

CS 410/586: Quiz 1, 29 March 2011

Name: _____

Non-counting. No books or notes. Work individually.

Scoring: 2 points per question, 10 points total.

These questions are about the Freq algorithm.

Question 1: What is this general style of algorithm called?

Divide and Conquer (a kind of recursive algorithm)

Question 2: What happens if Freq is called on collection S where two different elements both have the maximum frequency?

It might seem at first like the larger of the two is always returned, because f_{high} “wins” if there is a tie, in Step 4. However, if the larger element is chosen as e , then Freq will return the smaller element.

Question 3: Suggest a method to perform Steps 2 and 3 of Freq with one pass through S (not counting the recursive calls). What data structure for S would support that method?

We don't need to go through S twice to get S_{low} and S_{high} . We can make one pass, placing each element in the appropriate subcollection based on comparison with e . Linked lists would be good for this approach. We can also compute f and the sizes of S_{low} and S_{high} during this pass.

(It's also possible to do the partitioning into S_{low} and S_{high} in place in an array, as is done in some implementations of Quicksort.)

Question 4: Suggest an improvement to Step 3 that can avoid one or both recursive calls in some cases.

If $f > |S_{low}|$, we can just use $(null, 0)$ for (e_{low}, f_{low}) instead of calling $Freq(S_{low})$, because S_{low} can't contain an element that is more frequent than e . Similarly for S_{high}

Question 5: Suppose we are fortunate and the choice of e always splits S into subsets of approximately equal size (that is $|S_{low}| \approx |S_{high}|$). What is the maximum depth of recursion Freq will have on an input of size n ?

In the case described, at each level of recursion, we will have calls on collections that are roughly half the size of the input. The number of times we can divide n in half is roughly $\log_2 n$ (also known as $\lg n$).

The following routine finds the most frequent element in a collection S of integers, plus the frequency (number of occurrences) of that element (without having to sort S). Note that S can have duplicates.

$|S|$ means the number of elements in S . For example $|\{3, 9, 8, 3\}| = 4$.

Freq(S)

0. If $|S| = 0$, return (null, 0)

1. Pick some element $e \in S$.

2. Split S into two subsets:

$$S_{low} \leftarrow \{d \in S \mid d < e\}$$

$$S_{high} \leftarrow \{d \in S \mid d > e\}$$

3. Set $f \leftarrow |S| - |S_{low}| - |S_{high}|$

Set $(e_{low}, f_{low}) \leftarrow \text{Freq}(S_{low})$

Set $(e_{high}, f_{high}) \leftarrow \text{Freq}(S_{high})$

4. If $f > f_{low}$ **and** $f > f_{high}$, **return** (e, f)

else if $f_{low} > f_{high}$, **return** (e_{low}, f_{low})

else return (e_{high}, f_{high})

Show a sequence of recursive calls for

Freq($\{4, 2, 1, 6, 6, 1, 2, 2, 8, 4\}$)