

## Example of Naive 1

```


Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

$f_0 = \emptyset$

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## Example of Naive 2

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

$f_1 = I_p(f_0) =$   
*Local0(r, node1).*  
*Local0(s, node1).*  
*Local0(q, node1).*  
*Local0(u, node2).*  
*Op1(s1, r, node1).*  
*Op1(s2, u, node2).*  
*Op2(j1, s1, u1, node1).*  
*Op2(u1, s, q, node1).*  
*Op2(j2, j1, s2, node1).*

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## Example of Naïve 3

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

```

f2 = Ip(f1) =
Local(r, node1, 1).
Local(s, node1, 1).
Local(q, node1, 1).
Local(u, node2, 1).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```



## Example of Naïve 4

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

```

f3 = Ip(f2) =
Local(r, node1, 1).
Local(s, node1, 1).
Local(q, node1, 1).
Local(u, node2, 1).
Local(s1, node1, 2).
Local(u1, node1, 3).
Local(s2, node2, 2).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```



## Example of Naïve 5

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

f4 = Ip(f3) =
Local(r, node1, 1).
Local(s, node1, 1).
Local(q, node1, 1).
Local(u, node2, 1).
Local(s1, node1, 2).
Local(u1, node1, 3).
Local(j1, node1, 6).
Local(s2, node2, 2).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

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## Example of Naïve 6

```


Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

f5 = Ip(f4) = f4
Local(r, node1, 1).
Local(s, node1, 1).
Local(q, node1, 1).
Local(u, node2, 1).
Local(s1, node1, 2).
Local(u1, node1, 3).
Local(j1, node1, 6).
Local(s2, node2, 2).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

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## Problem with Computed Predicates

```


Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
                    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
                    Local(In1, Loc, M), Local(In2, Loc, N),
                    Sum(M, N, J), Sum(J, 1, K).

Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).

f1 = Ip(∅)
Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).
-----
f2 = Ip(f1)
Local(r, node1, 1).
Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).
-----
f3 = Ip(f2)
Local(r, node1, 1).
Local(s1, node1, 3).
Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).

```

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## Some Refinements of Naive

- Start with  $f = \text{facts}(P)$
- Use dependency graph
  - Nodes are the different predicates of  $P$
  - Have an edge  $p \rightarrow q$  if there is a rule of the form
 
$$q(\dots) \text{ :- } \dots, p(), \dots .$$

The paper uses a more detailed version of the dependency graph than the book, where there are nodes for rules as well.

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## Example Dependency Graph

```

Needed(D, M, C) :- MajAll(M, C), Offered(D, M).
Needed(D, M, C) :- DegAll(D, C), Offered(D, M).

MajAll(M, C) :- MajReq(M, C).
MajAll(M, C) :- MajAll(M, C1), Prereq(C, C1).

DegAll(D, C) :- DegReq(D, C).
DegAll(D, C) :- DegAll(D, C1), Prereq(C, C1).


```

Needed

MajAll          DegAll          Offered

MajReq          Prereq          DegReq

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## Evaluate in Order of Dependency Graph

- **Compute** DegAll from DegReq and Prereq
- **Compute** MajorAll from MajReq and Prereq
- **Compute** Needed from MajAll, DegAll, Offered

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## Semi-Naïve Evaluation

Naïve evaluation discovers the same fact over and over.

E.g., derive `local(r, node1, 1)` repeatedly

**Idea:** When evaluating a rule, make sure we use at least one newly discovered fact.

Otherwise, we derive a fact we've already seen



## Example

Consider the rule

```
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
  Local(In1, Loc, M), Local(In2, Loc, N),
  Sum(M, N, J), Sum(J, 1, K).
```

When computing  $f4 = I_p(f3)$ .

If we use facts `Op2(u1, s, q, node1)`,  
`Local(s, node1, 1)`, `Local(q, node1, 1)`, we  
 derive `Local(u1, node1, 3)`, which we already  
 have

If we use `Op2(j1, s1, u1, node1)`,  
`Local(s1, node1, 2)`, `Local(u1, node1, 3)`, we  
 get



## Implementing Semi-Naive

Notice after  $f_1$ , only new facts are `local` facts

Suppose at each step in Naïve we determine

`NewLocal(Id, Loc, K)` if `Local(Id, Loc, K)` in  $f - f_{old}$ .

We can work with a modified set of rules

```
Local(Id, Loc, K) :- Op1(Id, In, Loc),
  NewLocal(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
  NewLocal(In1, Loc, M), Local(In2, Loc, N),
  Sum(M, N, J), Sum(J, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
  Local(In1, Loc, M), NewLocal(In2, Loc, N),
  Sum(M, N, J), Sum(J, 1, K).
```



## One Other Modification

If  $P'$  is the modified program, then instead of

$$f = I_p(f)$$

we need

$$f = I_{p'}(f) \cup f.$$

Note: Semi-Naïve doesn't guarantee we won't generate a fact at multiple steps, if there are distinct "proofs" of it

```
MajReq(eng, wr310). MajReq(eng, wr408).
Prereq(wr309, wr310).
Prereq(wr401, wr408).
Prereq(wr309, wr401).
```



## Directed Methods

Suppose we have the query

```
:- Local(I, node2, K).
```

We could push the constant into the rules:

```
Local(Id, node2, 1) :- Local0(Id, node2).
```

```
Local(Id, node2, K) :- Op1(Id, In, node2),
  Local(In, node2, M), Sum(M, 1, K).
```

```
Local(Id, node2, K) :- Op2(Id, In1, In2, node2),
  Local(In1, node2, M), Local(In2, node2, N),
  Sum(M, N, J), Sum(J, 1, K).
```



## Top-Down Methods

Sometimes you can propagate constants statically (at "compile" time), before any computation.

However, not all constants are available at compile time.



## Example of Constant Propagation

### Consider query

```
:- (ba, eng, C).
```

### for our program

```
Needed(D, M, C) :- MajAll(M, C), Offered(D, M).
```

```
Needed(D, M, C) :- DegAll(D, C), Offered(D, M).
```

```
MajAll(M, C) :- MajReq(M, C).
```

```
MajAll(M, C) :- MajAll(M, C1), Prereq(C, C1).
```

```
DegAll(D, C) :- DegReq(D, C).
```

```
DegAll(D, C) :- DegAll(D, C1), Prereq(C, C1).
```



## Can Do Some Constant Propagation Right Away

### Modified program

```
Needed(ba, eng, C) :- MajAll(eng, C), Offered(ba, eng).
```

```
Needed(ba, eng, C) :- DegAll(ba, C), Offered(ba, eng).
```

```
MajAll(eng, C) :- MajReq(eng, C).
```

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```

```
DegAll(ba, C) :- DegReq(ba, C).
```

```
DegAll(ba, C) :- DegAll(ba, C1), Prereq(C, C1).
```



## Consider Evaluation of `MajAll`

Suppose we have the facts

```
MajReq(eng, wr310).
```

```
MajReq(eng, lit481).
```

We'll derive (by the first rule)

```
MajAll(eng, wr310).
```

```
MajAll(eng, lit481).
```

Specializing the second rule:

```
MajAll(eng, C) :- MajAll(eng, wr310), Prereq(C, wr310).
```

```
MajAll(eng, C) :- MajAll(eng, lit481),  
                  Prereq(C, lit481).
```



## New Sub-queries

Pose the new queries

```
:- Prereq(C, wr310).
```

```
:- Prereq(C, lit481).
```

which can be used to instantiate the rules  
together with the `MajAll` facts.



## The Query-Sub-Query (QSQ) Method

Combine queries with the same *query form* (pattern of free and bound positions) into a *generalized query*:

```
:- Prereq(C, {wr310, lit481}).
```

Use generalized query:

- Evaluate against extensional DB (if appropriate)
- Plug results into rules for new facts
- Use new facts to generate new sub-queries
- Combine sub-queries with existing generalized query, if possible




## Example

Suppose we have the facts:

```
Prereq(wr301, wr310).
Prereq(wr301, wr418).
Prereq(wr200, wr301).
Prereq(lit480, lit481).
Prereq(lit217, lit218).
```

Which match the generalized query?

```
:- Prereq(C, {wr310, lit481}).
```



## Example 2

**Plug**

```
Prereq(wr301, wr310).
Prereq(lit480, lit481).
```

**plus**

```
MajAll(eng, wr310).
MajAll(eng, lit481).
```


**into**

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```

**to get**

```
MajAll(eng, wr301).
MajAll(eng, lit480).
```

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## Example 3

**Use**

```
MajAll(eng, wr301).
MajAll(eng, lit480).
```

**with**

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```

**to get new sub-queries**

```
:- Prereq(C, wr301).
:- Prereq(C, lit480).
```

**Expand generalized query**

```
:- Prereq(C, {wr310, lit481, wr301, lit480}).
```

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## Example 4

### Evaluate generalized query again

```
Prereq(wr301, wr310).
Prereq(wr301, wr418).
Prereq(wr200, wr301).
Prereq(lit480, lit481).
Prereq(lit217, lit218).
```

### Which match the new generalized query?

```
:- Prereq(C, {wr310, lit481, wr301, lit480})
```



## Example 5

### Plug

```
Prereq(wr301, wr310).
Prereq(wr201, wr301).
Prereq(lit480, lit481).
```

### plus

```
MajAll(eng, wr310).
MajAll(eng, lit481).
MajAll(eng, wr301).
MajAll(eng, lit480).
```

### into

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```



## Example 6

We get

```
MajAll(eng, wr301).
MajAll(eng, lit480).
MajAll(eng, wr200).
```

Plus we expand the generalized query to

```
:- Prereq(C,
        {wr310, lit481, wr301, lit480, wr200}).
```



## Example 7

Evaluate generalized query again

```
Prereq(wr301, wr310).
Prereq(wr301, wr418).
Prereq(wr200, wr301).
Prereq(lit480, lit481).
Prereq(lit217, lit218).
```

Which match the new generalized query?

```
:- Prereq(C,
        {wr310, lit481, wr301, lit480, wr200}).
```



## Example 8

The answers we get are the same for the last version of the generalized query:

```
Prereq(wr301, wr310).  
Prereq(wr201, wr301).  
Prereq(lit480, lit481).
```

So we will get no new facts nor new sub-queries.