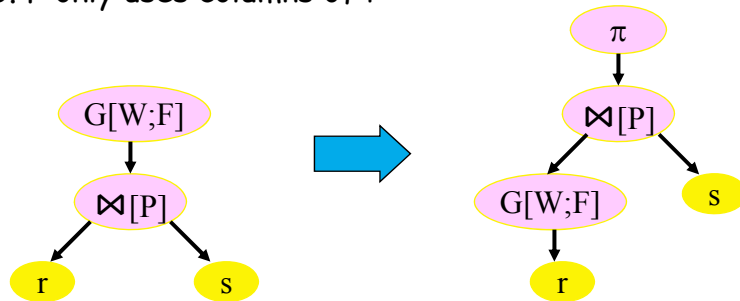
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Side Conditions

Have $r(R)$, $s(S)$, P mentions $X \subseteq R$

Group-by through join okay if

1. X is in W : join columns of R included in grouping columns
2. W contains a key of s
3. F only uses columns of r



Why?

1. Can't lose the join columns
2. Need to have only one s tuple per row of $G[W;F]$
Can relax for max and min
3. Only have r available

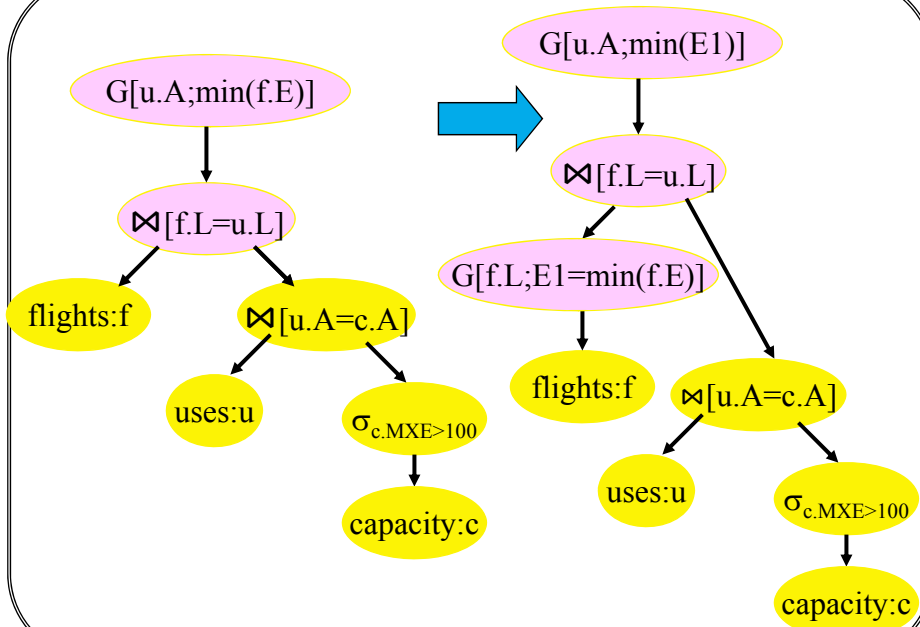
Transforming from right to left,
conditions 1 and 3 come for free
Given that the initial expression is well formed

Splitting a Group-by

Can do partial grouping (local group-by), then a join, then rest of grouping

```

select u.A, min(f.E)
from flight as f, uses as u,
     capacity as c
where f.L = u.L and u.A = c.A
     and c.MXE > 100
group-by u.A
    
```



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Min of Mins = Min of all

<u>u.A</u>	<u>f.L</u>	<u>f.E</u>	<u>u.A</u>	<u>f.L</u>	<u>E1</u>	<u>u.A</u>	<u>min(E1)</u>
737	115	65	737	115	65	737	65
737	115	70					
737	115	75					
737	80	95	737	80	75		
737	80	75					

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Works for Sum, Too

<u>u.A</u>	<u>f.L</u>	<u>f.E</u>	<u>u.A</u>	<u>f.L</u>	<u>E1</u>	<u>u.A</u>	<u>sum(E1)</u>
737	115	65	737	115	210	737	380
737	115	70					
737	115	75					
737	80	95	737	80	170		
737	80	75					

$E1 = \text{sum}(f.E)$

sum of counts = count of all

What about count? Does count of counts = count of all?
 What about average??

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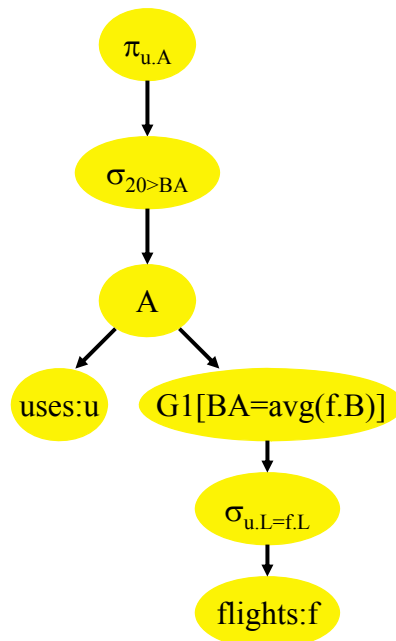
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Removing Sub-queries

We introduced "apply" to translate subqueries

```

select u.A
from uses as u
where 20 > [BA]
      (select avg(f.B) [as BA]
       from flights as f
       where f.L = u.L)
    
```



Apply Options

Can have a direct (for-loop) implementation of apply as a physical operator

But it is also possible to remove apply via rewriting

Strategy:

Use push-down transforms that move apply towards bottom of expression tree

Until its right input no longer depends on the left input

Then, convert apply to a join

Example: Apply through Scalar Group-by

$$r \text{ A } G1[F](e(x)) \rightarrow G[r.*, F](r \text{ A}^{LOJ} e(x))$$

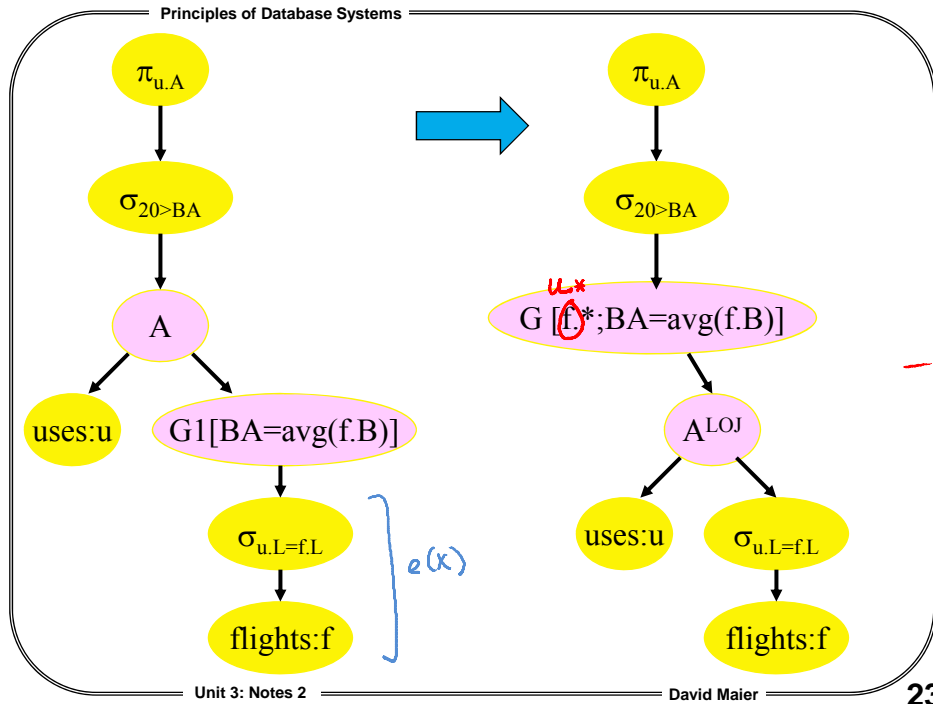
A^{LOJ} keeps every tuple t in r , even if $e(t)$ is empty

need it because $G1$ always produces a row

$$r \text{ A}^{LOJ} e(x) =$$

$$r \text{ LOJ } (DE(r) \text{ A } e(x))$$

collect tuples $t \in r$ where $e(t) = \emptyset$



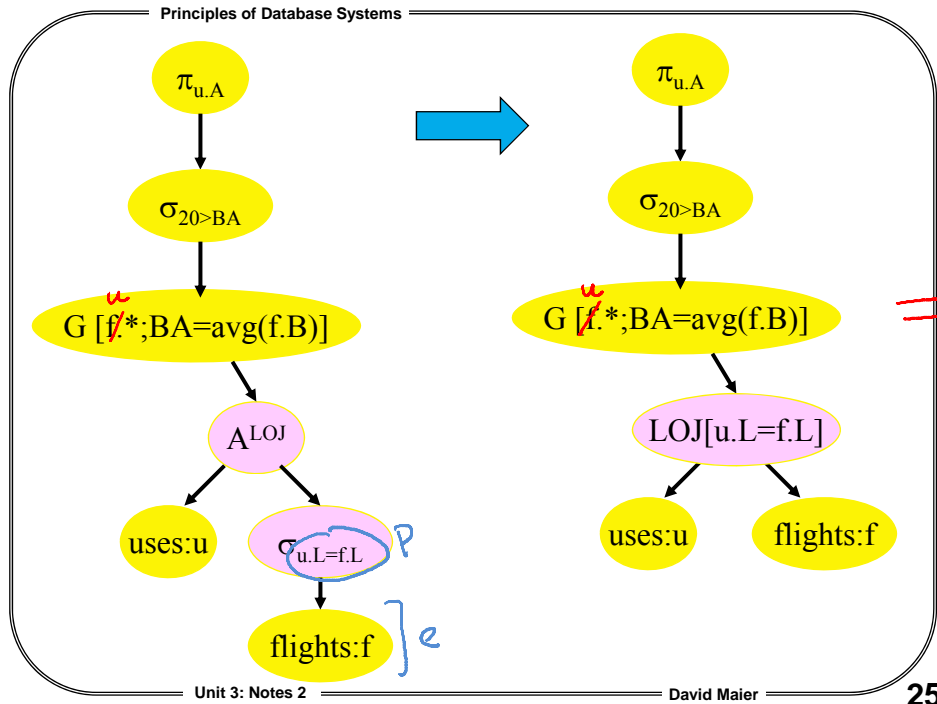
Example: Apply-to-Join Conversion

$$r A^{LOJ} \sigma_P(e) \rightarrow r LOJ[P] e$$

if e does not depend on r
(okay for P to depend on r)

Can think of apply as a "catalyst": used in transforming algebra, but doesn't show up in the end product

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What Was I Lying About?
Subqueries of optimal queries are optimal

Not all query plans are created equal
 They differ in *physical properties*

- sort order
- distribution
- compression

Some physical operators need inputs with certain properties in order to work correctly

Sort-merge join: needs inputs sorted on attributes in an equality join condition

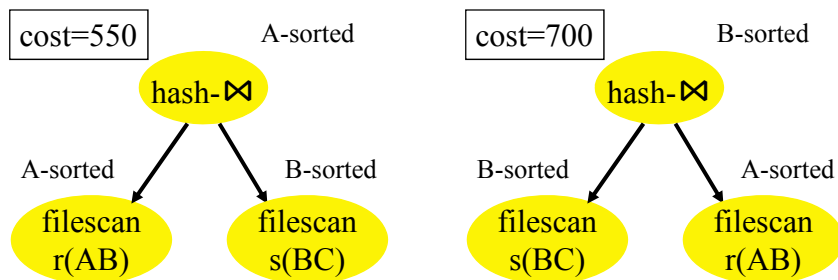
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Physical Properties Can Affect Cost

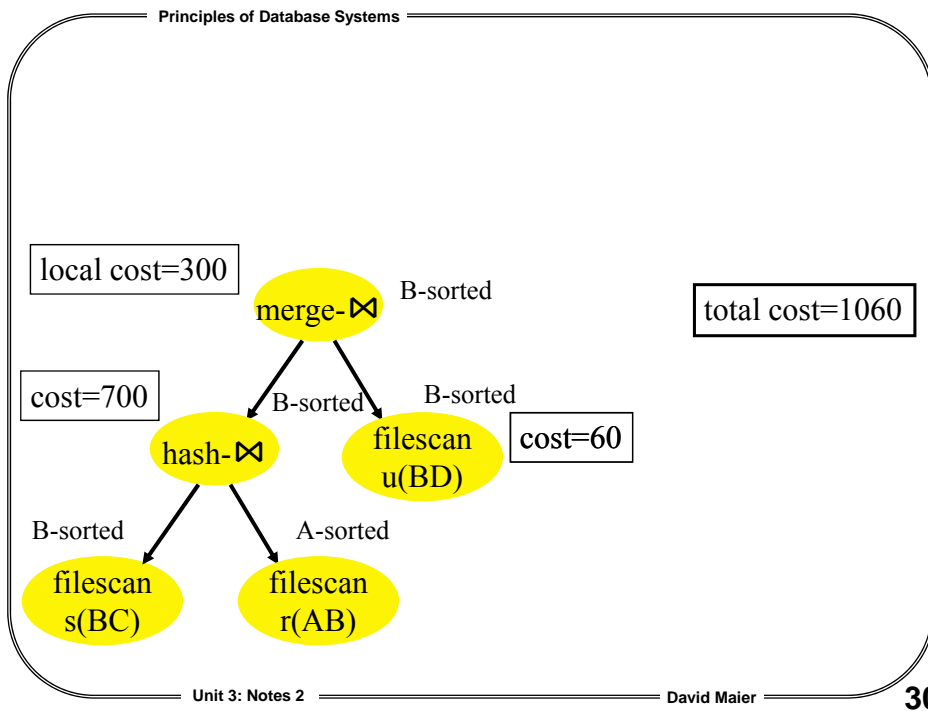
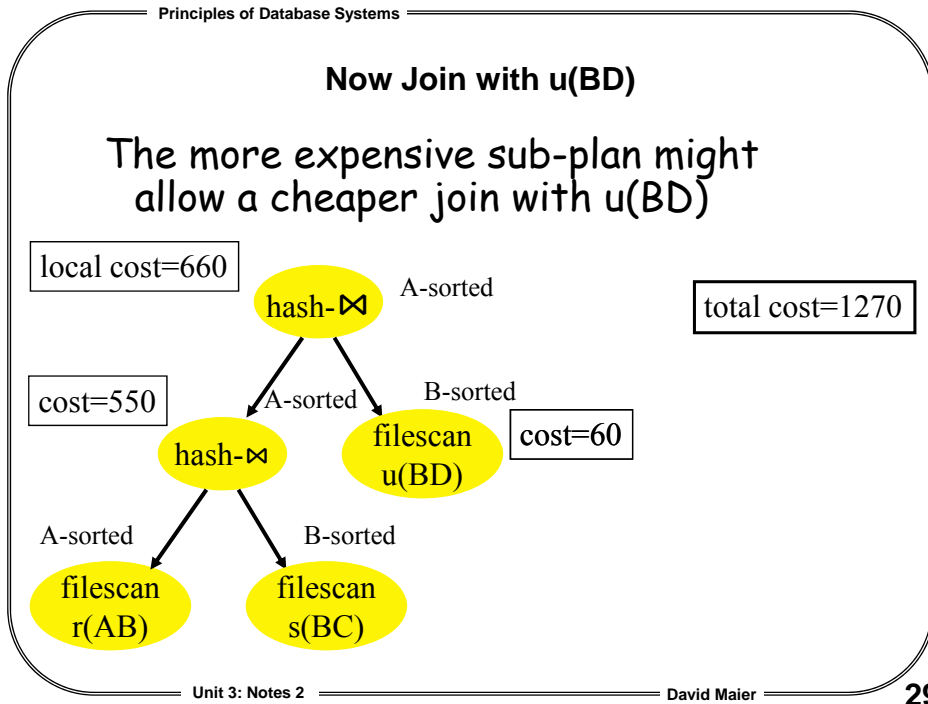
Union, project might be faster on compressed inputs
Fewer bytes to deal with

Physical Properties and Optimization

Because of differing physical properties, cheapest sub-plan might not be best choice in overall plan.
consider join of r, s and u



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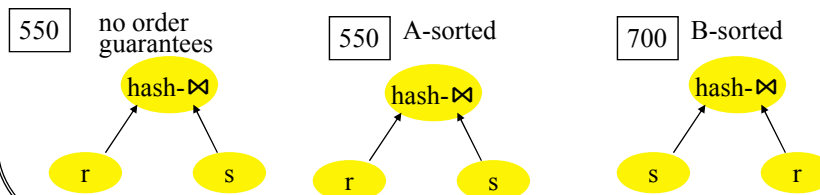


So, Do We Forget the Optimality Principle?

No - it still holds, relative to plans with the same physical properties.

So - need to keep best plan for each physical property

For example, dynamic programming approach keeps the best plan for each collection of physical properties

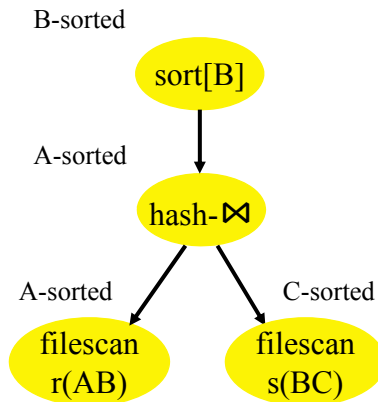


Do We Need Plans for Every Possible Property?

Probably not. For example, System R only keeps best plans for each *interesting order* (sort orders that might be useful in rest of query)

- attributes in a join predicate
- grouping attributes in an aggregate
- order-by clause in original query

Can Enforce Properties



Even if a sub-plan doesn't produce a particular physical property, usually possible to modify output to do so.

Enforcer: physical operator that enforces a physical property, but is a "no-op" at the logical level
Usually incurs some cost

Mapping Logical to Physical

Not a one-to-one correspondence generally between logical and physical operators

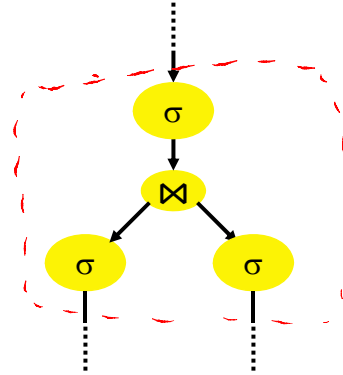
- Can have several physical ops that correspond to a given logical op

join: nested-loops join, index-nested-loops join, hash join, merge join

- One physical op might implement several logical ops

Have flags on hash join for regular join, outer-join, semi-join

- One physical op might translate a tree of logical ops



hash-join(join-pred, pre-left,
pre-right, post-pred)

Cost Estimation

Can view operators under alternative interpretations

Normal view

- $\sigma_{A>B}$: relation \rightarrow relation

Alternative interpretations

Properties

- $\sigma_{A>B}$: order \rightarrow order, FDs \rightarrow FDs

Stats

- $\sigma_{A>B}$: column histogram \rightarrow column histogram

Costs

- $\sigma_{A>B}$: CPU secs \rightarrow CPU secs

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Also for Selection and Join Predicates

Interpret predicates as a *selectivity*.
% of tuples (or tuple pairs) expected
to satisfy the condition

System R

$sel(A=5) = 10\%$

$sel(A<5) = 50\%$

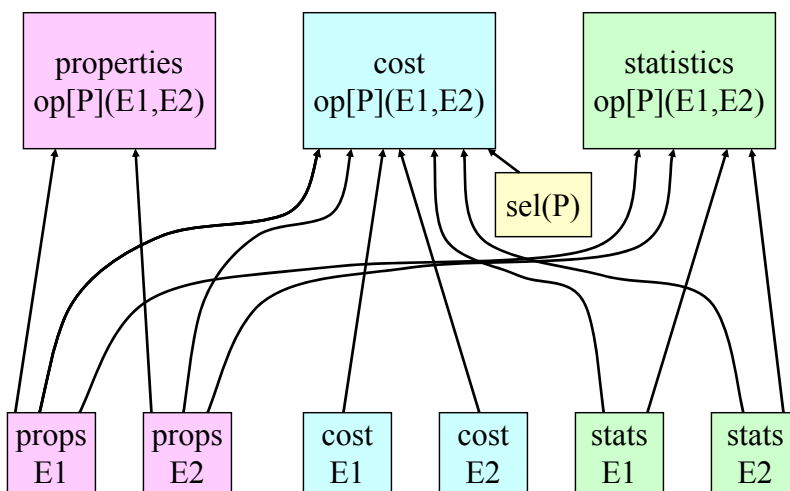
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These Interpretations Make Use of Each Other

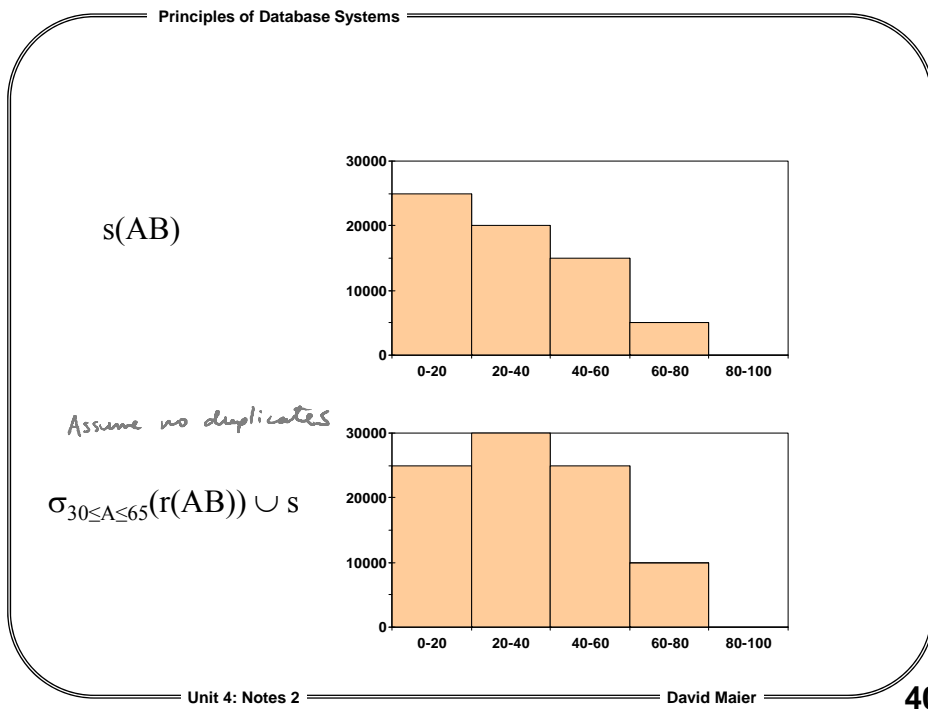
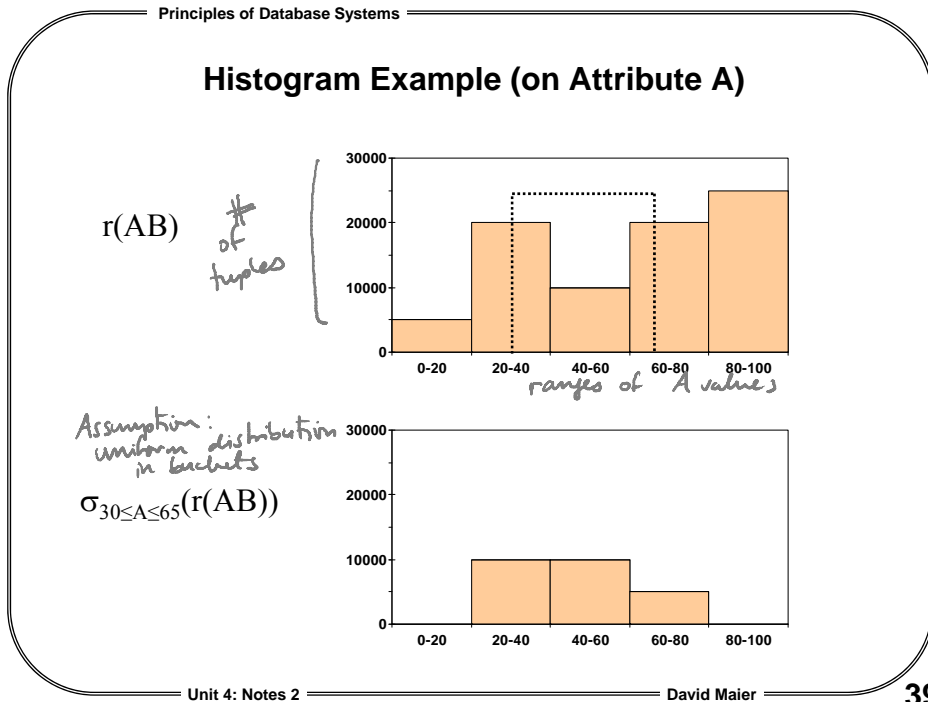


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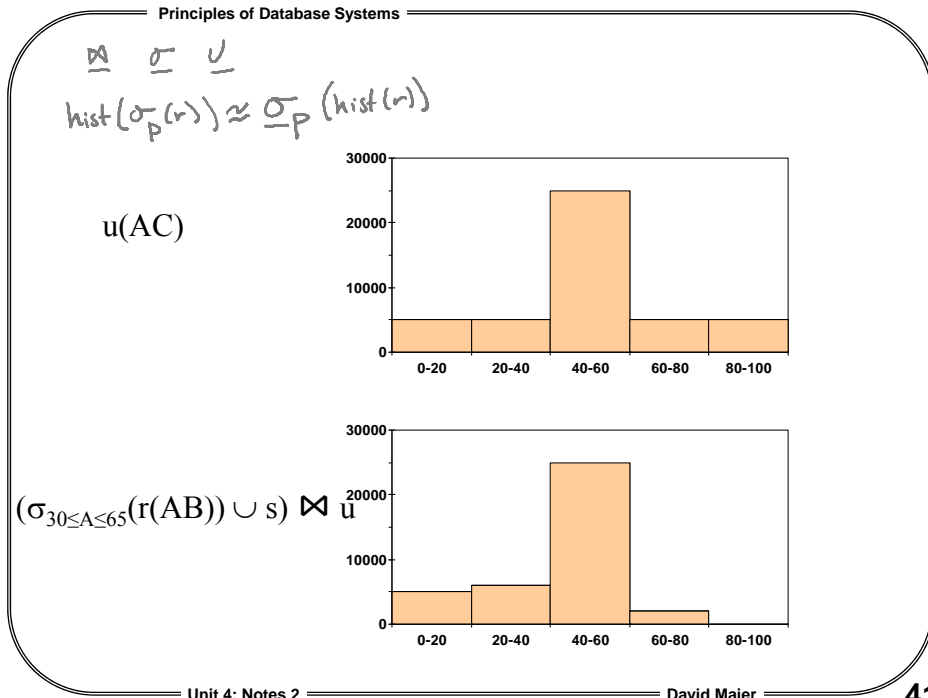
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