DNS - Domain Name System

TCP/IP class

outline

introduction

- naming scheme
- protocol
 - format
 - record types
 - how it works
- reverse lookup
- implementation named config files
- summary futuresJim Binkley

bibliography

RFCs

- 1034 Internet Domain Naming Philosophy (87)
- 1035 Protocol standard (87)
- updated in 1101/1183 in 1990
- Evi Nemeth's Unix Sys. Admin Handbook
- Sun Network Documentation
- BSD Bind docs
- Steven's TCP/Proto book gives good protocol explanation

Some DNS urls of interest

- ICANN Inet Corporation for Assigned Names and Numbers - www.icann.org
 - DNS administration
 - see the FAQ for recent info
- http://www.internic/net/regist.html
 - accredited list of registrars for .com/.net/.org
- http://www.isc.org/products/BIND

- where to get DNS BSD Bind software

intro

DNS - Domain Name System

- distributed database: (key, value) pairs include:
 - 1. map DNS names to ip number (ADDRESS)

» sirius.cs.pdx.edu to 131.1.2.3

- » connect(2) wants an ip number, so does sendto(2)
- 2. map ip number to DNS name (PTR)
- 3. mail routing, (MX)
 - » mail cat@cs.pdx.edu -> sirius.cs.pdx.edu
- 4. aliases (CNAME)
 - » www.cs.pdx.edu is actually sirius.cs.pdx.edu

intro

- done with distributed client/server implementation
- system is hierarchical and very scalable
- Internet-wide service, possibly the largest nameserver paradigm on the block
- DNS is of course both a protocol specification and an implementation

◆ BSD BIND/named implementation (4.9?) Jim Binkley 6

intro - why?

- computers use ip numbers, people want names, 1.2.3.4 not as good as foo.com
- all computers can't be named "fluffy" so we need a hierarchical naming system to avoid host name conflicts
- origin rooted in /etc/hosts file which has ip address to name mapping, but obviousally doesn't scale to Internet-wide

3 schemes for name lookup

- 1. put it in /etc/hosts and distribute that periodically (rdist/cron/rcp, whatever)
- 2. Sun's NIS (Network Info. System)
 - 1. not supported on all hosts
 - 2. can't go outside your admin. domain
 - » (Tek could do it, but can't manage it with Intel)
 - 3. can do more than DNS, passwd, groups too



Berkeley Internet Name Domain

parts include:

- named a name server, takes files as database including a list of names/ip pairs, DNS roots, master configuration file
- resolver client side library code for making
 DNS lookup, linked to apps like telnet, etc.,
 inside gethostbyname(3), gethostbyaddr(3)
 calls

debug tools: nslookup, dig (elsewhere)
 Jim Binkley

intro - protocol

- DNS basically send/recv style message protocol. Here's the question, what's the answer?
- uses UDP for query/response
- may switch to TCP if response too big
- uses TCP to xfer database between primary and secondary domain servers (zone transfer)

DNS naming scheme

problems to overcome...

- arpanet /etc/hosts file grew too big
- 30000+ hosts named venus
- UUCP gave us human-names mixed with route, we don't want the route in the name a!b!c!d from a to d ...

we need hierarchical names with distributed control of naming authority



sirius.cs.pdx.edu.

from right to left (top to bottom), right is general, left is specific

- root is on right (1st dot)
- namespace is a tree of domains, root domain, edu, pdx, cs, etc

- DNS labels acc. to RFC 1032 should be short (12 chars max), but 64 is actually allowed
- names are case-insensitive
- can be relative; e.g., sirius, but smart resolver
 code has to append local domainname correctly
- edu is TOP level domain
- pdx is 2nd-level. Apply to NIC registration service (http://rs.internic.net) for 2nd-level

they must guarantee uniqueness
 Jim Binkley

- internic manages root and 2nd-level
- local admins manage 3rd-level (or more) and can distribute that management locally if size warrants
- the rest of the world is NOT involved in how we manage our DNS names internally

names

 domain name - tree is made up of hierarchy of domains, each node is a domain and we go from right to left

edu, pdx.edu, cs.pdx.edu,

- labels are unique only within a domain
- absolute domain name has root dot and is called
 FQDN Fully Qualified Domain Name
- if relative, name must be completed somehow

names

- zone: namespace is partitioned into zones, zones are a set of names managed by 1 server, server will have > 1 zone to manage
- typically DNS server must have partner server elsewhere for redundancy, manages its zone and somebody else's too

DNS supports > 1 name hierarchy

organizational

reverse lookup

country

Liberian freighter registry? (not yet...)



organizational domains

mostly but not all U.S. com - commercial edu - educational gov - U.S. goverment int - international organizations mil - U.S. military net - networks org - other organizations

recent proposal - new domains

- arts entertainment/cultural
- .firm firms or businesses
- .info information providers
- .nom personal names
- .rec entertainment/recreation entities
- store bus offering goods for purchase
- .web WWW related entities

or alternative registries even

name.space

AlterNIC

- .exp, .ltd, .lnx, .med, .nic, and .xxx



Acc. To ICANN, recent new TLDs (wotsa TLD?)

- .aero, air transport industry
- ♦ .biz
- .coop, cooperatives
- .info
- .museum
- .name, individual humans
- .pro lawyers, Drs, accountants





cnri.reston.va.us

DNS record types

- adding a record types isn't easy
- name has type, client must ask for it by type
- Basic = ADDRESS, PTR, MX
- Zone=SOA (start of authority), NS (name server, server for a zone)
- optional = CNAME (alias), HINFO (host info), RP (responsible person), WKS (deprecated, but well known services), TXT (text)

record types

you can start to see that a query = (name, type); (sirius.cs.pdx.edu.,A)
a name may map to more than one item, e.g., sirius.cs.pdx.edu, type A, might have 1-n ip addresses associated with it

protocol header

- typically UDP request, reply
- request/reply format
- request = fixed header + question section
- reply = fixed header + query + answer sections
- question = (name, type, class = IP) with name in compressed format
- reply = set of **Resource Records** (RR)

DNS query/response format



node: id field used by client to match responses

header - flags field



QR - 0 if query, 1 if response

opcode - 0, standard query; 1, inverse; 2, server status request

AA - authoritative answer. Answer is from NS that maintains zone

TC - truncated. Using UDP, and answer was > 512 bytes

RD - recursion desired. If not set then iterative, NS may return list of other NS servers to try.

RA - recursion available. Set if server supports recursion rcode - error codes, 0 means no error. Only set if name invalid. Jim Binkley

question portion

- query part logically (name, type, class)
 class is IP
- type is A, NS, PTR, CNAME, MX, etc.
- name is encoded and variable length with a count (size) for each label in the name
- typically 1 question (and no answers yet...)

Resource Record format



RR format

domain name - same format as in query
TTL - in seconds, value often 86400, 1 day
resource data, the answer, for example if the query was foo.com, the answer might be 192.12.1.2

theory of operation/server-side

- person responsible for zone supplies 1 or more name servers
- primary name server and 1 or more secondary servers elsewhere necessary
- hope is to avoid single point of failure
- admin must change zone file and notify name server to reload (signal to named on UNIX)
- currently sans security, can't change zone info dynamically, must be manual

operation - root servers

- name server must contact other name servers for non-local info
- need IP address to contact... hmm....
- each name server has list of root servers, need IP addresses for them
- root servers know 2nd-level servers, refer name servers to them
- see how to get list of root names later

operation

- port 53 used for queries/answers
- may use UDP/TCP if answer too big
- TCP used for zone transfer between primary and secondary servers

how it works

• % telnet local.machine.com

- DNS is invoked to do name mapping to get IP address. (arp before udp before tcp...)
- client "resolver" code must have addresses for local name servers to prime the pump
- on UNIX, /etc/resolv.conf

 gethostbyname(3) calls resolver code which uses info in *resolv.conf* file Jim Binkley

/etc/resolv.conf

#

domain cs.pdx.edu # search cs.pdx.edu. ee.pdx.edu. # nameserver 131.252.20.183

nameserver 131.252.20.2 nameserver 128.95.120.1

operation - types of name resolution

- ◆ 2 ways to resolve a name used in DNS
- iterative contact name servers one at a time.
 You ask a root server and it gives you a list of contacts (NS + A records)
- recursive ask a name server to do the whole job
- resolvers are assumed to be simple and do things recursively
- root servers are busy and are iterative

when server gets query

checks to see if name lies in its zones if so

translates name to A and returns RRs if not, acc. to "howto" code if recursive contacts other servers (root/2ndary) until it gets A, and returns it else iterative (e.g., root server)

returns NS with A records

1 query - how many answers?

 if you ask for sirius.cs.pdx.edu, you get two A records (last time I checked)

 multi-homed host > 1 address record by definition

local lookup:

send request to local server

server returns authoritative response



remote lookup

- ask local NS (recursive)
- local NS may return cached hit. This is called a non-authoritative answer
- may ask root server in iterative fashion, get list of secondary servers
- ask secondary server, cache result
- return result to client

remote lookup picture



caching

- we don't want to work down the tree every time as root servers have a heavy load
- server data from remote systems thus kept locally with timeout (1 day typical)
- without timeout we would have remote systems with DNS, ip mapping where ip address might change

cached data is non-authoritative
 Jim Binkley

reverse address lookup

- is Alex Trebec involved with the DNS?
- pose an answer (IP address), get a question (DNS name)...
- gethostbyaddr(3), ip in, DNS out
- make a pointer query (PTR)
- special tree using in-addr.arpa domain
- class A, B, C net part at top, host at bottom

reverse address lookup

- if we didn't have tree organized by address, how would we do it?
- search through from domain roots? (take a lifetime...)
- gethostbyaddr(3) often used as weak security check, get DNS domain part and make sure it is in xyzzy domain (foo.com)
- gethostbyaddr(3) may automatically check that returned ip matches name by internal gethostbyname(3) call



nslookup - debug utility

- it's a shell, help and exit are commands
- from the command-line

% nslookup foo.com

- 1.2.3.4 <----- you (may) get an answer
- basically a resolver (client-side) for testing

to do reverse-pointer lookup
 -> set type=PTR
 -> 1.254.138.128.in-addr.arpa.

 yes, you have to reverse the address to put it in normal form, above is for 128.138.254.1

named - config

 /usr/somewhere/in.named - BSD named DNS server, started as server at boot /etc/named.boot - named configuration file not inetd based, start after syslog roughly: /usr/etc/in.named -r /etc/named.boot • named.boot tells named where to find database files Jim Binkley

named files

- besides named.boot
- /somewhere/named/named.ca (name may be different) - cache of root servers as RR records
- *hosts* primary RR (A records, MX, CNAME, etc), one or more
- *reverse pointer* PTR query records, one or more

named.boot - sample

- ; named.boot
- ; pdx-domain
- ; directory means where the db files hang out
- directory /usr/local/lib/named
- ; root cache
- cache named.ca
- ; what kind of server/domain name/filename
- primary 0.0.127.in-addr.arpa localhost
- primary cs.pdx.edu zone.pdx.cs
- primary 20.252.131.in-addr.arpa revp.131.252.20
- primary 21.252.131.in-addr.arpa revp.131.252.21

more details on boot file

- **primary** keyword indicates this named is primary server for zone info in that file
- secondary keyword says that host is secondary nameserver, need ip address + cache file name
- single name server can support more than one zone
- may have caching only server, no data files
- **forwarders**, can point interior NS at one NS that will cache all external queries
- **xfrnets** can restrict who can do zone tranfer

root cache

- get one from
 - ftp://rs.internic.net/domain/named.cache
- need to "prime the pump", 9 currently
- "dot" in NS record is for domain field and means
 "root domain"
- can use nslookup to check given 1 root server
 - -> server = a.root-servers.net.
 - -> set type=ns
 - -> .

root cache example (don't use)

```
; /domain/named.root
.
                     3600000 IN NS a.root-servers.net.
                     3600000 A 198.41.0.4
a.root-servers.net.
•
                     3600000 IN NS b.root-servers.net.
                     3600000 A 128.9.0.107
b.root-servers.net.
9
etc... through
i.root-servers.net.
```

RR - SOA

SOA - source of authority and start of zone db

- ; start of authority
- ; @ current zone symbol
- @ IN SOA foo.com. joebob@foo.com. (
 - 1001; serial #, increment when change zone
 - ; used to make zone xfer happen
 - 21600; 6 hours, time to zone xfer

1800; 30 minutes, 2nd retry if primary missing1209600; expire, 2 weeks for 2nd if no 1st432000); minimum, 5 days, ttl on RRs

RR - NS

- NS name server record. Assume foo.com
- ; ns, name servers for this domain, all primarys and
- ; secondaries
- foo.com. IN NS a.foo.com.
- foo.com. IN NS z.somethingelse.com.
- ; ns records for the current domain are not necessary
- ; needed however in parent for subdomains
- ; don't forget the A records for the NS records admin IN NS a.admin.foo.com. marketing IN NS z.mrkt.foo.com.

RR - A and other records

Address red	cords	are heart of DNS
; the data p	art	
localhost	IN	A 127.0.0.1
foobar	IN	A 192.0.0.1
snarp	IN	A 192.0.0.2
flozzle	IN	A 192.0.0.3
WWW	IN	CNAME foobar.foo.com.
foo.com	IN	MX foobar.foo.com.
multi	IN	A 192.0.0.4
	IN	A 192.0.0.5

update of zone files

- change the serial # in the SOA record
- add the A records or whatever
- kill -HUP signal to named
- secondary won't change until refresh time passed, will check serial # first, cache current data in "data file"
- zone xfer done with TCP, port 53

DNS futures

- types are attribute OF the name, not in the name
- not easy to change types since you must change code everywhere, but it happens
- e.g., there are ISDN/X.25 record types
- would like DNS to be dynamic but 1st
- DNS needs a secure protocol

DNSSEC

- IETF working group
- add security to DNS
- dynamic DNS updates
- KEY/SIG records using RSA public keys
- zone authority may sign all DNS records
- new USER record as well
- hence can authenticate user@dns