1.(a) Yes, the grammar is ambiguous, because there are expressions with two distinct parse trees. For example, the expression \( x + y + z \) has these two parse trees:

\[
\begin{array}{c}
\text{expr} \\
\text{term} \\
\text{id} + \text{term} \\
\text{id} \text{ term} + \text{term}
\end{array}
\quad
\begin{array}{c}
\text{expr} \\
\text{term} \\
\text{term} + \text{term} \\
\text{id} \text{ id} + \text{term}
\end{array}
\]

(b) No such program can exist. The only way evaluation order can affect the result is because of side-effects, and the only side-effect in the language is assignment. The operands of \(' + ' and \(' * ' are always terms; evaluating terms never causes an assignment. So it can't matter which order the operands are evaluated.

2.(a,b,c) Marked below. For (e), the idea is that to achieve dynamic scope, we simply set the initial environment for evaluating functions to the caller's environment (rather than the lexical environment of the callee, which for this language is empty).

(c) Under static scoping, the use of \( y \) in the body of \( g \) is free, so interpreting the program raises the exception \texttt{InterpException("undefined variable:y")}.

(d) Under dynamic scoping, the binding of \( y \) is visible inside \( g \), so the program evaluates to \( 21 + 21 = 42 \).
3. (a) The result of \( h \) is the result of \( g \), which is the value of \( b \). But \( b \) was never initialized and its values is therefore undefined.

(b) In the generated code, the stack location of \( a \) in \( f \) is the same as the location of local variable \( b \) in \( g \). The call to \( f \) sets this location to 42; the code in \( g \) reads this location and uses it as the value for \( b \), which is then returned. (Note: Without knowing what assembly code \( gcc \) generates, you can’t be sure this is why the program behaves as it does, but it’s the most probable explanation.)

(c) A Java compiler would reject the program (during semantic analysis), because the use of \( b \) before it has been defined violates Java’s “definite assignment” property.

(d) We would need to change the declaration of \( int b \) to \( var b : \text{Int} = \ldots \) or \( val b : \text{Int} = \ldots \), for some expression \( \ldots \), because Scala does not allow uninitialized local variables.

4. OK, this problem is pretty hard! The key idea is that two pairs are reference-equal iff a change to an element in one pair is visible in the other. So our function should use \( \text{setFst} \) (or \( \text{setSnd} \)) to alter one pair and \( \text{fst} \) (or \( \text{snd} \)) to see if that has also changed the other pair. The rest is just details: we must do the update with a value that is different from the existing element value (in both pairs), and we must reset the original element value after completing the test (no matter how it comes out).

Here’s one concrete implementation:

\[
\text{eqpair \ (a \ b)}
\begin{align*}
\text{let \ q \ (fst \ a)}
\quad \text{if \ (eq \ q \ (fst \ b))}
\quad \{ \ (fst \ a) = (fst \ b) = q \}
\quad \text{(let \ p \ (if \ (eq \ q \ 42) 43 42)} \ (\text{any \ p \neq \ q \ will \ do})
\quad \text{(block)}
\quad \ (\text{setFst} \ a \ p)
\quad \ (\text{let \ r \ (eq \ (fst \ b) \ p)} \ (\text{set \ r \ = \ true \ iff \ changing \ a \ also \ changed \ b})
\quad \text{(block)}
\quad \ (\text{setFst} \ a \ q) \ (\text{undo \ the \ change \ on \ a \ (perhaps \ also \ on \ b!}) \ (\text{r}))
\end{align*}
\]
\[
\{ \ (fst \ a) \neq (fst \ b), \text{so pairs cannot be the same} \}
\]

5. (a)
\[
\text{while} \ (\neg \ E) \ \text{do} \ S
\]

(b)
\[
\frac{\{P\} S \{Q\} \ {Q \land \neg E} S \{Q\}}{\{P\} \ \text{repeat} \ S \ \text{until} \ E \ \{Q \land E\}} \ \text{(REPEAT-UNTIL)}
\]

6. \[
\{e_1, E, S\} \Downarrow \langle v_1, S' \rangle \quad \{e_2, E, S'\} \Downarrow \langle v_2, S'' \rangle \quad v_1 \neq 0 \quad \langle e_3, E, S'' \rangle \Downarrow \langle v_3, S''' \rangle
\]
\[
\{ (\text{gifnz} e_1 e_3 e_2 e_4), E, S \} \Downarrow \langle v_4, S''' \rangle \ \text{(Gif-1)}
\]
\[
\{e_1, E, S\} \Downarrow \langle v_1, S' \rangle \quad \{e_2, E, S'\} \Downarrow \langle v_2, S'' \rangle \quad v_2 \neq 0 \quad \langle e_3, E, S'' \rangle \Downarrow \langle v_3, S''' \rangle
\]
\[
\{ (\text{gifnz} e_1 e_3 e_2 e_4), E, S \} \Downarrow \langle v_4, S''' \rangle \ \text{(Gif-2)}
\]