

# **An Introduction to Programming in Haskell**

Mark P Jones

Portland State University

# Haskell Resources:

# Haskell Resources:

- ◆ The focal point for information about Haskell programming, implementations, libraries, etc... is [www.haskell.org](http://www.haskell.org)
- ◆ I'll be using:
  - the Hugs interpreter ([haskell.org/hugs](http://haskell.org/hugs))
  - the Glasgow Haskell compiler, GHC, and interpreter, GHCi ([haskell.org/ghc](http://haskell.org/ghc))
- ◆ Online tutorials/references:
  - [learnyouahaskell.com](http://learnyouahaskell.com)
  - [book.realworldhaskell.org](http://book.realworldhaskell.org)

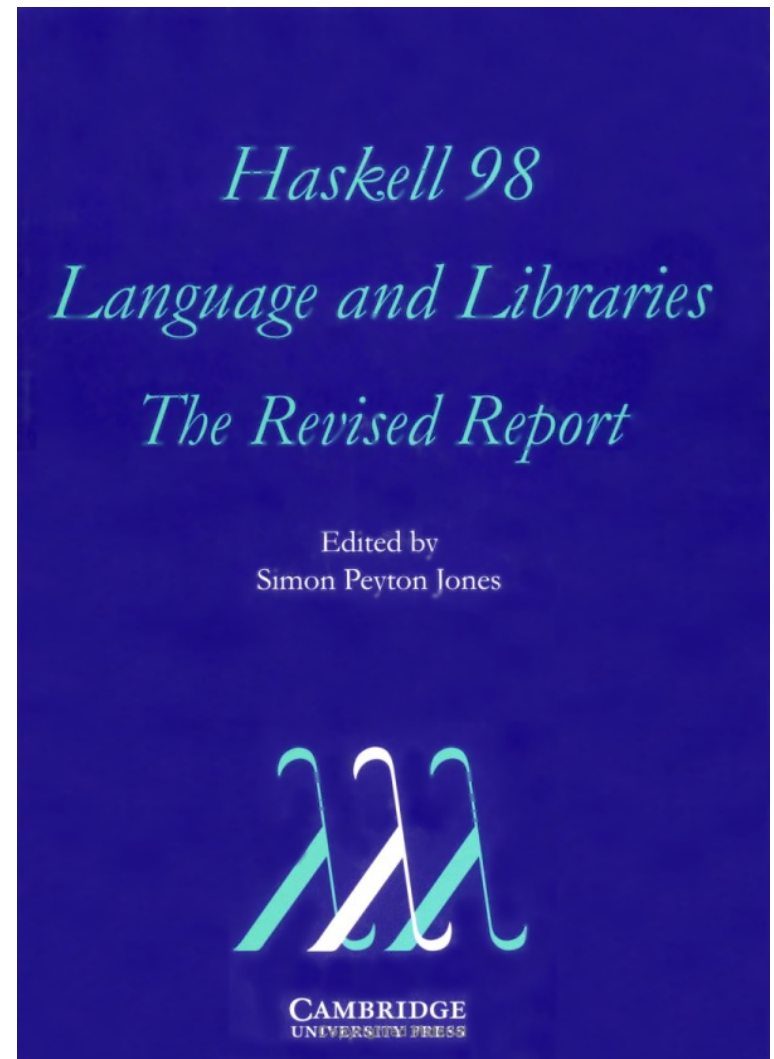
# The Language Report:

The definition of the Haskell 98 standard

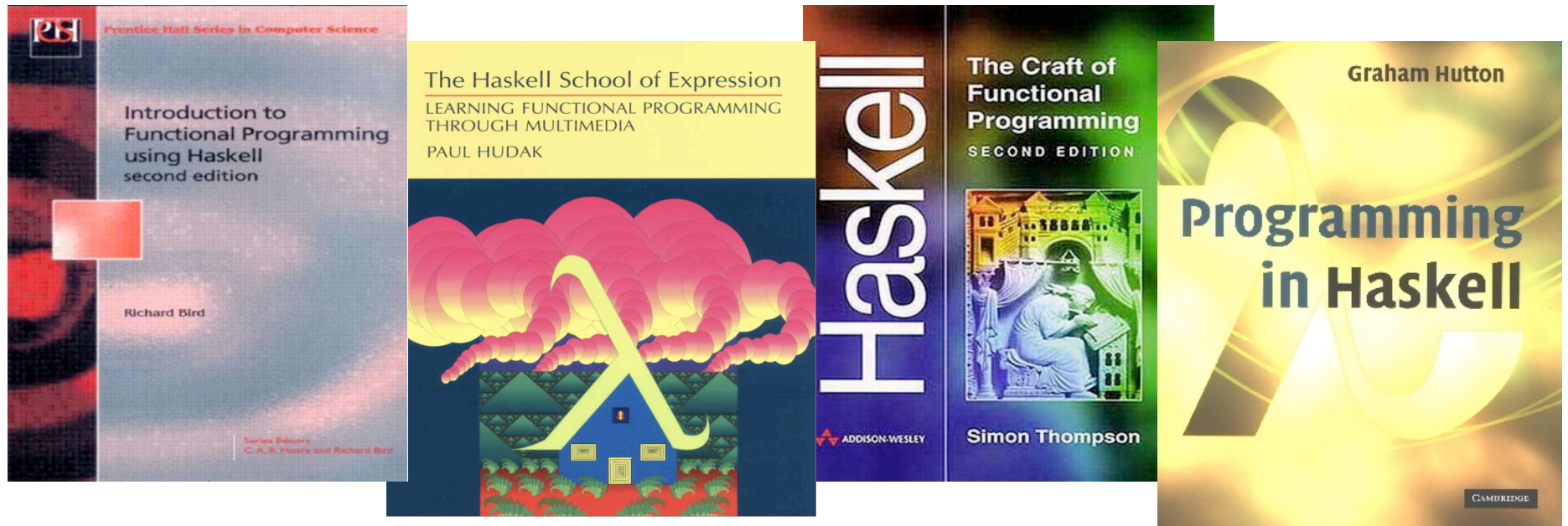
Lots of technical details ... not a great read!

Available in hard copy from Cambridge University Press

Or in pdf/html/etc... from [www.haskell.org/definition](http://www.haskell.org/definition)



# Textbooks:



- ◆ *Introduction to Functional Programming using Haskell* (2nd edition), Richard Bird
- ◆ *The Haskell School of Expression*, Paul Hudak
- ◆ *Haskell: The Craft of Functional Programming* (2nd edition), Simon Thompson
- ◆ *Programming in Haskell*, Graham Hutton

# What is Functional Programming?

# What is Functional Programming?

- ◆ Functional programming is a style of programming that emphasizes the evaluation of expressions, rather than execution of commands
- ◆ Expressions are formed by using functions to combine basic values
- ◆ A functional language is a language that supports and encourages programming in a functional style

# Functions:

In a pure functional language:

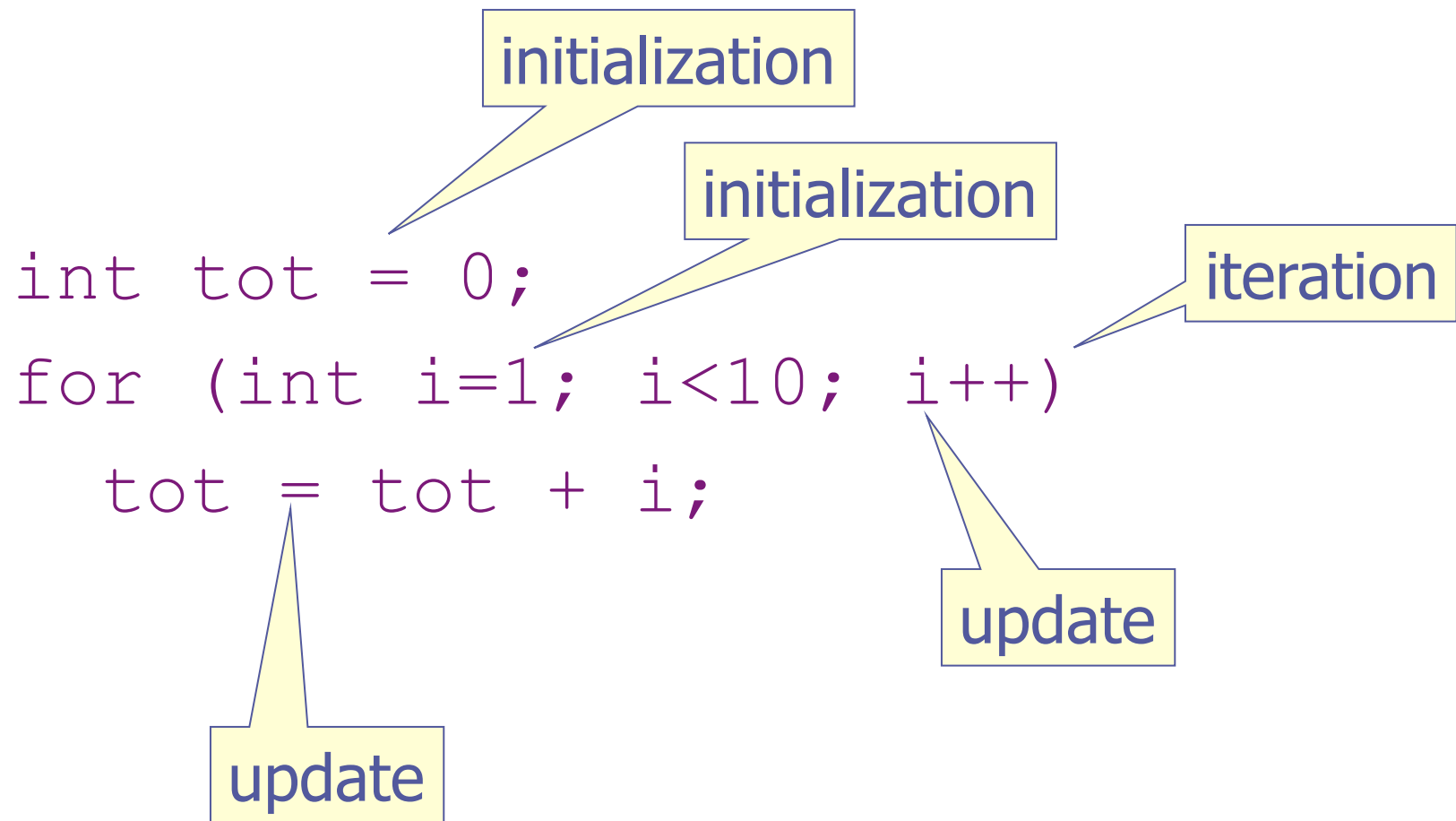
- ◆ The result of a function depends *only* on the values of its inputs:
  - Like functions in mathematics
  - No global variables / side-effects
- ◆ Functions are first-class values:
  - They can be stored in data structures
  - They can be passed as arguments or returned as results of other functions



# Example:

- ◆ Write a program to add up the numbers from 1 to 10

# In C, C++, Java, C#, ... :



implicit result returned in the variable `tot`

# In ML:

accumulating parameter

```
let fun sum i tot
    = if i > 10
      then tot
      else sum (i+1) (tot+i)
in sum 1 0
end
```

initialization

(tail) recursion

result is the value of this expression

# In Haskell:

```
sum [1..10]
```

combining  
function

the list of numbers to add

result is the value of this expression

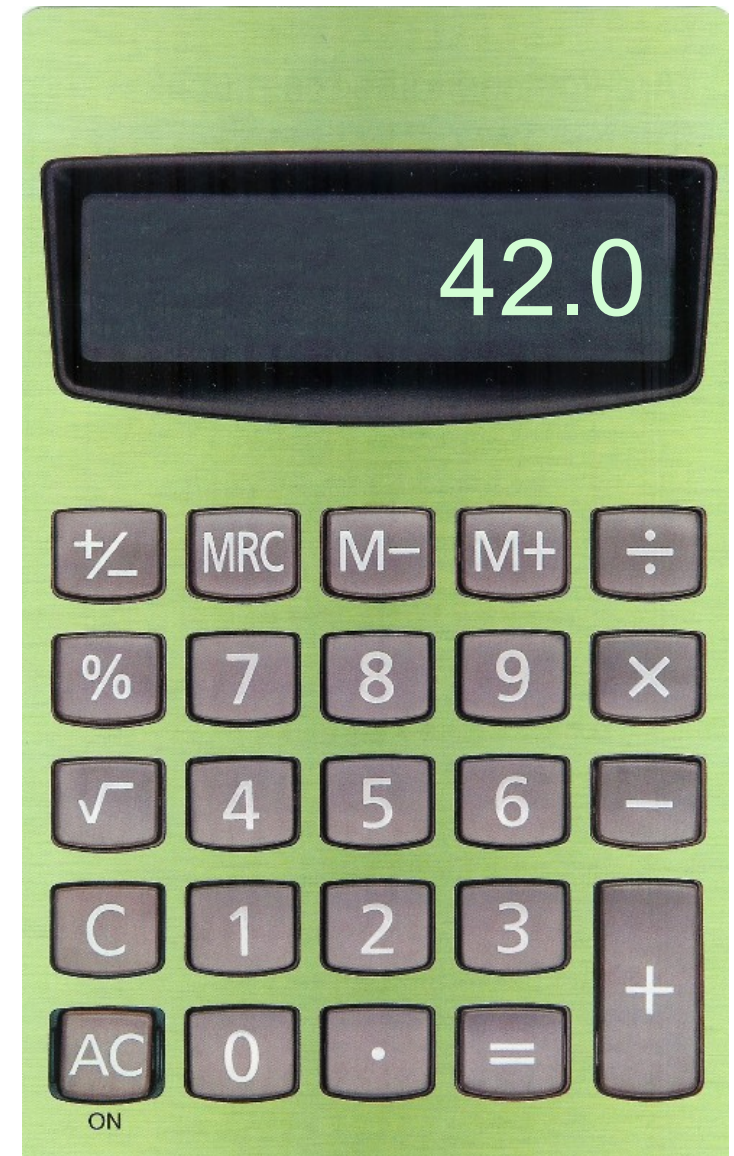
# Raising the Level of Abstraction:

"If you want to reduce [design time], you have to stop thinking about something you used to have to think about." (Joe Stoy, recently quoted on the Haskell mailing list)

- ◆ Example: memory allocation
- ◆ Example: data representation
- ◆ Example: order of evaluation
- ◆ Example: (restrictive) type annotations

# Computing by Calculating:

- ◆ Calculators are a great tool for manipulating numbers
- ◆ Buttons for:
  - entering digits
  - combining values
  - using stored values
- ◆ Not so good for manipulating large quantities of data
- ◆ Not good for manipulating other types of data



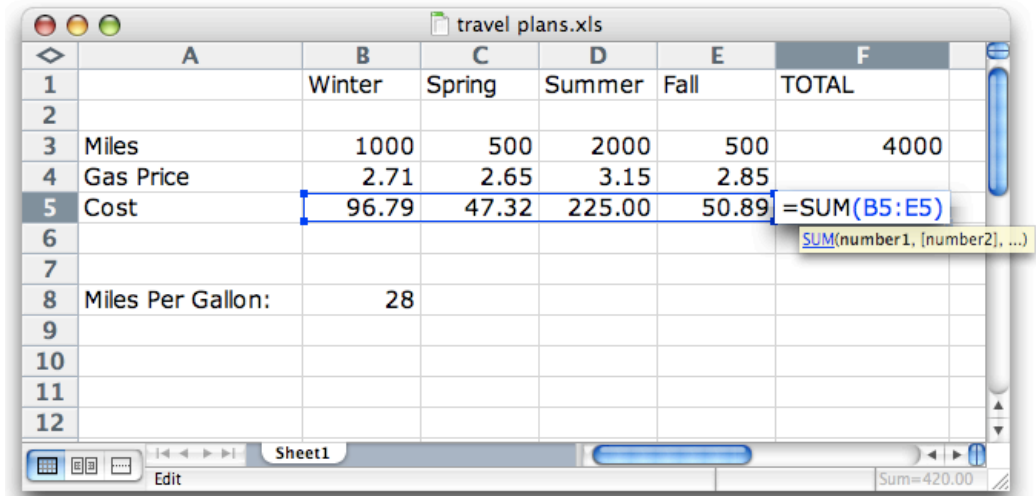
# Computing by Calculating:

- ◆ What if we could “calculate” with other types of value?
- ◆ Buttons for:
  - entering pixels
  - combining pictures
  - using stored pictures
- ◆ I wouldn't want to calculate a whole picture this way!
- ◆ I probably want to deal with *several different types of data at the same time*



# Computing by Calculating:

- ◆ Spreadsheets are better suited for dealing with larger quantities of data
- ◆ Values can be named (but not operations)
- ◆ Calculations (i.e., programs) are recorded so that they can be repeated, inspected, modified
- ◆ Good if data fits an “array”
- ◆ Not so good for multiple types of data



The screenshot shows a spreadsheet window titled "travel plans.xls". The data is organized as follows:

	A	B	C	D	E	F
1		Winter	Spring	Summer	Fall	TOTAL
2						
3	Miles	1000	500	2000	500	4000
4	Gas Price	2.71	2.65	3.15	2.85	
5	Cost	96.79	47.32	225.00	50.89	=SUM(B5:E5)
6						
7						
8	Miles Per Gallon:	28				
9						
10						
11						
12						

The formula bar at the bottom shows the formula for cell F5: `=SUM(B5:E5)`. A tooltip above the formula bar displays the syntax: `SUM(number1, [number2], ...)`. The status bar at the bottom right indicates "Sum=420.00".



# Functional Languages:

- ◆ Multiple types of data
  - Primitive types, lists, functions, ...
  - Flexible user defined types ...
- ◆ Operations for combining values to build new values (combinators)
- ◆ Ability to name values and operations (abstraction)
- ◆ Scale to arbitrary size and shape data
- ◆ “Algebra of programming” supports reasoning

# Getting Started with Haskell

# Starting Hugs:

```
user$ hugs
```

```
_____  
||  ||  ||  ||  ||  ||  ||  ||__  
||__||  ||__||  ||__||  __||  
||---||          __||  
||  ||  
||  || Version: September 2006 _____  
Hugs 98: Based on the Haskell 98 standard  
Copyright (c) 1994-2005  
World Wide Web: http://haskell.org/hugs  
Bugs: http://hackage.haskell.org/trac/hugs
```

```
Haskell 98 mode: Restart with command line option -98 to enable extensions
```

```
Type :? for help  
Hugs>
```

## The most important commands:

- `:q` quit
- `:l file` load file
- `:e file` edit file
- `Expr` evaluate expression

# The read-eval-print loop:

1. Enter expression at the prompt
2. Hit return
- 3. The expression is read, checked, and evaluated*
- 4. Result is displayed*
5. Repeat at Step 1

# Simple Expressions:

Expressions can be constructed using:

◆ The usual arithmetic operations:

$1 + 2 * 3$

◆ Comparisons:

$1 == 2$

$'a' < 'z'$

◆ Boolean operators:

$\text{True \&\& False}$

$\text{not False}$

◆ Built-in primitives:

$\text{odd } 2$

$\text{sin } 0.5$

◆ Parentheses:

$\text{odd } (2 + 1)$

$(1 + 2) * 3$

◆ Etc ...

# Expressions Have Types:

- ◆ The type of an expression tells you what kind of value you will get when you evaluate that expression:
- ◆ In Haskell, read “`::`” as “has type”
- ◆ Examples:
  - `1 :: Int`, `'a' :: Char`, `True :: Bool`, `1.2 :: Float`, ...
- ◆ You can even ask Hugs for the type of an expression: `:t expr`

# Type Errors:

```
Hugs> 'a' && True
ERROR - Type error in application
*** Expression      : 'a' && True
*** Term           : 'a'
*** Type           : Char
*** Does not match : Bool
```

```
Hugs> odd 1 + 2
ERROR - Cannot infer instance
*** Instance       : Num Bool
*** Expression    : odd 1 + 2
```

```
Hugs>
```

# Pairs:

- ◆ A pair packages two values into one

(1, 2)

('a', 'z')

(True, False)

- ◆ Components can have different types

(1, 'z')

('a', False)

(True, 2)

- ◆ The type of a pair whose first component is of type **A** and second component is of type **B** is written **(A,B)**

- ◆ What are the types of the pairs above?



# Operating on Pairs:

- ◆ There are built-in functions for extracting the first and second component of a pair:

`fst (True, 2) = True`

`snd (0, 7) = 7`

# Lists:

- ◆ Lists can be used to store zero or more elements, in sequence, in a single value:

`[]`   `[1, 2, 3]`   `['a', 'z']`   `[True, True, False]`

- ◆ All of the elements in a list must have the same type
- ◆ The type of a list whose elements are of type **A** is written as `[A]`
- ◆ What are the types of the lists above?

# Operating on Lists:

- ◆ There are built-in functions for extracting the head and the tail components of a list:
  - $\text{head } [1,2,3,4] = 1$
  - $\text{tail } [1,2,3,4] = [2,3,4]$
- ◆ Conversely, we can build a list from a given head and tail using the “cons” operator:
  - $1 : [2, 3, 4] = [1, 2, 3, 4]$

# More Operations on Lists:

◆ Finding the length of a list:

`length [1,2,3,4,5] = 5`

◆ Finding the sum of a list:

`sum [1,2,3,4,5] = 15`

◆ Finding the product of a list:

`product [1,2,3,4,5] = 120`

◆ Applying a function to the elements of a list:

`map odd [1,2,3,4] = [True, False, True, False]`

# Continued ...

◆ Selecting an element (by position):

`[1,2,3,4,5] !! 3 = 4`

◆ Taking an initial prefix (by number):

`take 3 [1,2,3,4,5] = [1,2,3]`

◆ Taking an initial prefix (by property):

`takeWhile odd [1,2,3,4,5] = [1]`

◆ Checking for an empty list:

`null [1,2,3,4,5] = False`

# More ways to Construct Lists:

## ◆ Concatenation:

$[1,2,3] ++ [4,5] = [1,2,3,4,5]$

## ◆ Arithmetic sequences:

$[1..10] = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]$

$[1,3..10] = [1, 3, 5, 7, 9]$

## ◆ Comprehensions:

$[ 2 * x \mid x <- [1,2,3,4,5] ] = [2, 4, 6, 8, 10]$

$[ y \mid y <- [1,2,3,4], \text{odd } y ] = [ 1, 3 ]$

# Strings are Lists:

◆ A String is just a list of Characters

```
['w', 'o', 'w', '!'] = "wow!"
```

```
['a'..'j'] = "abcdefghij"
```

```
"hello, world" !! 7 = 'w'
```

```
length "abcdef" = 6
```

```
"hello, " ++ "world" = "hello, world"
```

```
take 3 "functional" = "fun"
```

# Functions:

- ◆ The type of a function that maps values of type  $A$  to values of type  $B$  is written  $A \rightarrow B$
- ◆ Examples:
  - $\text{odd} :: \text{Int} \rightarrow \text{Bool}$
  - $\text{fst} :: (a, b) \rightarrow a$  ( $a, b$  are type variables)
  - $\text{length} :: [a] \rightarrow \text{Int}$



# Operations on Functions:

◆ Function Application. If  $f :: A \rightarrow B$  and  $x :: A$ , then  $f x :: B$

◆ Notice that function application associates more tightly than any infix operator:

$$f x + y = (f x) + y$$

◆ In types, arrows associate to the right:

$$A \rightarrow B \rightarrow C = A \rightarrow (B \rightarrow C)$$

Example:  $\text{take} :: \text{Int} \rightarrow [a] \rightarrow [a]$

$$\text{take } 2 [1,2,3,4] = (\text{take } 2) [1,2,3,4]$$

# Sections:

◆ If  $\oplus$  is a binary op of type  $A \rightarrow B \rightarrow C$ , then we can use “sections”:

- $(\oplus) \quad :: A \rightarrow B \rightarrow C$
- $(\text{expr } \oplus) :: B \rightarrow C$  (assuming  $\text{expr} :: A$ )
- $(\oplus \text{ expr}) :: A \rightarrow C$  (assuming  $\text{expr} :: B$ )

◆ Examples:

- $(1+)$ ,  $(2^*)$ ,  $(1/)$ ,  $(<10)$ , ...

# Higher-order Functions:

◆  $\text{map} :: (a \rightarrow b) \rightarrow [a] \rightarrow [b]$

■  $\text{map } (1+) [1..5] = [2,3,4,5,6]$

◆  $\text{takeWhile} :: (a \rightarrow \text{Bool}) \rightarrow [a] \rightarrow [a]$

■  $\text{takeWhile } (<5) [1..10] = [1,2,3,4]$

◆  $(.) :: (a \rightarrow b) \rightarrow (c \rightarrow a) \rightarrow c \rightarrow b$

■  $(\text{odd} . (1+)) 2 = \text{True}$

“composition”

# Definitions:

- ◆ So far, we've been focusing on expressions that we might want to evaluate.
- ◆ What if we wanted to:
  - Define a new constant (i.e., Give a name to the result of an expression)?
  - Define a new function?
  - Define a new type?
- ◆ Definitions are placed in files with a `.hs` suffix that can be loaded into the interpreter

# Simple Definitions:

Put the following text in a file "defs.hs":

```
greet name = "hello " ++ name
```

```
square x = x * x
```

```
fact n = product [1..n]
```

# Loading Defined Values:

Pass the filename as a command line argument to Hugs, or use the `:l` command from inside Hugs:

```
Main> :l defs
```

```
Main> greet "everybody"
```

```
"hello everybody"
```

```
Main> square 12
```

```
144
```

```
Main> fact 32
```

```
263130836933693530167218012160000000
```

```
Main>
```

# Using Libraries:

- ◆ Many useful functions are provided as part of the “Prelude”
- ◆ Many more are provided by libraries that must be imported before they can be used
- ◆ Example:

```
import Char
nextChar c = chr (1 + ord c)
```
- ◆ (The `Char` library also provides functions for converting to upper/lower case, testing for alphabetic or numeric chars, etc...)

# Typeful Programming:

- ◆ Types are an inescapable feature of programming in Haskell
  - Programs, definitions, and expressions that do not type check are not valid Haskell programs
  - Compilation of Haskell code depends on information that is obtained by type checking
- ◆ Haskell provides several predefined types:
  - Some built-in (functions, numeric types, ...)
  - Some defined in the Prelude (**Bool**, lists, ...)
- ◆ What if you need a type that isn't built-in?



# Type Synonyms:

# Type Synonym:

◆ A type synonym (or type abbreviation) gives a new name for an existing type.

◆ Examples:

**type** String = [Char]

**type** Length = Float

**type** Angle = Float

**type** Radius = Length

**type** Point = (Float, Float)

**type** Set a = a -> Bool

# Algebraic Datatypes:

# In Haskell Notation:

**data** Bool = False | True

introduces:

- A type, **Bool**
- A constructor function, **False :: Bool**
- A constructor function, **True :: Bool**

**data** List a = Nil | Cons a (List a)

introduces

- A type, **List t**, for each type **t**
- A constructor function, **Nil :: List a**
- A constructor function, **Cons :: a -> List a -> List a**

# More Enumerations:

```
data Rainbow = Red | Orange | Yellow  
              | Green | Blue | Indigo | Violet
```

introduces:

- A type, `Rainbow`
- A constructor function, `Red :: Rainbow`
- ...
- A constructor function, `Violet :: Rainbow`

# More Recursive Types:

```
data Shape      = Circle Radius  
                | Rect Length Length  
                | Transform Transform Shape
```

```
data Transform  
      = Translate Point  
      | Rotate Angle  
      | Compose Transform Transform
```

introduces:

- Two types, Shape and Transform
- Circle :: Radius -> Shape
- Rect :: Length -> Length -> Shape
- Transform :: Transform -> Shape -> Shape
- ...

# Using New Data Types:

- ◆ Building values of these new types is easy:

Nil :: List Rainbow

Cons Red Nil :: List Rainbow

Cons Blue (Cons Red Nil) :: List Rainbow

- ◆ But how do we inspect them or take them apart?

# Pattern Matching:

- ◆ In addition to introducing a new type and a collection of constructor functions, each data definition also adds the ability to pattern match over values of the new type

- ◆ Example:

```
first           :: (a, b) -> a
first (x, y)    = x
```

```
wavelengths    :: Rainbow -> (Length, Length)
wavelengths Red    = (620*nm, 750*nm)
wavelengths Orange = (590*nm, 620*nm)
```

...

```
nm = 1e-9 :: Float
```



# More Examples:

head :: [a] -> a  
head [] = error "head of []"  
head (x:xs) = x

length :: [a] -> Int  
length [] = 0  
length (x:xs) = 1 + length xs

area :: Shape -> Float  
area (Circle r) = pi \* r \* r  
area (Rect w h) = w \* h  
area (Transform t s) = area s

# Pattern Matching & Substitution:

- ◆ The result of a pattern match is either:
  - A failure
  - A success, accompanied by a substitution that provides a value for each of the values in the pattern
  
- ◆ Examples:
  - `[]` does not match the pattern `(x:xs)`
  - `[1,2,3]` matches the pattern `(x:xs)` with `x=1` and `xs=[2,3]`

# Patterns:

More formally, a pattern is either:

◆ An identifier

- Matches any value, binds result to the identifier

◆ An underscore (a “wildcard”)

- Matches any value, discards the result

◆ A constructed pattern of the form  $C p_1 \dots p_n$ , where  $C$  is a constructor of arity  $n$  and  $p_1, \dots, p_n$  are patterns of the appropriate type

- Matches any value of the form  $C e_1 \dots e_n$ , provided that each of the  $e_i$  values matches the corresponding  $p_i$  pattern.

# Other Pattern Forms:

For completeness:

## ◆ “Sugared” constructor patterns:

- Tuple patterns  $(p_1, p_2)$
- Cons patterns  $(ph : pt)$
- List patterns  $[p_1, p_2, p_3]$
- Strings, for example: `"hi" = ('h' : 'i' : [])`

## ◆ Character and numeric Literals:

- Can be considered as constructor patterns, but the implementation uses equality (`==`) to test for matches

# Function Definitions:

- ◆ In general, a function definition is written as a list of adjacent equations of the form:

$$f\ p_1 \ \dots\ p_n = rhs$$

where:

- $f$  is the name of the function that is being defined
  - $p_1, \dots, p_n$  are patterns, and  $rhs$  is an expression
- ◆ All equations in the definition of  $f$  must have the same number of arguments (the arity of  $f$ )

## ... continued:

- ◆ Given a function definition with  $m$  equations:

$$f\ p_{1,1} \ \dots \ p_{n,1} = \text{rhs}_1$$

$$f\ p_{1,2} \ \dots \ p_{n,2} = \text{rhs}_2$$

...

$$f\ p_{1,m} \ \dots \ p_{n,m} = \text{rhs}_m$$

- ◆ The value of  $f\ e_1 \ \dots \ e_n$  is  $S\ \text{rhs}_i$ , where  $i$  is the smallest integer such that the expressions  $e_j$  match the patterns  $p_{j,i}$  and  $S$  is the corresponding substitution.

# Example: filter

filter :: (a -> Bool) -> [a] -> [a]

filter p [] = []

filter p (x:xs)

| p x = x : rest

| otherwise = rest

where rest = filter p xs



guards



"where" clause

# Example: Binary Search Trees

**data** Tree = Leaf | Fork Tree Int Tree

insert :: Int -> Tree -> Tree

insert n Leaf = Fork Leaf n Leaf

insert n (Fork l m r)  
| n <= m = Fork (insert n l) m r  
| otherwise = Fork l m (insert n r)

lookup :: Int -> Tree -> Bool

lookup n Leaf = False

lookup n (Fork l m r)  
| n < m = lookup n l  
| n > m = lookup n r  
| otherwise = True



# Summary:

- ◆ An appealing, high-level approach to program construction in which independent aspects of program behavior are neatly separated
- ◆ It is possible to program in a similar compositional / calculational manner in other languages ...
- ◆ ... but it seems particularly natural in a functional language like Haskell ...

# Assignment #1

- ◆ Your goal is to write a function:
  - `toInt :: String -> Int`
- ◆ To accomplish this, consider the following functions:
  - `explode :: String -> [Char]`
  - `digitValue :: [Char] -> [Int]`
  - `reverse :: [Int] -> [Int]`
  - `pairedWithPowersOf10 :: [Int] -> [(Int,Int)]`
  - `pairwiseProduct :: [(Int,Int)] -> [Int]`
  - `sum :: [Int] -> Int`
- ◆ Write definitions for four of these functions (`reverse` and `sum` are built-in), using pattern matching and recursion where necessary
- ◆ Turn in an elegant program that communicates your solution well, including appropriate tests for each part.