# Finding the Roots of $f(x)=0$ 

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## Overview

Topics covered in this chapter

- Preliminary considerations and bracketing.
- Fixed Point Iteration
- Bisection
- Newton's Method
- The Secant Method
- Hybrid Methods: the built in fzero function
- Roots of Polynomials


## Example: Picnic Table Leg

Computing the dimensions of a picnic table leg involves a root-finding problem.

Leg assembly


Detail of one leg


## Example: Picnic Table Leg

Dimensions of a the picnic table leg satisfy

$$
w \sin \theta=h \cos \theta+b
$$

Given overall dimensions $w$ and $h$, and the material dimension, $b$, what is the value of $\theta$ ? An analytical solution for $\theta=f(w, h, b)$ exists, but is not obvious.

Use a numerical root-finding procedure to find the value of $\theta$ that satisfies

$$
f(\theta)=w \sin \theta-h \cos \theta-b=0
$$

## Roots of $f(x)=0$

Any function of one variable can be put in the form $f(x)=0$.

## Example:

To find the $x$ that satisfies

$$
\cos (x)=x
$$

find the zero crossing of

$$
f(x)=\cos (x)-x=0
$$



## General Considerations

- Is this a special function that will be evaluated often?
- How much precision is needed?
- How fast and robust must the method be?
- Is the function a polynomial?
- Does the function have singularities?

There is no single root-finding method that is best for all situations.

## Root-Finding Procedure

The basic strategy is

1. Plot the function.
$>$ The plot provides an initial guess, and an indication of potential problems.
2. Select an initial guess.
3. Iteratively refine the initial guess with a root-finding algorithm.

## Bracketing

A root is bracketed on the interval $[a, b]$ if $f(a)$ and $f(b)$ have opposite sign. A sign change occurs for singularities as well as roots


Bracketing is used to make initial guesses at the roots, not to accurately estimate the values of the roots.

## Bracketing Algorithm (1)

## Algorithm 6.1 Bracket Roots

$$
\begin{aligned}
& \text { given: } f(x), x_{\min }, x_{\max }, n \\
& d x=\left(x_{\max }-x_{\min }\right) / n \\
& x_{\text {left }}=x_{\min } \\
& \text { i }=0 \\
& \text { while } i<n \\
& \quad i \leftarrow i+1 \\
& \quad x_{\text {right }}=x_{\text {left }}+d x \\
& \text { if } f(x) \text { changes sign in }\left[x_{\text {left }}, x_{\text {right }}\right] \\
& \quad \text { save }\left[x_{\text {left }}, x_{\text {right }}\right] \text { for further root-finding } \\
& \quad \text { end } \\
& \quad x_{\text {left }}=x_{\text {right }} \\
& \text { end }
\end{aligned}
$$

## Bracketing Algorithm (2)

```
A simple test for sign change: }f(a)\timesf(b)<0\mathrm{ ?
or in Matlab
    if
fa = ..
fb = ...
if fa*fb < 0
    save bracket
end
```

but this test is susceptible to underflow.

## Bracketing Algorithm (3)

A better test uses the built-in sign function

```
fa = ..
fb = ...
if sign(fa)~}=\operatorname{sign(fb)
    save bracket
end
```

See implementation in the brackPlot function

## The brackPlot Function

brackPlot is a NMM toolbox function that

- Looks for brackets of a user-defined $f(x)$
- Plots the brackets and $f(x)$
- Returns brackets in a two-column matrix


## Syntax:

```
brackPlot('myFun',xmin,xmax)
brackPlot('myFun',xmin, xmax, nx)
```

where

| myFun | is the name of an $m$-file that evaluates $f(x)$ |
| :--- | :--- |
| $x m i n, x m a x$ | define range of $x$ axis to search |
| $n x$ | is the number of subintervals on [xmin, xmax] used to <br> check for sign changes of $f(x)$. Default: $n x=20$ |

## Apply brackPlot Function to $\sin (x)$

```
>> Xb = brackPlot('sin', 4*pi,4*pi)
Xb =
    -12.5664 -11.2436
    -9.9208 -8.5980
    -7.2753 -5.9525
    -3.3069 -1.9842
    -0.6614 0.6614
        1.9842 3.3069
        5.9525 7.2753
        8.5980 9.9208
    11.2436 12.5664
```



## Apply brackPlot to a user-defined Function (1)

To solve

$$
f(x)=x-x^{1 / 3}-2=0
$$

we need an m -file function to evaluate $f(x)$ for any scalar or vector of $x$ values.

File fx3.m:
Note the use of the array operator.

```
function f = fx3(x)
% fx3 Evaluates f(x) = x - x^(1/3) - 2
f = x - x.^(1/3) - 2;
```

Run brackPlot with fx3 as the input function

```
>> brackPlot('fx3',0,5)
ans =
    3.4000 3.6000
```


## Apply brackPlot to a user-defined Function (2)

```
>> Xb = brackPlot('fx3',0,5)
Xb =
    3.4211 3.6842
```



## Apply brackPlot to a user-defined Function (3)

Instead of creating a separate m-file, we can use an in-line function object.

```
>> f = inline('x - x.^(1/3) - 2')
f =
    Inline function:
    f(x) = x - x.^(1/3) - 2
>> brackPlot(f,0,5)
ans =
    3.4000 3.6000
```

Note: When an inline function object is supplied to brackPlot, the name of the object is not surrounded in quotes:

$$
\text { brackPlot }(f, 0,5) \quad \text { instead of brackPlot('fun',0,5) }
$$

## Root-Finding Algorithms

We now proceed to develop the following root-finding algorithms:

- Fixed point iteration
- Bisection
- Newton's method
- Secant method

These algorithms are applied after initial guesses at the root(s) are identified with bracketing (or guesswork).

## Fixed Point Iteration

Fixed point iteration is a simple method. It only works when the iteration function is convergent.

Given $f(x)=0$, rewrite as $x_{\text {new }}=g\left(x_{\text {old }}\right)$

## Algorithm 6.2 Fixed Point Iteration

$$
\begin{aligned}
& \text { initialize: } x_{0}=\ldots \\
& \text { for } k=1,2, \ldots \\
& \quad x_{k}=g\left(x_{k-1}\right) \\
& \text { if converged, stop } \\
& \text { end }
\end{aligned}
$$

## Convergence Criteria

An automatic root-finding procedure needs to monitor progress toward the root and stop when current guess is close enough to the desired root.

- Convergence checking will avoid searching to unnecessary accuracy.
- Convergence checking can consider whether two successive approximations to the root are close enough to be considered equal.
- Convergence checking can examine whether $f(x)$ is sufficiently close to zero at the current guess.

More on this later . . .

## Fixed Point Iteration Example (1)

To solve

$$
x-x^{1 / 3}-2=0
$$

rewrite as

$$
x_{\text {new }}=g_{1}\left(x_{\text {old }}\right)=x_{\mathrm{old}}^{1 / 3}+2
$$

or

$$
x_{\text {new }}=g_{2}\left(x_{\text {old }}\right)=\left(x_{\text {old }}-2\right)^{3}
$$

or

$$
x_{\mathrm{new}}=g_{3}\left(x_{\mathrm{old}}\right)=\frac{6+2 x_{\mathrm{old}}^{1 / 3}}{3-x_{\mathrm{old}}^{2 / 3}}
$$

Are these $g(x)$ functions equally effective?

## Fixed Point Iteration Example (2)

$$
\begin{aligned}
& g_{1}(x)=x^{1 / 3}+2 \\
& g_{2}(x)=(x-2)^{3} \\
& g_{3}(x)=\frac{6+2 x^{1 / 3}}{3-x^{2 / 3}}
\end{aligned}
$$

| $k$ | $g_{1}\left(x_{k-1}\right)$ | $g_{2}\left(x_{k-1}\right)$ | $g_{3}\left(x_{k-1}\right)$ |
| :---: | :---: | :---: | :---: |
| 0 | 3 | 3 | 3 |
| 1 | 3.4422495703 | 1 | 3.5266442931 |
| 2 | 3.5098974493 | -1 | 3.5213801474 |
| 3 | 3.5197243050 | -27 | 3.5213797068 |
| 4 | 3.5211412691 | -24389 | 3.5213797068 |
| 5 | 3.5213453678 | $-1.451 \times 10^{13}$ | 3.5213797068 |
| 6 | 3.5213747615 | $-3.055 \times 10^{39}$ | 3.5213797068 |
| 7 | 3.5213789946 | $-2.852 \times 10^{118}$ | 3.5213797068 |
| 8 | 3.5213796042 | $\infty$ | 3.5213797068 |
| 9 | 3.5213796920 | $\infty$ | 3.5213797068 |

Summary: $g_{1}(x)$ converges, $g_{2}(x)$ diverges, $g_{3}(x)$ converges very quickly

## Bisection

Given a bracketed root, halve the interval while continuing to bracket the root


## Bisection (2)

For the bracket interval $[a, b]$ the midpoint is

$$
x_{m}=\frac{1}{2}(a+b)
$$

A better formula, one that is less susceptible to round-off is

$$
x_{m}=a+\frac{b-a}{2}
$$

## Bisection Algorithm

## Algorithm 6.3 Bisection

$$
\begin{aligned}
& \text { initialize: } a=\ldots, b=\ldots \\
& \text { for } k=1,2, \ldots \\
& \quad x_{m}=a+(b-a) / 2 \\
& \text { if } \operatorname{sign}\left(f\left(x_{m}\right)\right)=\operatorname{sign}\left(f\left(x_{a}\right)\right) \\
& \quad a=x_{m} \\
& \text { else } \\
& \quad b=x_{m} \\
& \text { end } \\
& \text { if converged, stop } \\
& \text { end }
\end{aligned}
$$

## Bisection Example

Solve with bisection:

$$
x-x^{1 / 3}-2=0
$$

| $k$ | $a$ | $b$ | $x_{\text {mid }}$ | $f\left(x_{\text {mid }}\right)$ |
| :---: | :--- | :--- | :--- | ---: |
| 0 | 3 | 4 |  |  |
| 1 | 3 | 4 | 3.5 | -0.01829449 |
| 2 | 3.5 | 4 | 3.75 | 0.19638375 |
| 3 | 3.5 | 3.75 | 3.625 | 0.08884159 |
| 4 | 3.5 | 3.625 | 3.5625 | 0.03522131 |
| 5 | 3.5 | 3.5625 | 3.53125 | 0.00845016 |
| 6 | 3.5 | 3.53125 | 3.515625 | -0.00492550 |
| 7 | 3.51625 | 3.53125 | 3.5234375 | 0.00176150 |
| 8 | 3.51625 | 3.5234375 | 3.51953125 | -0.00158221 |
| 9 | 3.51953125 | 3.5234375 | 3.52148438 | 0.00008959 |
| 10 | 3.51953125 | 3.52148438 | 3.52050781 | -0.00074632 |

## Analysis of Bisection (1)

Let $\delta_{n}$ be the size of the bracketing interval at the $n^{t h}$ stage of bisection. Then
$\delta_{0}=b-a=$ initial bracketing interval
$\delta_{1}=\frac{1}{2} \delta_{0}$
$\delta_{2}=\frac{1}{2} \delta_{1}=\frac{1}{4} \delta_{0}$
:
$\delta_{n}=\left(\frac{1}{2}\right)^{n} \delta_{0}$

$$
\begin{array}{rlrl}
\Longrightarrow & \frac{\delta_{n}}{\delta_{0}} & =\left(\frac{1}{2}\right)^{n}=2^{-n} \\
\text { or } & & n & =\log _{2}\left(\frac{\delta_{n}}{\delta_{0}}\right)
\end{array}
$$

## Analysis of Bisection (2)

$$
\begin{gathered}
\frac{\delta_{n}}{\delta_{0}}=\left(\frac{1}{2}\right)^{n}=2^{-n} \quad \text { or }
\end{gathered} \begin{array}{cc}
n=\log _{2}\left(\frac{\delta_{n}}{\delta_{0}}\right) \\
& \begin{array}{ccc}
n & \frac{\delta_{n}}{\delta_{0}} & \begin{array}{c}
\text { function } \\
\text { evaluations }
\end{array} \\
\hline 5 & 3.1 \times 10^{-2} & 7 \\
10 & 9.8 \times 10^{-4} & 12 \\
20 & 9.5 \times 10^{-7} & 22 \\
30 & 9.3 \times 10^{-10} & 32 \\
40 & 9.1 \times 10^{-13} & 42 \\
50 & 8.9 \times 10^{-16} & 52 \\
\hline
\end{array}
\end{array}
$$

## Convergence Criteria

An automatic root-finding procedure needs to monitor progress toward the root and stop when current guess is close enough to the desired root.

- Convergence checking will avoid searching to unnecessary accuracy.
- Check whether successive approximations are close enough to be considered the same:

$$
\left|x_{k}-x_{k-1}\right|<\delta_{x}
$$

- Check whether $f(x)$ is close enough zero.

$$
\left|f\left(x_{k}\right)\right|<\delta_{f}
$$

## Convergence Criteria on $x$



Absolute tolerance: $\left|x_{k}-x_{k-1}\right|<\delta_{x}$
Relative tolerance: $\left|\frac{x_{k}-x_{k-1}}{b-a}\right|<\hat{\delta}_{x}$
$x_{k}=$ current guess at the root
$x_{k-1}=$ previous guess at the root

## Convergence Criteria on $f(x)$



Absolute tolerance: $\left|f\left(x_{k}\right)\right|<\delta_{f}$
Relative tolerance:

$$
\left|f\left(x_{k}\right)\right|<\hat{\delta}_{f} \max \left\{\left|f\left(a_{0}\right)\right|,\left|f\left(b_{0}\right)\right|\right\}
$$

where $a_{0}$ and $b_{0}$ are the original brackets

## Convergence Criteria on $f(x)$

If $f^{\prime}(x)$ is small near the root, it is easy to satisfy a tolerance on $f(x)$ for a large range of $\Delta x$. A tolerance on $\Delta x$ is more conservative.


If $f^{\prime}(x)$ is large near the root, it is possible to satisfy a tolerance on $\Delta x$ when $|f(x)|$ is still large. A tolerance on $f(x)$ is more conservative.


## Newton's Method

For a current guess $x_{k}$, use $f\left(x_{k}\right)$ and the slope $f^{\prime}\left(x_{k}\right)$ to predict where $f(x)$ crosses the $x$ axis.


## Newton's Method (2)

Expand $f(x)$ in Taylor Series around $x_{k}$

$$
f\left(x_{k}+\Delta x\right)=f\left(x_{k}\right)+\left.\Delta x \frac{d f}{d x}\right|_{x_{k}}+\left.\frac{(\Delta x)^{2}}{2} \frac{d^{2} f}{d x^{2}}\right|_{x_{k}}+\ldots
$$

Substitute $\Delta x=x_{k+1}-x_{k}$ and neglect second order terms to get

$$
f\left(x_{k+1}\right) \approx f\left(x_{k}\right)+\left(x_{k+1}-x_{k}\right) f^{\prime}\left(x_{k}\right)
$$

where

$$
f^{\prime}\left(x_{k}\right)=\left.\frac{d f}{d x}\right|_{x_{k}}
$$

## Newton's Method (3)

Goal is to find $x$ such that $f(x)=0$.
Set $f\left(x_{k+1}\right)=0$ and solve for $x_{k+1}$

$$
0=f\left(x_{k}\right)+\left(x_{k+1}-x_{k}\right) f^{\prime}\left(x_{k}\right)
$$

or, solving for $x_{k+1}$

$$
x_{k+1}=x_{k}-\frac{f\left(x_{k}\right)}{f^{\prime}\left(x_{k}\right)}
$$

## Newton's Method Algorithm

Algorithm 6.4<br>initialize: $x_{1}=\ldots$<br>for $k=2,3, \ldots$<br>$x_{k}=x_{k-1}-f\left(x_{k-1}\right) / f^{\prime}\left(x_{k-1}\right)$<br>if converged, stop<br>end

## Newton's Method Example (1)

Solve:

$$
x-x^{1 / 3}-2=0
$$

First derivative is

$$
f^{\prime}(x)=1-\frac{1}{3} x^{-2 / 3}
$$

The iteration formula is

$$
x_{k+1}=x_{k}-\frac{x_{k}-x_{k}^{1 / 3}-2}{1-\frac{1}{3} x_{k}^{-2 / 3}}
$$

## Newton's Method Example (2)

|  | $x_{k+1}=x_{k}-\frac{x_{k}-x_{k}^{1 / 3}-2}{1-\frac{1}{3} x_{k}^{-2 / 3}}$ |  |  |
| :---: | :---: | :---: | :---: |
|  |  | $f^{\prime}\left(x_{k}\right)$ | $f(x)$ |
| $k$ | $x_{k}$ | 0.83975005 | -0.44224957 |
| 0 | 3 | 0.85612976 | 0.00450679 |
| 1 | 3.52664429 | $3.771 \times 10^{-7}$ |  |
| 2 | 3.52138015 | 0.85598641 | $3.7510^{-15}$ |
| 3 | 3.52137971 | 0.85598640 | $2.664 \times 10^{2}$ |
| 4 | 3.52137971 | 0.85598640 | 0.0 |

## Conclusion

- Newton's method converges much more quickly than bisection
- Newton's method requires an analytical formula for $f^{\prime}(x)$
- The algorithm is simple as long as $f^{\prime}(x)$ is available.
- Iterations are not guaranteed to stay inside an ordinal bracket.


## Divergence of Newton's Method



Since

$$
x_{k+1}=x_{k}-\frac{f\left(x_{k}\right)}{f^{\prime}\left(x_{k}\right)}
$$

the new guess, $x_{k+1}$, will be far from the old guess whenever $f^{\prime}\left(x_{k}\right) \approx 0$

## Secant Method (1)

Given two guesses $x_{k-1}$ and $x_{k}$, the next guess at the root is where the line through $f\left(x_{k-1}\right)$ and $f\left(x_{k}\right)$ crosses the $x$ axis.


## Secant Method (2)

Given

$$
\begin{aligned}
x_{k} & =\text { current guess at the root } \\
x_{k-1} & =\text { previous guess at the root }
\end{aligned}
$$

Approximate the first derivative with

$$
f^{\prime}\left(x_{k}\right) \approx \frac{f\left(x_{k}\right)-f\left(x_{k-1}\right)}{x_{k}-x_{k-1}}
$$

Substitute approximate $f^{\prime}\left(x_{k}\right)$ into formula for Newton's method

$$
x_{k+1}=x_{k}-\frac{f\left(x_{k}\right)}{f^{\prime}\left(x_{k}\right)}
$$

to get

$$
x_{k+1}=x_{k}-f\left(x_{k}\right)\left[\frac{x_{k}-x_{k-1}}{f\left(x_{k}\right)-f\left(x_{k-1}\right)}\right]
$$

## Secant Method

Two versions of this formula are equivalent in exact math:

$$
x_{k+1}=x_{k}-f\left(x_{k}\right)\left[\frac{x_{k}-x_{k-1}}{f\left(x_{k}\right)-f\left(x_{k-1}\right)}\right]
$$

and

$$
x_{k+1}=\frac{f\left(x_{k}\right) x_{k-1}-f\left(x_{k-1}\right) x_{k}}{f\left(x_{k}\right)-f\left(x_{k-1}\right)}
$$

Equation $(\star)$ is better since it is of the form $x_{k+1}=x_{k}+\Delta$. Even if $\Delta$ is inaccurate the change in the estimate of the root will be small at convergence because $f\left(x_{k}\right)$ will also be small.

Equation ( $\star \star$ ) is susceptible to catastrophic cancellation:

- $f\left(x_{k}\right) \rightarrow f\left(x_{k-1}\right)$ as convergence approaches, so cancellation error in the denominator can be large.
- $|f(x)| \rightarrow 0$ as convergence approaches, so underflow is possible


## Secant Algorithm

## Algorithm 6.5

```
initialize: \(x_{1}=\ldots, x_{2}=\ldots\)
for \(k=2,3 \ldots\)
    \(x_{k+1}=x_{k}\)
        \(-f\left(x_{k}\right)\left(x_{k}-x_{k-1}\right) /\left(f\left(x_{k}\right)-f\left(x_{k-1}\right)\right)\)
    if converged, stop
    end
```


## Secant Method Example

Solve:

$$
x-x^{1 / 3}-2=0
$$

| $k$ | $x_{k-1}$ | $x_{k}$ | $f\left(x_{k}\right)$ |
| :---: | :---: | :---: | :---: |
| 0 | 4 | 3 | -0.44224957 |
| 1 | 3 | 3.51734262 | -0.00345547 |
| 2 | 3.51734262 | 3.52141665 | 0.00003163 |
| 3 | 3.52141665 | 3.52137970 | $-2.034 \times 10^{-9}$ |
| 4 | 3.52137959 | 3.52137971 | $-1.332 \times 10^{-15}$ |
| 5 | 3.52137971 | 3.52137971 | 0.0 |

## Conclusions

- Converges almost as quickly as Newton's method.
- No need to compute $f^{\prime}(x)$.
- The algorithm is simple.
- Two initial guesses are necessary
- Iterations are not guaranteed to stay inside an ordinal bracket.


## Divergence of Secant Method



Since
$x_{k+1}=x_{k}-f\left(x_{k}\right)\left[\frac{x_{k}-x_{k-1}}{f\left(x_{k}\right)-f\left(x_{k-1}\right)}\right]$
the new guess, $x_{k+1}$, will be far from the old guess whenever $f^{\prime}\left(x_{k}\right) \approx f\left(x_{k-1}\right)$ and $|f(x)|$ is not small.

## Summary of Basic Root-finding Methods

- Plot $f(x)$ before searching for roots
- Bracketing finds coarse interval containing roots and singularities
- Bisection is robust, but converges slowly
- Newton's Method
$\triangleright$ Requires $f(x)$ and $f^{\prime}(x)$.
$\triangleright$ Iterates are not confined to initial bracket.
$\triangleright$ Converges rapidly.
$\triangleright$ Diverges if $f^{\prime}(x) \approx 0$ is encountered.
- Secant Method
$\triangleright$ Uses $f(x)$ values to approximate $f^{\prime}(x)$.
$\triangleright$ Iterates are not confined to initial bracket.
$\triangleright$ Converges almost as rapidly as Newton's method.
$\triangleright$ Diverges if $f^{\prime}(x) \approx 0$ is encountered.


## fzero Function (1)

fzero is a hybrid method that combines bisection, secant and reverse quadratic interpolation

## Syntax:

$r=$ fzero('fun', $x 0$ )
$r=$ fzero('fun', x0,options)
$r=$ fzero('fun', $x 0$,options, arg1, arg2,...)
x 0 can be a scalar or a two element vector

- If $x 0$ is a scalar, fzero tries to create its own bracket.
- If x 0 is a two element vector, fzero uses the vector as a bracket.


## Reverse Quadratic Interpolation

Find the point where the $x$ axis intersects the sideways parabola passing through three pairs of $(x, f(x))$ values.


## fzero Function (2)

fzero chooses next root as

- Result of reverse quadratic interpolation (RQI) if that result is inside the current bracket.
- Result of secant step if RQI fails, and if the result of secant method is in inside the current bracket.
- Result of bisection step if both RQI and secant method fail to produce guesses inside the current bracket.


## fzero Function (3)

Optional parameters to control fzero are specified with the optimset function.

## Examples:

Tell fzero to display the results of each step:

```
>> options = optimset('Display','iter');
>> x = fzero('myFun',x0,options)
```

Tell fzero to use a relative tolerance of $5 \times 10^{-9}$ :

```
>> options = optimset('TolX',5e-9);
>> x = fzero('myFun',x0,options)
```

Tell fzero to suppress all printed output, and use a relative tolerance of $5 \times 10^{-4}$ :

```
>> options = optimset('Display','off','TolX',5e-4);
>> x = fzero('myFun',x0,options)
```


## fzero Function (4)

Allowable options (specified via optimset):

\left.| Option type | Value | Effect |
| :--- | :--- | :--- |
| 'Display' | 'iter' | Show results of each iteration |
|  | 'final' | Show root and original bracket |$\right]$| Suppress all print out |
| :--- |

The default values of 'Display' and 'TolX' are equivalent to

```
options = optimset('Display','iter','TolX',eps)
```


## Roots of Polynomials

Complications arise due to

- Repeated roots
- Complex roots
- Sensitivity of roots to small perturbations in the polynomial coefficients (conditioning).



## Algorithms for Finding Polynomial Roots

- Bairstow's method
- Müller's method
- Laguerre's method
- Jenkin's-Traub method
- Companion matrix method


## roots Function (1)

The built-in roots function uses the companion matrix method

- No initial guess
- Returns all roots of the polynomial
- Solves eigenvalue problem for companion matrix

Write polynomial in the form

$$
c_{1} x^{n}+c_{2} x^{n-1}+\ldots+c_{n} x+c_{n+1}=0
$$

Then, for a third order polynomial

```
>> c = [lc1 c2 c3 c4];
>> r = roots(c)
```


## roots Function (2)

The eigenvalues of

$$
A=\left[\begin{array}{cccc}
-c_{2} / c_{1} & -c_{3} / c_{1} & -c_{4} / c_{1} & -c_{5} / c_{1} \\
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 0
\end{array}\right]
$$

are the same as the roots of

$$
c_{5} \lambda^{4}+c_{4} \lambda^{3}+c_{3} \lambda^{2}+c_{2} \lambda+c_{1}=0 .
$$

## roots Function (3)

The statements

```
c = ... % vector of polynomial coefficients
r = roots(c);
```

are equivalent to

```
c = ...
n = length(c);
A = diag(ones(1,n-2),-1); % ones on first subdiagonal
A(1,:) = -c(2:n) ./ c(1); % first row is -c(j)/c(1), j=2..n
r = eig(A);
```


## roots Examples

Roots of

$$
\begin{aligned}
& f_{1}(x)=x^{2}-3 x+2 \\
& f_{2}(x)=x^{2}-10 x+25 \\
& f_{3}(x)=x^{2}-17 x+72.5
\end{aligned}
$$

are found with

```
>> roots([1 -3 2])
ans =
        2
        1
>> roots([1 -10 25])
ans =
        5
        5
>> roots([\begin{array}{lll}{1}&{-17}&{72.5}\end{array}])
ans =
    8.5000 + 0.5000i
    8.5000-0.5000i
```

