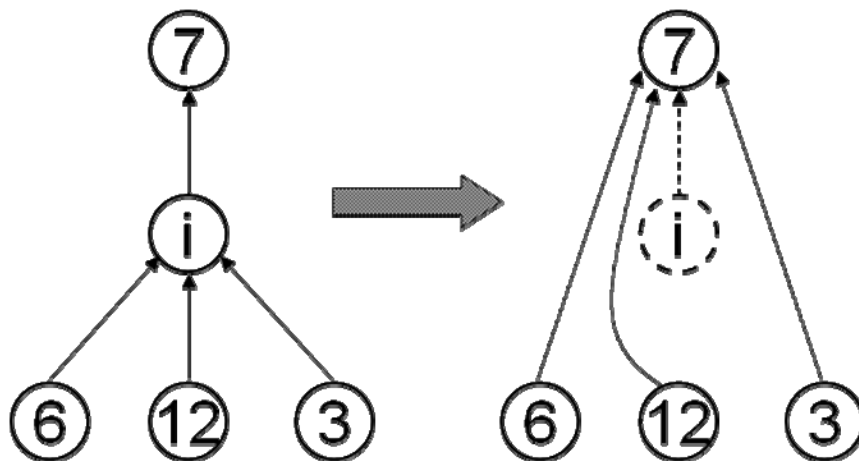


**CS 410/586: Quiz 4, 20 April 2009 Name: \_\_\_\_\_ KEY \_\_\_\_\_**

No books or notes. Work individually.

Prof. Mumble claims he has a method to support DELETE's in the tree-structured UNION-FIND algorithm in a way that allows a sequence of  $n$  operations (UNION, FIND and DELETE) in  $O(n G(n))$  time. His method is described by the following pseudo-code and figure.

```
DELETE( $i$ )
for each child  $c$  of the node for  $i$ 
    Parent( $c$ ) ← Parent( $i$ )
remove the node for  $i$ .
```



You point out two problems with Prof. Mumble's method.

Question 1 (4 points): DELETE( $i$ ) won't work correctly for all items. Why?  
*If node  $i$  happens to be the root, then the children of  $i$  won't be linked correctly.*

Question 2 (6 points): A sequence of  $n$  operations might actually take  $O(n^2)$  time. Why?  
*Observe that in the tree structure for UNION-FIND, each node knows its parent, but not its children. So to find all the children of node  $i$ , we actually have to iterate through every node, checking its parent. So DELETE is actually  $O(n)$  if there are  $n$  elements in the universe. (There are also arguments that can be made relative to a certain shape of tree, and that one node can have  $O(n)$  children that have to be updated on a DELETE.)*

Note that if you wanted to include DELETE and keep the same time complexity, you could leave the node in place, but just have a 1-bit flag to mark the node as "dead".