

Object-Oriented Programming

Language



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5



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ECOOP 2010, Maribor, Slovenia

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- Ewan Tempero, The University of Auckland
- Dave Thomas, Bedarra Research Labs
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Design by a really small Committee

Supported by a wide community

Public Blog: http://www.gracelang.org

Obvious Questions:

What is an educational language? Why not use a "real" language? Why not Java? Scala? Python? Why now? Non-Questions: why start with objects? why teach objects at all?

What is an educational programming language?

Designed specifically for novices
 Can have limited or broad domain of application

We are interested in broad domain

Main focus is on programming in the small, but some modularity features.

Teach Industrial - Strength Languages?

Too much conceptual redundancy
High overhead for simple programs
Too hard to read and write
Conceptual clarity sacrificed for practicality
Saddled w/backward compatibility

Why Not Java?

Overloading Confusing subtyping with inheritance No user-defined operators Ø Primitive & Object types No lambdas

 Weak support for Generics

Covariant arrays

 Equality not automatic

 No definable control structures

Synchronized

Why Not Scala?

Too complex for novices
Multiple ways of doing everything
Weak generics
Powerful, but complex, type system

Why Not Python?

Weak encapsulation

Can't teach typed programming

Mismatch between method declarations & message sends

Implicit creation of fields

Why Now?

Happy teaching Java next 3-5 years In 2015, Java will be 20 years old State of the art has advanced - patches look like ... patches New languages bring good ideas ... but are for professionals, not students To be ready in 2015, we need to start now.

Our User Model

First year students in OO CS1 or CS2 objects early or late, static or dynamic types, functionals first or scriptings first or ... Second year students Faculty & TAs — assignments and libraries

We are in the dog food business

User model: Beginning students

Customer: experienced instructors



The consumer is not the customer

The Big Question

What do we hope the students learn?

- 1. To program well in Grace?
- 2. To understand and use the o-o model?
- 3. To be prepared for other languages and models?

My position: 3 is less important than 1 and 2

Features

Outputtered code; layout significant Structural typing Local type inference
 Subtyping separated from inheritance Ser-definable operators Sensible generics Lambdas

Allows both static and dynamic typing Parallel programming Equals & hashcode
 work automatically ø v instead of getV() for access Minimize "incantations" public static void main

Design is in early phases

Design is in early phasesAmbitious goals

Design is in early phases
Ambitious goals
Still disagree on many details

Grace Fundamentals

Severything is an object Simple method dispatch Single inheritance via cloning and concatenation Language levels for teaching
 Extensible via Libraries (control & data)
 Java / C / Python / Scala programmers should be able to read Grace programs

Might appear in a stream object

- except for methods
- Functions are objects
 - as in Smalltalk, lambda expressions create objects that mimic functions
- const welcomeAction := { print "Hello" }

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- const welcomeAction := { print "Hello" }

object { method apply { print "Hello" } }

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- const welcomeAction := { print "Hello" }

except for methods

Functions are objects

as in Smalltalk, lambda expressions create objects that mimic functions

const orderingFunction := $\{a, b \rightarrow a.name \leq b.name \}$

except for methods

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- @ except for methods
- Functions are objects
 - as in Smalltalk, lambda expressions create objects that mimic functions
- const orderingFunction := $\{a, b \rightarrow a.name \leq b.name \}$

object { method apply(a, b) { a.name ≤ b.name } }

except for methods

Functions are objects

as in Smalltalk, lambda expressions create objects that mimic functions

const orderingFunction := $\{a, b \rightarrow a.name \leq b.name \}$

Everything is an Object

except for methods

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Everything is an Object

except for methods

Functions are objects

as in Smalltalk, lambda expressions create objects that mimic functions

const orderingFunction := $\{a, b \rightarrow a.name \leq b.name \}$

if orderingFunction.apply(x, y) then { ... }

Everything is an Object

except for methods

Functions are objects

as in Smalltalk, lambda expressions create objects that mimic functions

const orderingFunction := $\{a, b \rightarrow a.name \leq b.name \}$

Everything is an Object But every object is not an instance of a class Instead: objects are self-contained Objects are created by executing an object constructor:

object { const x:Number := 2 const y:Number := 3 method distanceTo other:Point → Number { ((x - other.x)^2 + (y - other.y)^2) } } object { const x:Number := 2 const y:Number := 3 method distanceTo other:Point → Number { ((x - other.x)^2 + (y - other.y)^2) } }

object { const x:Number := 2 const y:Number := 3 method distanceTo other:Point → Number { ((x - other.x)^2 + (y - other.y)^2) } }

×	У	distance To	2	3

object { const x:Number := 2 const y:Number := 3 method distanceTo other:Point → Number { ((x - other.x)^2 + (y - other.y)^2) } }

Design Decisions:

- fields and methods share the same namespace
- ø p.x might be a field access or a method request
- the implementation can replace a field by a method without the client knowing

What about classes?

Pro

- Instructors are familiar with classes
- Classes capture a common pattern: a "factory" object that makes similar "instance" objects

Brevity

Con

- Unnecessary just
 use objects
- The common pattern usually lies in some way
- Restrictive, e.g.
 Smalltalk's parallel hierarchies

Compromise Design

Grace has classes; they resemble a block containing an object constructor

We try to make the syntax familiar, but not so familiar that we lie

Classes are restrictive, but the full power of object constructors is available to implement the general case

const Point := class { x': Number, y':Number →
 const x:Number := x'
 const y:Number := y'
 method distanceTo other:Point → Number {
 ((x - other.x)^2 + (y - other.y)^2) }
}

const Point := class { x': Number, y':Number →
 const x:Number := x'
 const y:Number := y'
 method distanceTo other:Point → Number {
 ((x - other.x)^2 + (y - other.y)^2) }
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const Point := class { x': Number, y':Number →
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 method distanceTo other:Point → Number {
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}



const Point := class { x': Number, y':Number →
 const x:Number := x'
 const y:Number := y'
 method distanceTo other:Point → Number {
 ((x - other.x)^2 + (y - other.y)^2) }
}

```
const Point = object {
   method new (x':Number, y':Number) {
      object {
        const x:Number := x'
        const y:Number := y'
        method distanceTo other:Point→Number {
            ((x - other.x)^2 + (y - other.y)^2) }}}
```

}

Class: Summary

```
const Point := class { x': Number, y':Number →
   const x:Number := x'
   const y:Number := y'
   method distanceTo other:Point → Number {
      ((x - other.x)^2 + (y - other.y)^2) }
}
```

const Point := object {
 method new (x':Number, y':Number) {
 object {
 const x:Number := x'
 const y:Number := y'
 method distanceTo other:Point→Number {
 ((x - other.x)^2 + (y - other.y)^2) }}}

One true message send

Schelter Like Smalltalk and Self:

- no overloading
- "method request" names the method and provides the arguments
- "dynamic dispatch" selects the correspondinglynamed method in the receiver
- "method execution" occurs in the receiver
- field access is via methods

One true message send

Schelter Like Smalltalk and Self:

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- "method execution" occurs in the receiver
- field access is via methods

(I'm trying to learn <u>not</u> to say "message-send" or "method call".)

Example: a Contact Object

}

const andrewInfo := object {
 var firstName := "Andrew"
 const lastName := "Black"
 method printOn s:Stream {
 s.puts firstName
 s.puts lastName }

Example: a Contact Object

}

const andrewInfo := object {
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 s.puts lastName }

Creates a method lastname

Example: a Contact Object

}

Creates 2 methods: firstName and firstName:=

Creates a method lastname const andrewInfo := object {
 var firstName := "Andrew"
 const lastName := "Black"
 method printOn s:Stream {
 s.puts firstName
 s.puts lastName }

Contact Object Expanded

const andrewInfo := object { privar ¿firstName? method firstName -> String { ¿firstName? } method firstName:= s:String { ¿firstName? := s } ¿firstname? := "Andrew" priconst ¿lastName? method lastName -> String { ¿lastName? } ¿lastname? := "Black" method printOn s:Stream { s.puts firstName s.puts s.puts lastName }

}

Contact Object Expanded

const andrewInfo := object { privar ¿firstName? method firstName -> String { ¿firstName? } method firstName:= s:String { ¿firstName? := s } ¿firstname? := 'Andrew" priconst ¿lastName? method lastName -> String { ¿lastName? } ¿lastname? := "Black" method printOn s:Stream Not s.puts firstName proposed for s.puts surface syntax s.puts lastName }

}

```
const contact := object {
  method named (first, last) -> Contact {
     object {
       var firstName:String := first
       var lastName:String := last
       method printOn s:Stream {
          s.puts firstName
          s.puts
          s.puts lastName } } }
  const database := MutableSequence.empty
  method add c:Contact {
     database addLast c }
```

```
const contact := object {
  method named (first, last) -> Contact {
     object {
       var firstName:String := first
       var lastName:String := last
       method printOn s:Stream {
          s.puts firstName
          s.puts
          s.puts lastName } } }
  const database := MutableSequence.empty
  method add)c:Contact {
     database addLast c }
```

attributes of the outer "factory" object

```
const contact := object {
  method named (first, last) -> Contact {
     object {
       var firstName:String := first
       var lastName:String := last
       method printOn s:Stream {
          s.puts firstName
          s.puts
          s.puts lastName } } }
  const database := MutableSequence.empty
  method add c:Contact {
     database addLast c }
```

const contact := object {
 method named (first, last) -> Contact {

object {

var firstName:String := first
var lastName:String := last
method printOn s:Stream {
 s.puts firstName
 s.puts ''
 s.puts lastName } }
const database := MutableSequence.empty
method add c:Contact {

database addLast c }

returns a contact object initialized to (first, last)

```
const contact := object {
  method named (first, last) -> Contact {
     object {
       var firstName:String := first
       var lastName:String := last
       method printOn s:Stream {
          s.puts firstName
          s.puts
          s.puts lastName } } }
  const database := MutableSequence.empty
  method add c:Contact {
     database addLast c }
```

Sample client code

const host := contact.named("Graham", "Hutton")
const guest := contact.named("Andrew", "Black")
contact.database.add host
contact.database.add guest

Inheritance

 Grace's
 inheritance story is based on an old idea of Taivalsaari

Delegation versus concatenation cloning is inheritance too

> Antero Taivalsaari University of Jyväskylä, Finland¹ tsaari@jyu.fi

In this paper a simple prototype-based model of object-oriented programming is introduced. Unlike previous prototype-based systems, which use delegation to achieve incremental modification of objects, the suggested model is based on concatenation: linear composition of object interfaces. The model eliminates the notions of delegation and parent slots from prototype-based programming, and shows that the essence of object-oriented programming can be captured using only a small number of user-level language constructs.

1. Introduction

Object-oriented systems are typically based on classes. Classes are descriptions of objects

capable of serving as templates from which instances, the actual objects described by classes can be instantiated. In class-based systems new kinds of other

trunted 1

Cloning + Concatenation = inheritance

ACM SIGPLAN OOPS Messenger Volume 6 Issue 3, July 1995

andrewInfo

first Name	last Name	last Name:=	printOn	"Andrew"	"Black"
---------------	--------------	----------------	---------	----------	---------

andrewInfo

first last Name Nam	Constraint and an	printOn	"Andrew"	"Black"
------------------------	---	---------	----------	---------

We want to add a telephone number. The "delta" is:

andrewInfo

first Name	last Name	last Name:=	printOn	"Andrew"	"Black"
---------------	--------------	----------------	---------	----------	---------

We want to add a telephone number. The "delta" is: andrewPhone

work	work	"+1 503
phone	phone:=	725 2411"

andrewInfo

first Name	last Name	last Name:=	printOn	"Andrew"	"Black"
---------------	--------------	----------------	---------	----------	---------

We want to add a telephone number. The "delta" is: andrewPhone

work	work	"+1 503
phone	phone:=	725 2411"

1. Clone both objects

andrewInfo

Cast	Look	Look			100 M
first Name	Contraction of the second s	last Name:=	printOn	"Andrew"	"Black"

We want to add a telephone number. The "delta" is: andrewPhone

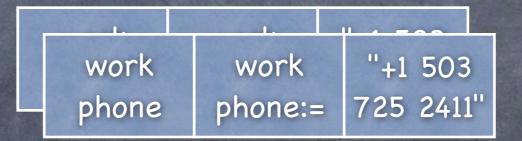


1. Clone both objects

andrewInfo

Current	Lash	Look			100 M
first Name	last Name	last Name:=	printOn	"Andrew"	"Black"

We want to add a telephone number. The "delta" is: andrewPhone



1. Clone both objects

2. Concatenate the copies

andrewInfo



andrewPhone

work	work	"+1 503
phone	phone:=	725 2411"

andrewPhone extends andrewInfo

work	work	"+1 503	first	last	last	printOn	"Andrew"	"Black"
phone	phone:=	725 2411"	Name	Name	Name:=			

In code:

const andrewInfo := object {
 var firstName := "Andrew"
 const lastName := "Black"
 method printOn s:Stream {
 s.puts firstName
 s.puts lastName }

}

}

const andrewPhone := object {
 var officePhone := "503 725 2411"

const and rewPhoneInfo := and rewPhone extends and rewInfo

No need to name intermediate objects:

const andrewPhoneInfo := object {
 var officePhone := "503 725 2411"
} extends contact.named ("Andrew", "Black")

Notice what this means:

Icone means shallow copy:

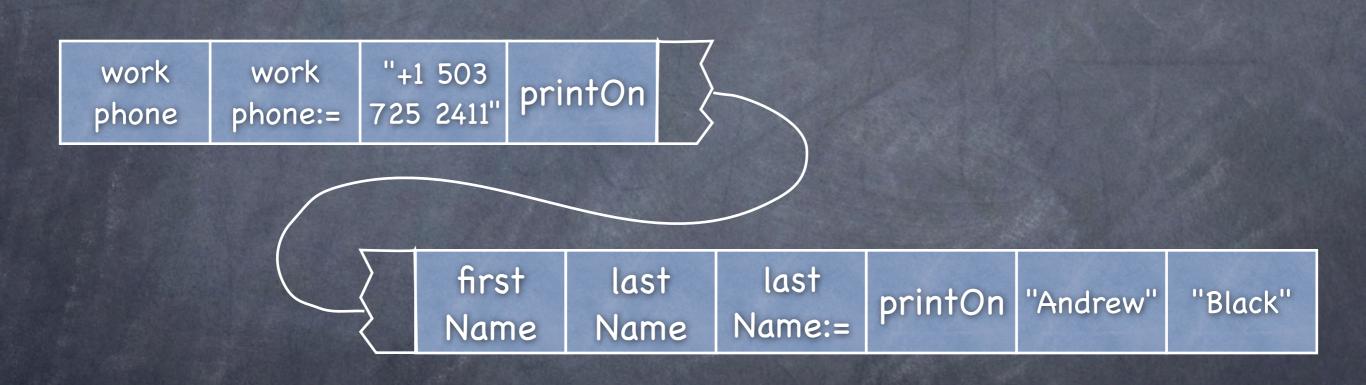
new object gets copies of the fields and the methods of the orignal objects

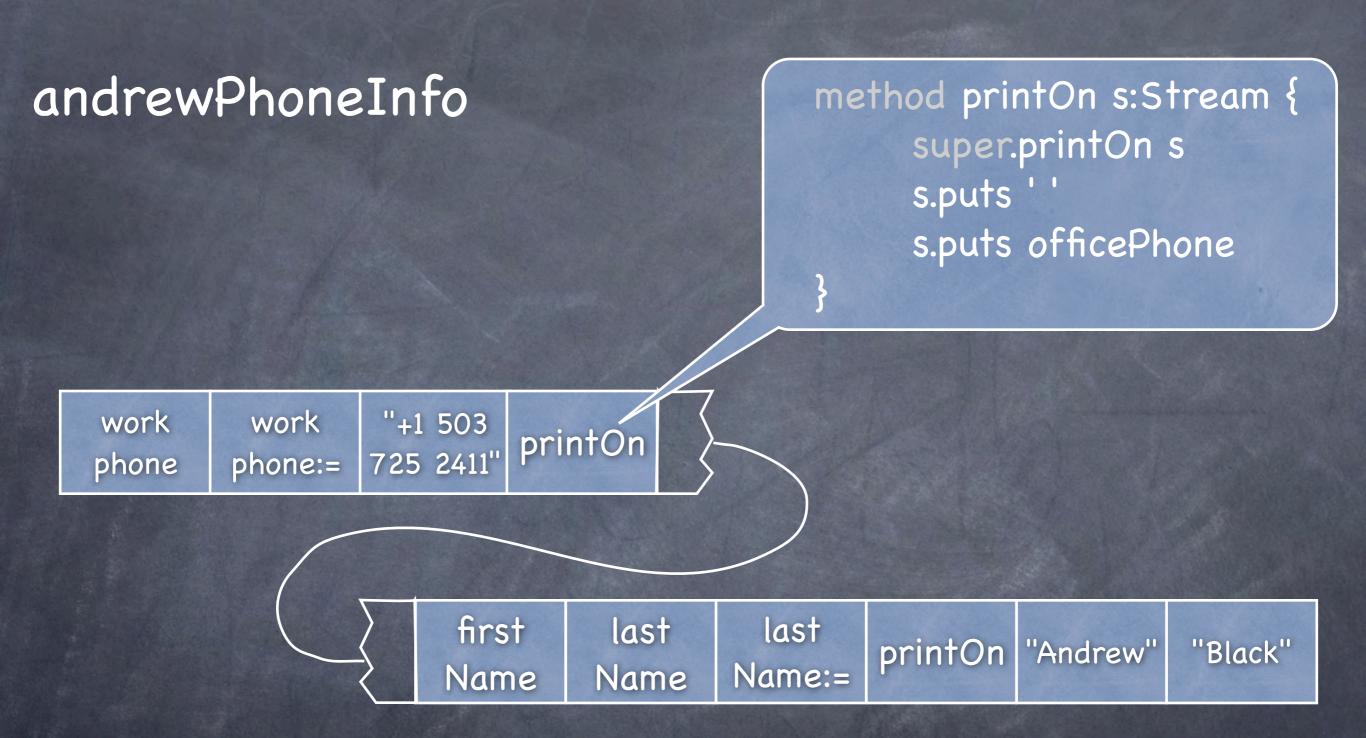
it's possible for an object to have two or more methods with the same name

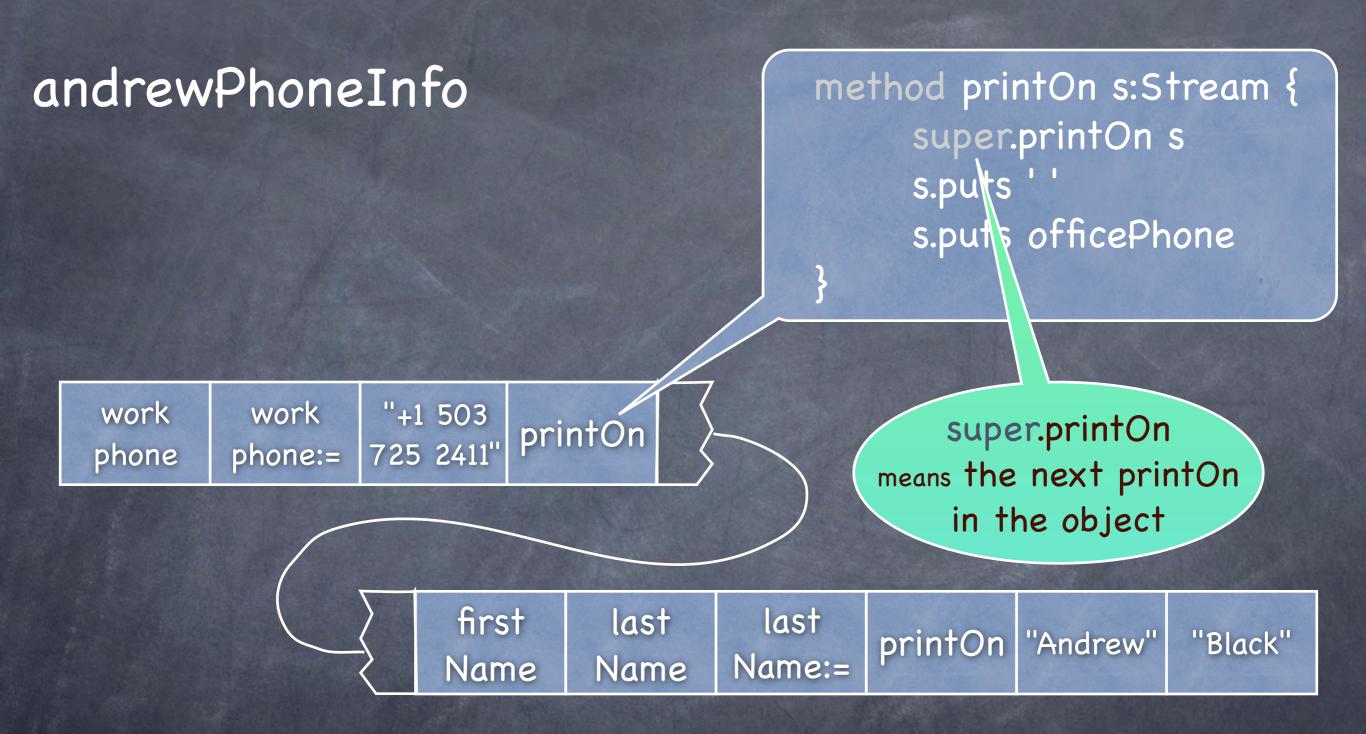
Let's fix printOn...

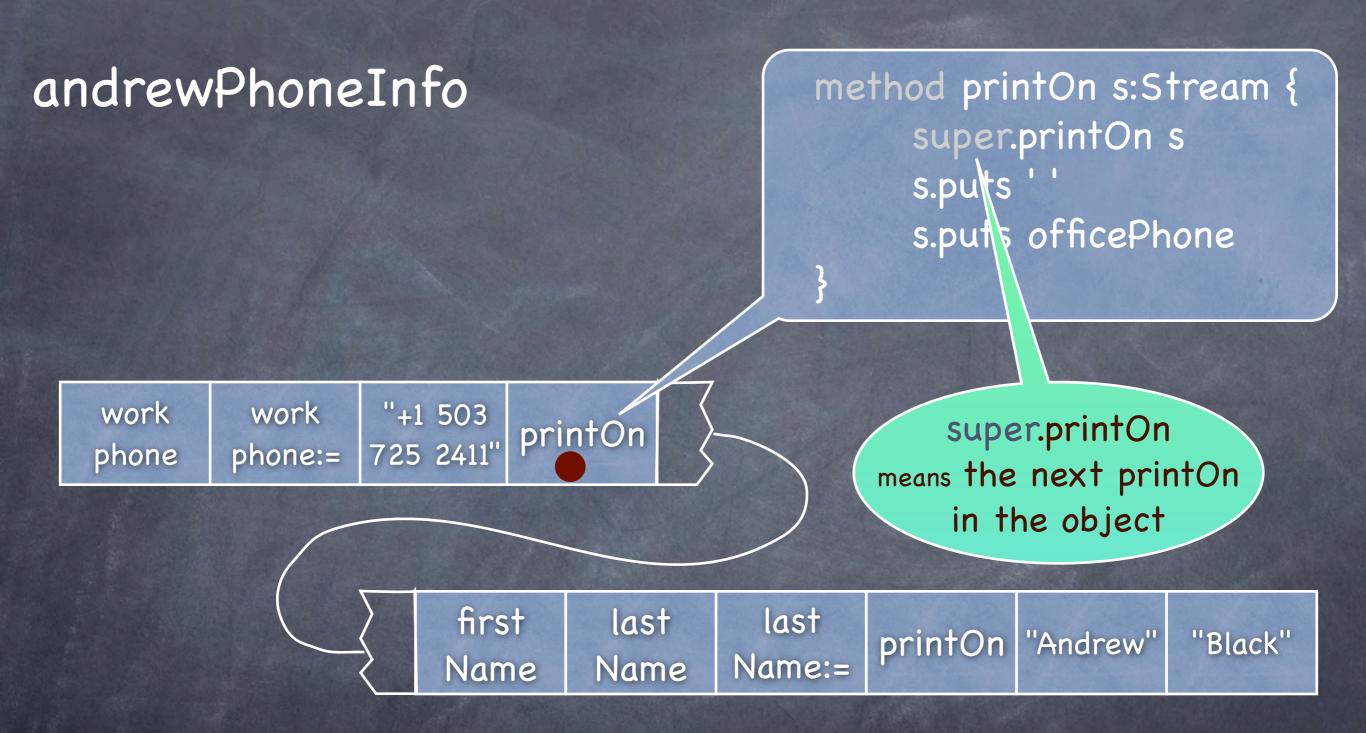
const andrewPhoneInfo := object {
 var officePhone := "503 725 2411"
 method printOn s:Stream {
 super.printOn s
 s.puts ''
 s.puts officePhone
 }
} extends contact.new ("Andrew", "Black")

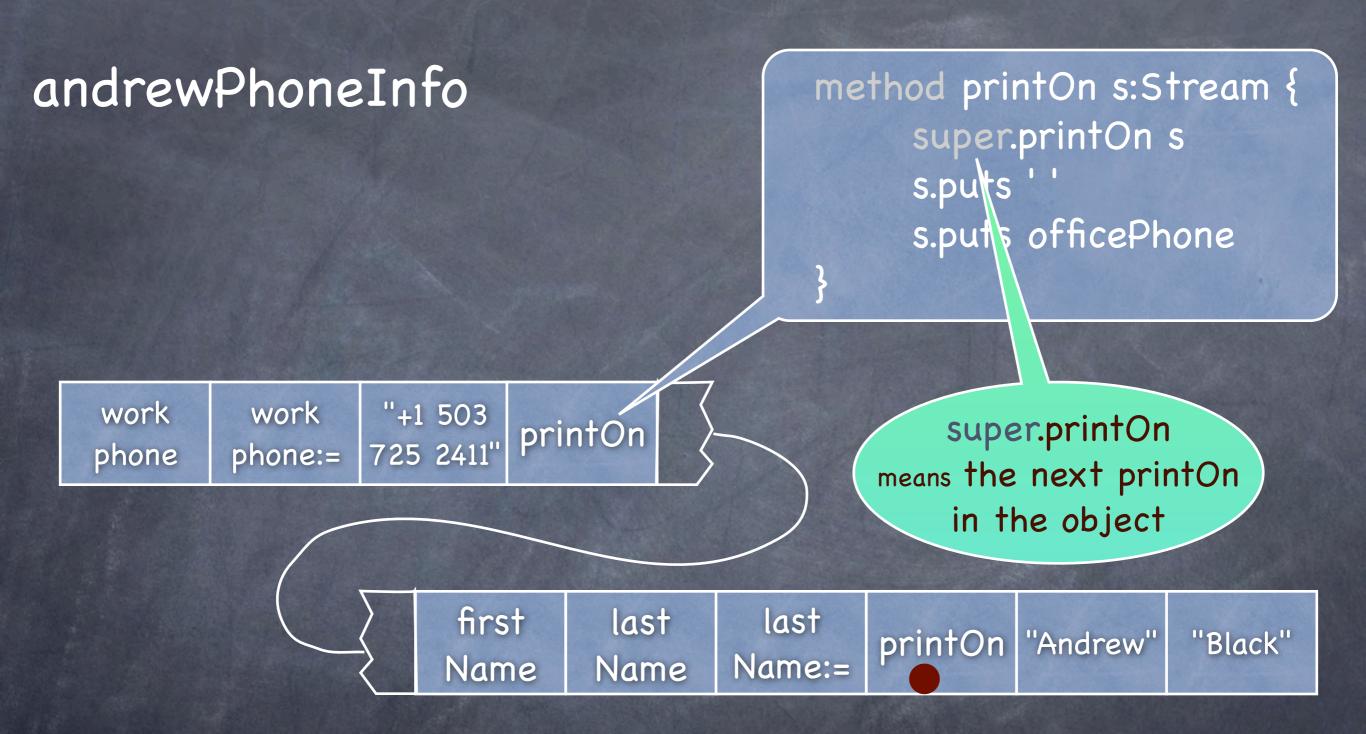
andrewPhoneInfo

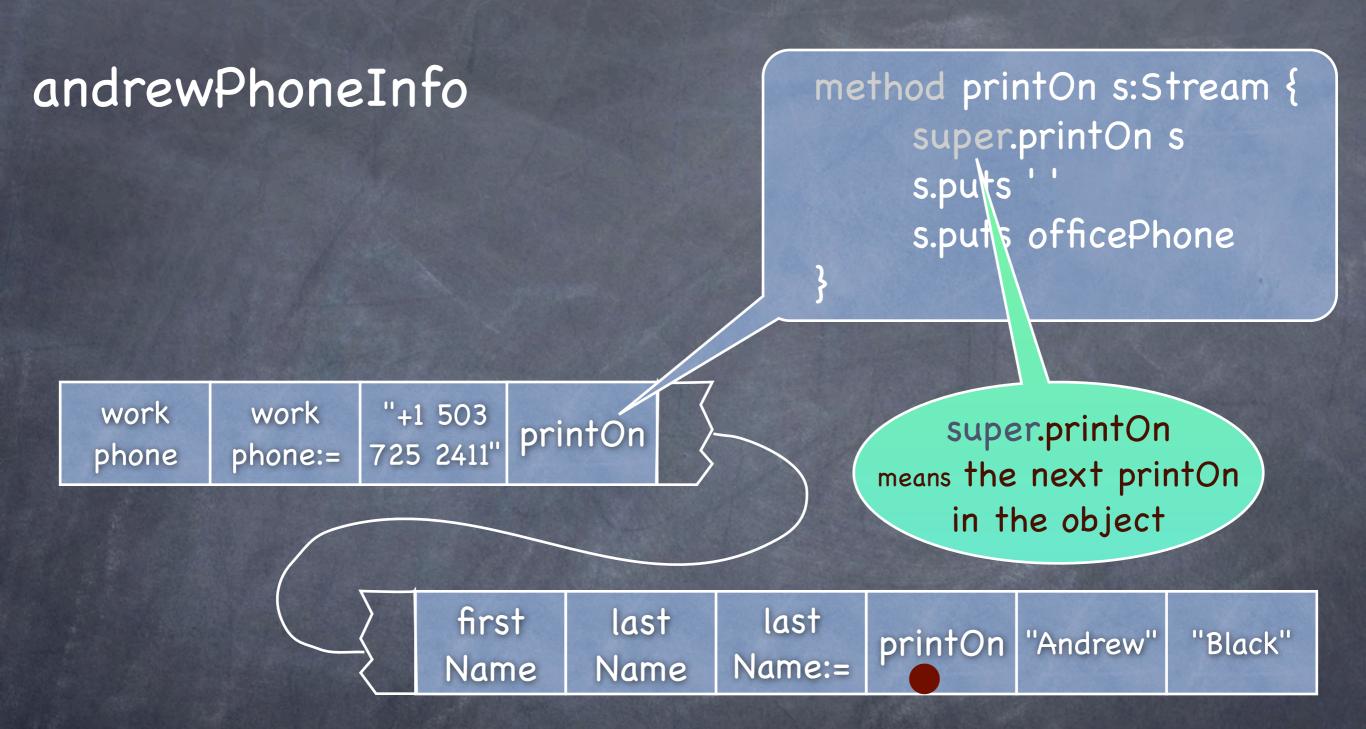












next might be a better keyword than super

Not your Grandfather's Inheritance

Not your Grandfather's Inheritance

What's more important:

To have a simple, explicable, inheritance story?

To have an inheritance story that's like the mainstream languages of the 1990's?

Not your Grandfather's Inheritance

What's more important:

To have a simple, explicable, inheritance story?

To have an inheritance story that's like the mainstream languages of the 1990's?

o Or:

Teachability, vs. familiar to instructors

Extensible via libraries

Objects include data and actions So it's essential to be able to define new control operations on new objects Second Example: a new kind of dictionary It must be possible to define new iterators, lookups, etc. with a syntax that's as convenient for the user as the "built in" objects

We achieve this!

Nothing is built in!

As in SELF, all built-in objects are really defined in libraries, including the Booleans.

 if <condition> then <block> else <block>, while <block> do <block>, and with <collection> do <block> are all method requests.

The methods are defined on object Grace, and inherited by all other objects object Grace := { method if c:Boolean then t:Block[→α] else f:Block[→α] → α { c ifTrue t ifFalse f }

}

object true := { method ifTrue t:Block[→α] ifFalse f:Block[→α] → α { t.apply } }

object false := { method ifTrue t:Block[→α] ifFalse f:Block[→α] → α { f.apply } } object Grace = { ... method with c:Collection[ɛ] do a:Block[ɛ→void] { c do a }

method with c:Collection[ε]
 map a:Block[ε→α] → Collection[α] {
 c collect a }

method with c: Collection[ε]
 select a:Block[ε→Boolean]
 → Collection[ε] {
 c.select a }

}

class interval = {
 const start:Number
 const stop:Number
 const step:Number

method do action:Block[Number→void] {
 var element
 var index := 0
 while {index < self size}
 do { element := start + (index × step)
 index := index + 1
 action.apply element } }</pre>

What about case?

Pro

 Instructors are familiar with case

Case is concise

 Students will meet case in other languages Con

Unnecessary — just use method dispatch
Assume "open classes"
Case violates object encapsulation
"Tell, don't ask"

What about case?

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Unnecessary — just use method dispatch
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Can we devise a simple, object-oriented dispatch?

What about case?

Pro

 Instructors are familiar with case

Case is concise

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Unnecessary — just use method dispatch
Assume "open classes"
Case violates object encapsulation
"Tell, don't ask"

Can we devise a simple, object-oriented dispatch? Should we? How can we teach case without case?
Add algebraic types and pattern matching?
Adopt Newspeak-style quad-dispatch case?
Scala case-classes?

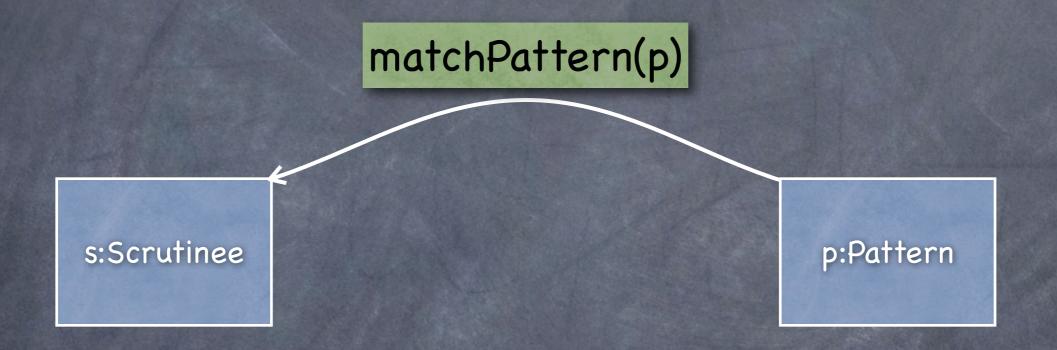
The last 2 men standing...

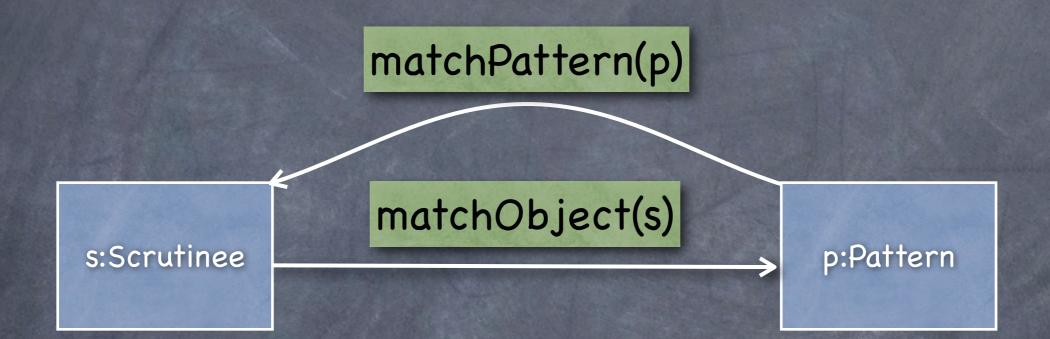
 Pattern-matching through method dispatch (James Noble, via Gilad Bracha)

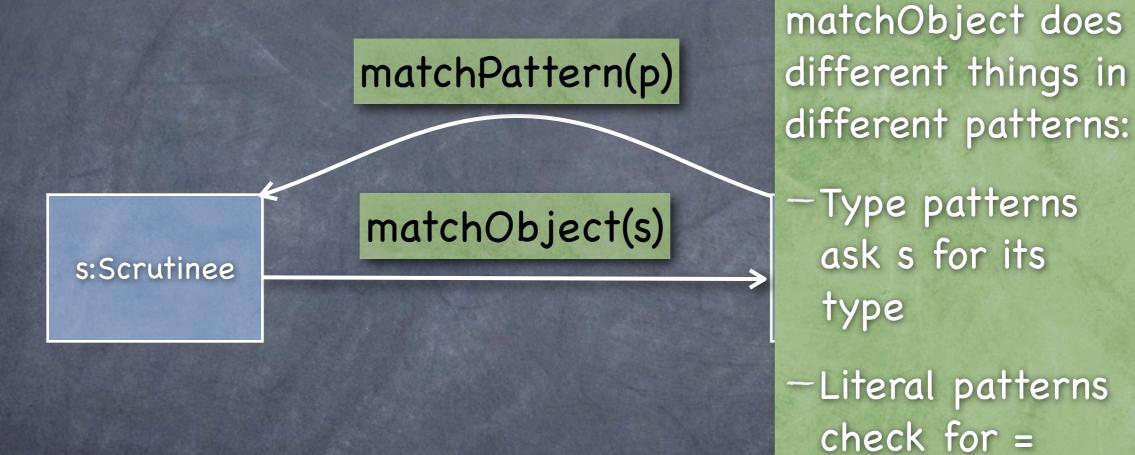
Scase as object (Andrew Black, via Blume, Acar & Chae)



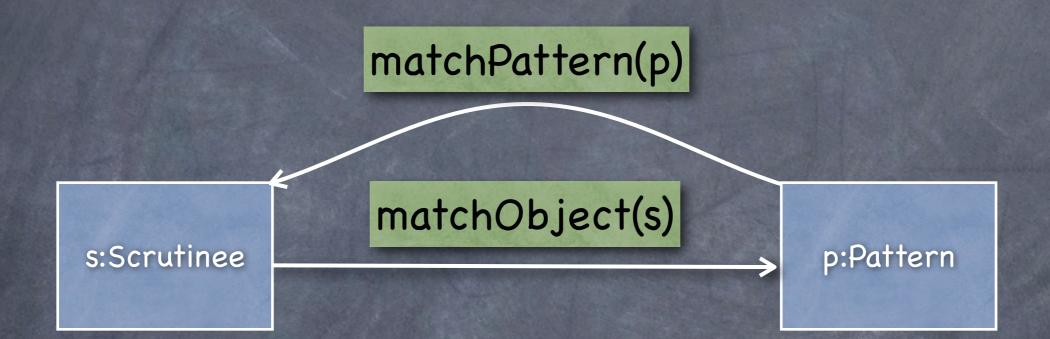


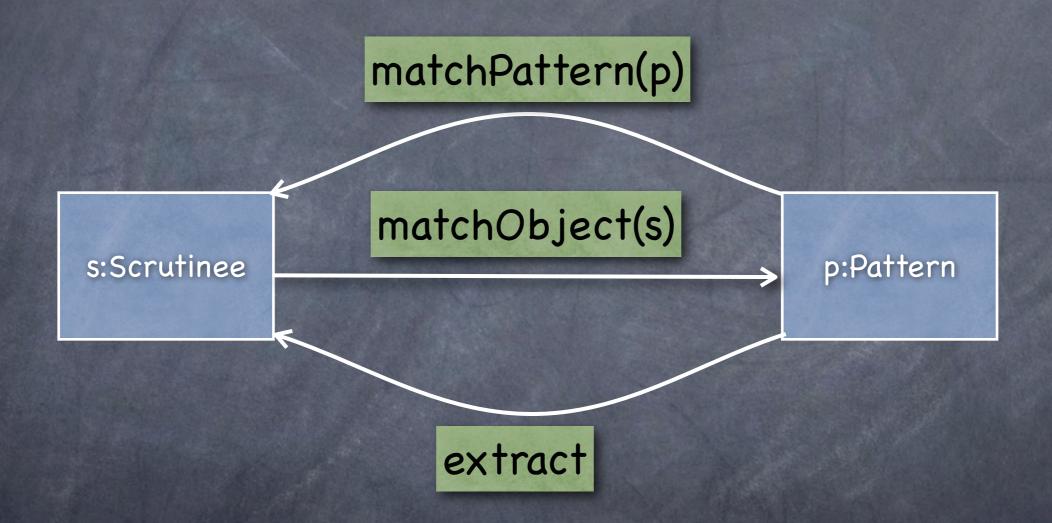


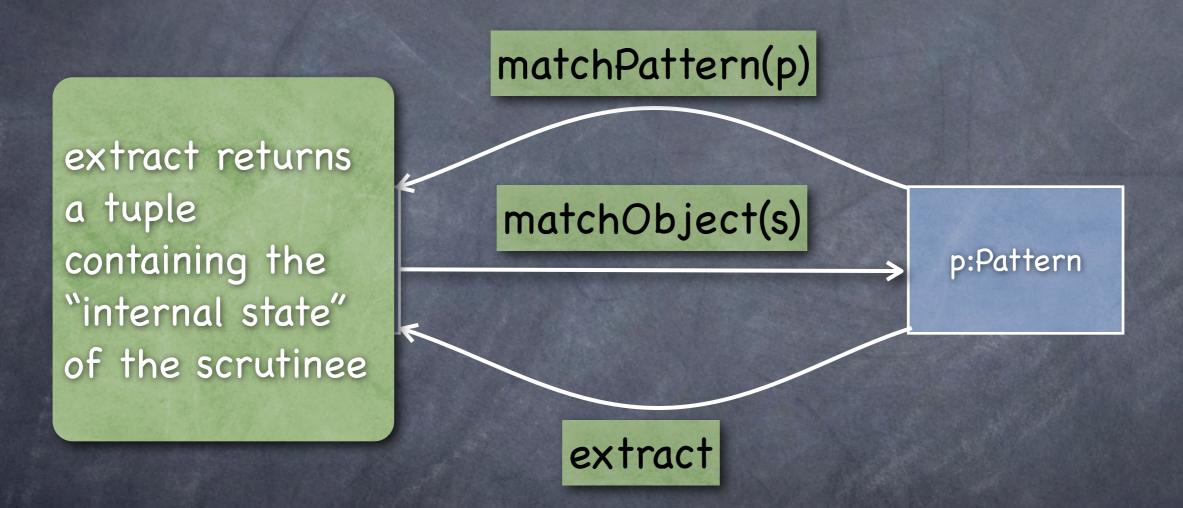


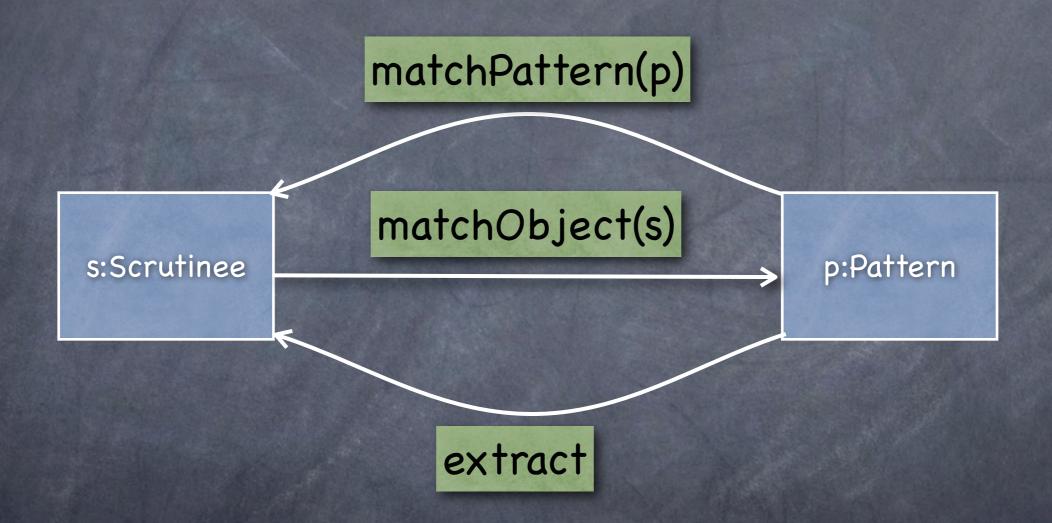


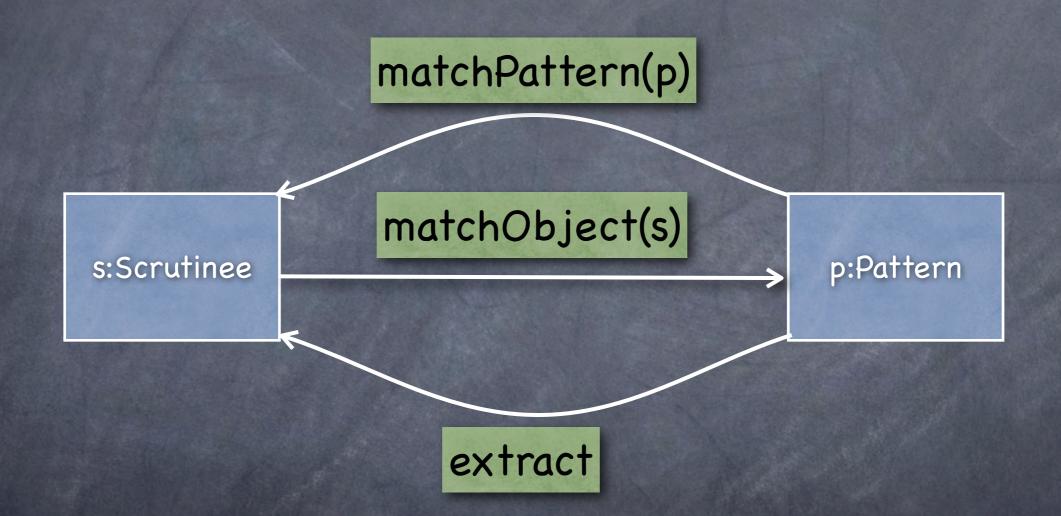
- etc











Treat any lambda-expression as a pattern

Case as Object

s:Scrutinee

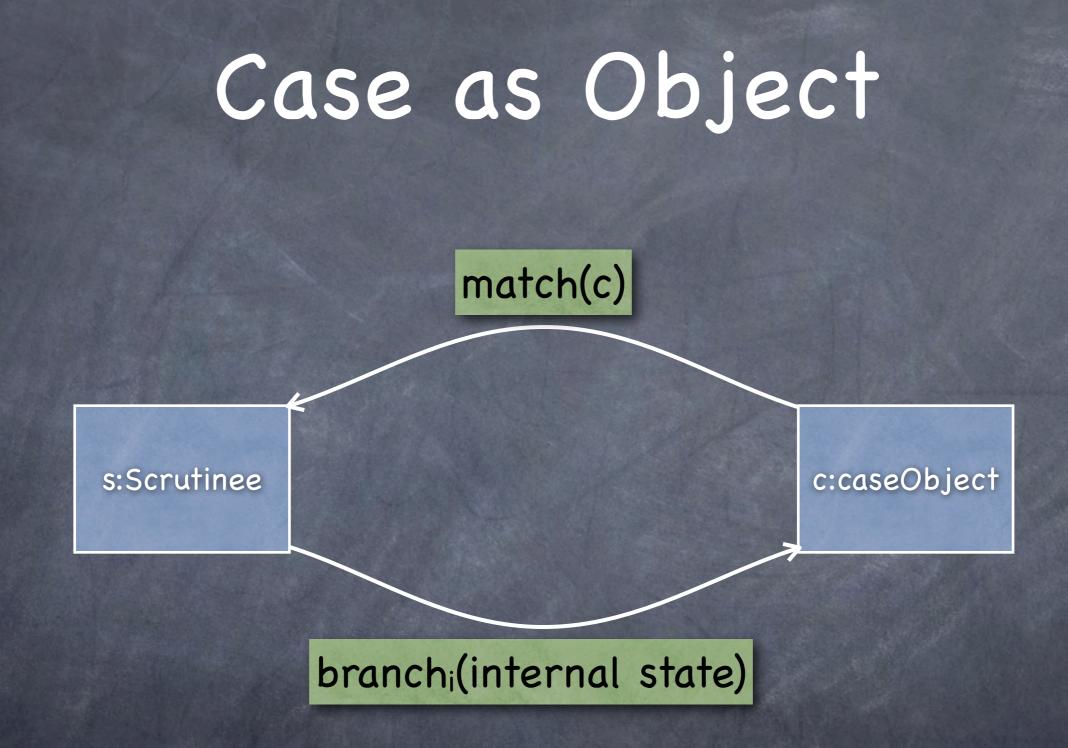
c:caseObject

Case as Object



s:Scrutinee

c:caseObject



Statements as well as Expressions?case statement?

What have we decided?

- Types are optional
- Lambdas-expressions are supported
- Sextensibility via libraries
- Types (= interfaces) are structural
- Classes define an interface corresponding to the operations on their instances
- Support for immutable objects
- Classes are open

Types are Optional

This means more than inferring type declarations:

"Untyped semantics": types don't change the semantics of correct programs — and not the syntax either!

explicit type annotations are assertions

just like assert s.notEmpty

The Laissez faire or George W. Bush interpretation:

If you mess up, the PDIC will bail you out.

The Laissez faire or George W. Bush interpretation:

If you mess up, the PDIC will bail you out.

Program debugger and interactive checker

The Laissez faire or George W. Bush interpretation:

If you mess up, the PDIC will bail you out.

The "The Nanny State" or Harold Wilson interpretation.

We will look after you. If it is even remotely possible that something may go wrong, we won't let you try.

The Laissez faire or George W. Bush interpretation:

The "The Nanny State" or Harold Wilson interpretation.

A third interpretation is useful:

The Laissez faire or George W. Bush interpretation:

The "The Nanny State" or Harold Wilson interpretation.

The "Proceed with caution", or Edward R. Murrow, interpretation.

The checker has been unable to prove that there are no type errors in your program. It may work; it may give you a run-time error. Good night and good luck.

Three interpretations

Onder all three interpretations, an error-free program has the same meaning.

Output the Wilson interpretation:

some error-free programs won't be permitted to run

an erroneous program will result in a checked run-time error.

Three interpretations

Onder all three interpretations, an error-free program has the same meaning.

Onder the Bush interpretation, all checks will be performed at runtime.

Even those that are guaranteed to fail because a counter-example is more useful than a type-error message

Three interpretations

- Onder all three interpretations, an error-free program has the same meaning.
- Onder the Bush interpretation, all checks will be performed at runtime.
- Onder the Murrow interpretation, you will get a mix of compile-time warnings and run-time checks.
- Onder the Wilson interpretation, you won't be permitted to run a program that <u>might</u> have a type-error

Help!

Supporters Programmers Implementers Library Writers
 IDE Writers Testers

- Teachers
- Students
- Tech Writers
- Textbook Authors
- Blog editors
- Community Builders

Schedule

- 2011: 0.1, 0.2 and 0.5 language releases, hopefully prototype implementations
- 2012 0.8 language spec, some mostly complete implementations
- 2013 0.9 language spec, reference implementation, experimental classroom use
- 2014 1.0 language spec, robust implementations, textbooks, initial adopters for CS1/CS2
- 2015 ready for general adoption?

No conclusions we aren't done yet

Questions

Comments Suggestions Brickbats