

# HW1 KEY - CS 589 - FALL 2008

(2pts) 1.1.  $N_2 \leq N \leq N_1 + N_2$ ,  $R_1 \neq R_2$  must be union compat.  
 $\underbrace{\hspace{10em}}_{\text{no tuples in common}} \quad \underbrace{\hspace{10em}}_{R_1 \subseteq R_2}$

(2pts) 1.2.  $\emptyset \leq N \leq N_1$ , and  $R_1 \neq R_2$  union compat  
 $\underbrace{\hspace{10em}}_{R_1 \subseteq R_2} \quad \underbrace{\hspace{10em}}_{R_1 \not\subseteq R_2 \text{ (i.e. } R_1 \cap R_2 = \emptyset)}$

(2pts) 1.3.  $\emptyset \leq N \leq N_1$ , and  $R_1 \neq R_2$  union compat.  
 $\underbrace{\hspace{10em}}_{R_1 \cap R_2 = \emptyset} \quad \underbrace{\hspace{10em}}_{R_1 \subseteq R_2}$

(2pts) 1.4.  $N \equiv N_1 \times N_2$ , and for completeness,  $\mathbb{R} \text{ schema}(R_1) \cap \text{schema}(R_2) = \emptyset$

(2pts) 1.5.  $\emptyset \leq N \leq N_1$ , and  $a \in \text{schema}(R_1)$ , and  $a$ 's type must be comparable to "5"  
 $\underbrace{\hspace{10em}}_{\text{all quality}} \quad \underbrace{\hspace{10em}}_{\text{none quality}}$

(2pts) 1.6.  $1 \leq N \leq N_1$ , and  $a \in \text{schema}(R_1)$   
 $\underbrace{\hspace{10em}}_{a \text{ is a candidate key for } R_1} \quad \underbrace{\hspace{10em}}_{a \text{ is the same in all } t \in R_1}$

(2pts) 1.7.  $N \equiv \emptyset$ , and  $\text{schema}(R_1) \supseteq \text{schema}(R_2)$   
 $\underbrace{\hspace{10em}}_{\text{since } N_2 > N_1, R_2 \text{ can never have enough tuples such that it "covers" the rows in } R_1}$

(5pts) 2.1. Names of suppliers of parts in the catalog that cost < \$100.

Explanation:

$$\begin{aligned} & \pi_{sname} (\pi_{sid} (\sigma_{color=red} Parts) \bowtie (\sigma_{cost < 100} Catalog) \bowtie Suppliers) \\ &= \pi_{sname} (\emptyset \bowtie (\sigma_{cost < 100} Catalog) \bowtie Suppliers) \\ &= \pi_{sname} ((\sigma_{cost < 100} Catalog) \bowtie Suppliers) \end{aligned}$$

(5pts) 2.2. Names of suppliers who supply both green parts under \$100 and red parts under \$100

(5pts) 3.1.  $\{ T \mid \exists C \in Catalog \exists S \in Suppliers$

$$(C.cost < 100 \wedge C.sid = S.sid \wedge T.sname = S.sname) \}$$

(5pts) 3.2  $\{ T \mid \exists P_1, P_2 \in Parts \exists C_1, C_2 \in Catalog \exists S_1, S_2 \in Suppliers$

$$\begin{aligned} & ( \\ & \quad ((P_1.color = red) \wedge (C_1.cost < 100) \wedge (P_1.pid = C_1.pid) \wedge (C_1.sid = S_1.sid)) \\ & \quad \wedge \\ & \quad ((P_2.color = green) \wedge (C_2.cost < 100) \wedge (P_2.pid = C_2.pid) \wedge (C_2.sid = S_2.sid)) \\ & \quad \wedge \\ & \quad (T.sname = S_1.sname) \\ & \quad \wedge \\ & \quad (T.sname = S_2.sname) \\ & \quad ) \} \end{aligned}$$

(5pts) 4.1.  $\{x_1: SNAME \mid \exists x_2: SID (\exists x_6: sid (\exists x_3: ADDRESS$   
 $(\exists x_4: PID (\exists x_5: COST$

$( (x_5 < 100) \wedge (x_2 = x_6) \wedge$

$SUPPLIERS(x_2, x_1, x_3) \wedge$

$CATALOG(x_6, x_4, x_5) ) ) ) ) ) \} (d)$

(5pts) 4.2  $\{x: SNAME \mid \exists p_1: pid (\exists p_n: pname (\exists c_1: cost (\exists a_1: address$   
 $(\exists s: sid (\exists p_2: pid (\exists p_{n2}: pname (\exists c_2: cost$

$\exists a_2: address$

$\wedge Parts(p_1, n_1, 'red') \wedge Catalog(s, p_1, c_1) \wedge c_1 < 100$   
 $\wedge Suppliers(s, x, a_1)$

$\wedge Parts(p_2, n_2, 'green') \wedge Catalog(s, p_2, c_2)$   
 $\wedge c_2 < 100 \wedge Suppliers(s, x, a_2) ) ) ) ) ) ) \} (d)$

NOTE: shorthanding with constants, e.g.  $Parts(p_2, n_2, 'green')$ ,  
 might be considered cheating a bit, but it's  
 shorter.

REF: THIS answer adapted from a student paper.  
 Thanks! to J.N.

5. FOLLOWING THE EXAMPLE IN THE BOOK.

(2 pts) a) schema (Supply chain) = {Supplier, Part}  
schema (Project Repts) = {Part, Project}

Supply Chain		Project Repts	
<u>Supplier</u>	<u>Part</u>	<u>Part</u>	<u>Project</u>
Bob	Widget	widget	A
Jim	Whatsit	widget	B
Bob	Thingy	thingy	A
		thingy	B
		whatsit	A

(2 pts) b) { Bob, B }

(2 pts) c) Show the list of suppliers that supply all parts for some project.