# Overview and History of Operating Systems

These are the notes for lecture 1.

Please review the "Syllabus" notes before these.

## **Overview / Historical Developments**

An Operating System... Sits between hardware and users **Provides "environment" to execute programs** Like a government No useful work **Regulates workers** Manages, allocates resources **CPU** (execution time) **Memory Space Disk / File storage I/O Devices Control Prevent incorrect use of hardware Security / Protection** 

## Goals

Make computer *easy* to use Make computer more *efficient* Help user *solve problems / do work* 

> Ease of use Efficiency

## Overview





## **Early Computers**

**Input Devices:** • Card Reader

**Output Devices:** • Printer

• Card Puncher

	11				ŝ																																		Ĩ	ġ
					1						9				ä			0	T.	A	-	4	15	2	æ											ŝ	1	ł.		
-	2	F.C.	-	-	11	-	111	11	-	1	-			11	-		-	-	-	11	-	Tim	F				11	1	FIF	-	11	11	10	11		-			-	
	-	111		14	-	11	-		11			1			-	11	i							1		10	1	0	1		11	1	13	11	1	ŝ		-	1	1
1221	ee.	111	00	13	11	11	33	63	1)		100	1	1	P	1	l	1	N	9	1	1	1	1	13	1	-	1	9	9	1	8	Ņ	13	1	1	2	1	11	1	1
ALC: NO	12								-											***					**											8	-		1	l
	4	10	1	-		i	-		1.1	1.1	1.1.7										-			15		LI		1		4		11		4.1		į,	17	1		
011		111	ġ,		11	ñ	U.		Î			Ē	ie.	ö	ŝ	ő	ö				8		9	Ū.	8	1	0	0	8	Ö,	ñ	0	ŋ	0	18	8		0)	g	h
	17	-	1	-	2	I		1.1	0	1				-	1	1	2		1	11	-		9	11	0	-	1.1	1	11	0	-	2	ņ		1	2	2			1
2222	84		8		2	8	11	8	5	Đ			8	H	2		8		B	13			1	H	33	23	1	10	13	ų	11	1	8	8	U.	르	H	8	8	h

Punch card

## (No disk, no secondary storage)

A "job"

**User prepares input cards** (**Program, Data**) User gets time on the machine Loads the program **Executes it** Study the output and come back tomorrow

## **Early Computers**

The **O.S....** 

(Simply a "control program")

- Read cards
- Load memory
- Transfer control to user code
- Print out contents of memory ("*core dump*")
- Loop to next "job"

## **Batch Operating System**

Read jobs from cards Some jobs are "control cards"

\$end,dump	
	the data
Şdata	
\$execute	
\$load	
	f the source code
\$compile,fortran	

## **New Technology: Magnetic Tape**

Idea:Read cards onto tapeTo output a card...Write to tape & punch it laterSame for printing



## **New Technology: Magnetic Tape**

**Idea:** Read cards onto tape To output a card... Write to tape & punch it later Same for printing

Concept: SPOOLING

- Use tape as a "buffer".
- Allows I/O from one job to overlap the computation from another job!

(When disks were invented, they were first used like this.)

## A Mainframe System (circa 1960)



**New Technology: The Disk** 

The first disks were used for spooling. (The "file" was invented later.)

**Concept: <u>The Job Pool</u>** 

Several jobs are waiting to be executed

- One job in memory
- Future jobs sitting on disk

**New Technology: The Disk** 

The first disks were used for spooling. (The "file" was invented later.)

**Concept: <u>The Job Pool</u>** 

Several jobs are waiting to be executed

- One job in memory
- Future jobs sitting on disk

The O.S. can make its first decision!

Which job to run next? **First-come, first served** VS. **Job Scheduling** 

## **Multi-programming**

Idea: Keep several jobs in memory at once!

When one job waits on I/O... another jobs can use the CPU!

#### **Increases CPU utilization**

Don't keep entire job pool in memory (just select a few) Job X starts I/O... OS selects another job to run. CPU does not sit idle. When job X finishes... OS selects another job from job pool

and loads it into memory.

13

## **Main Memory**



14

## **Main Memory**

#### **Simplest Approach:**

- The entire job is laoded into a contiguous range of memory
- No protection between programs
- A job runs until it requests I/O No "*time-slicing*"

If it loops... Oh, well...

• When complete, it transfers back to the OS.

#### **Memory Management**

A big topic (a chapter in textbook)

#### Job Scheduling

When job X blocks (due to I/O)...... which job will be run next?"Process Management" (chapter 2)

## **New Concept: Terminals**

Combines with multiprogramming! Users want to interact with the programs *while they are executing!* 

**Initial motivations for interaction:** 

- Deal with contingencies during job execution.
- Debugging programs.

The good 'ole days...

- Long turn-around times.
- Programs were smaller.
- Written in assembly / FORTRAN.
- Programmers were very careful.
- Bugs were intolerable.

**Batch Operating Systems** 

Each job runs non-stop (until it decides to perform I/O)

#### **Problems when used with interactive terminals:**

- User response times are...
  - Long Unpredictable
- Bugs in one program...
  - Crash the entire system
  - **Really annoyed people using terminals** (With batch jobs, just restart the programs)

## **New Concept: Time-Sharing**

"Multi-tasking"

Goal:

Make it seem like every user has a dedicated computer!

The Idea:

Switch the CPU rapidly between jobs so that every running job seems to be making regular progress.

New hardware support required:

**Periodic "timer interrupts"** 

Ways to protect one program from the other programs.

**<u>Goal:</u>** isolate bugs / loops to just that program!

## **Historical Context**

Each user has a Cathode Ray Tube (CRT) terminal. Each user types a command & it is executed. "Response time" -- should be < 1 second. Many users "online", sharing the CPU. Demonstrated in 1960, common in 1970s. Files kept on disks...

but lots of implementation details still visible. (blocking, file formats, sectors, etc...)

**UNIX** (1969, spreading in mid-1970's)

## **Time-Sharing**

Several processes in memory. Memory management & protection are required. If size of user job is really large... Swap other users' jobs out to disk (temporarily) *"Virtual Memory"* Don't load the entire job into memory

**Disk Management:** 

Directory Structures Ease-of-use Protection from other users

**Communication between users:** 

Still an active area Synchronization, coordination

20

## **Personal Computers**

Began appearing in the 1970s.

**Shift of emphasis:** 

- Hardware utilization is less important
- Maximize: User convenience Fast response time Ease of programming
- Security vs. Communication

Early opinions: Protection is unimportant Each person's computer is separate & isolated.

### Trends

Hardware costs will continue to fall. OS's will appear embedded in more devices! More kinds of OS's will be needed!

**Computer users will become more sophisticated.** *Better OS's will be demanded!* 

**Features on research / high-end systems will become common on small, inexpensive systems.**  *Parallel processing / multi-processing Real-time control Distributed systems* 

**Programs will become more complex.** *The OS will need to promote ease of programming / use!* 

More malicious programs will be created. More security will be needed!

## **Kinds of Operating Systems**

Personal Computers Large, super-computers Parallel multi-processors Embedded computers Real-time computers Distributed systems Highly reliable systems Super-low cost Research platforms

#### **Applications**

- Chess playing
- Dishwasher / microwave
- Flight control / space shuttle
- Military command
- Automobile control
- Lab / factory automation
- Assembly robot control
- Corporation management
- Web server
- Web search engine
- Nuclear reactor control
- Toys
- Artificial Intelligence

... etc ...

## **Parallel Systems**

Single processor vs. multi-processor systems

**Goals:** 

• Increase *"Throughput"* 

Get more work done, per hour

- Utilize small, inexpensive processors ... To get more horsepower (giga-flops) Example: Graphics co-processor
- Save money by sharing expensive peripheral devices.
- Reliability

**Graceful degradation Fault-tolerance** 



### **Fault-Tolerance**

# **Example:** Tandem System

**Two identical processors** 

- Primary
- Backup

Each has its own memory

**Operating in lock-step** 

Failure detected? Backup becomes the primary *"hot backup"* 

## **Types of Multiprocessing Systems**

#### **Tightly-Coupled Systems**

- Share the system bus
- Share peripheral I/O devices
- Share memory (sometimes)
- Share a common clock (sometimes)

**Example: Graphics Co-processor** 

**Lossely-Coupled Systems** 

26

**Types of Tightly-Coupled Systems** 

"Symmetric Multiprocessing"

- Each processor runs identical copy of OS
- OS code resides in shared memory
- Shared data structures (for concurrency control)

"Asymmetric Multiprocessing"

- Master-slave relationship
- One processor assigns tasks to others
- Increased specialization
  - ---> decreased reliability
- Trend: cheap processors

**Offload tasks to slaves or back-ends** 

**Graphics processor** 

**Disk processor** 

**Processor in keyboard / mouse** 

• Communication/interfacing becomes paramount.

27

## **Distributed Systems**

**Processors do not share** Memory, Clock, System bus, Devices "Loosely-coupled systems" Communication Slow (internet)  $\leftarrow$  Fast (specialized bus) **Motivations Communication & sharing** Email, shared data, group work **Reliability** No single failure should crash the entire system Data should remain accessible **Computation speed-up** Load-sharing **Recource sharing** Access to specialized or unique hardware

## **Real-Time Operating Systems**

Used to control a physical process Sensors collect data ("input"?) Actuators control the physical process ("output"?)

Examples:

- Avionics (space shuttle control system)
- Medical systems (heart monitor)
- Industrial system (oil refinery)
- Military (anti-missile laser control)
- Consumer products (automobile controller)

The key...

**Rigid**, well-defined, fixed time contraints

Time-sharing systems <u>should</u> repond quickly! Real-time systems <u>must</u> respond quickly!

## **Real-Time Operating Systems**

## **Soft Real-Time Systems**

Some processes are given higher "*priorities*." Example: video & music playback

**Common in many OS's.** 

Adequate for many applications, ... but too risky for some!

Even though response is often fast enough, it is never guaranteed! **Real-Time Operating Systems** 

**Hard Real-Time Systems** 

Guarantees that task will complete by their deadlines.

All potential delays are <u>bounded</u>. Want to avoid:

- **Disks** (highly variable latencies)
- Virtual Memory (complex & unpredictable)

OS tends to be low-level, minimal, close to hardware.

A specialized sub-field of OS.