## CS558 Programming Languages - Fall 2023 - Suggested Study Question Solutions for Lecture 3a

1. (a) Binding of identifiers occurs for a at line $01, g$ and $c$ at line 02 , a at line $03, d$ at line 04 (but not line $05!$ ), $h$ and $e$ at line $06, f$ and $b$ at line $09, j$ at line 11 , and $b$ at line 12 .
(b) Uses of identifiers occur for c at line 03 , f and a at line $04, \mathrm{~d}$ (twice!) at line 05 , e and d at line $07, \mathrm{~h}$ at line 08 , a and b at line $10, \mathrm{~g}$ and a at line 12 , and b and a at line 13 .
(c)
```
01: {}
02: {a: line 01}
03: {a: line 01, c: line 02, g: line 02, f: line 09 }
04: {a : line 03, c: line 02, g: line 02, f: line 09 }
05: {a : line 03, c: line 02, d: line 04, g: line 02, f: line 09}
06: {a : line 03, c: line 02, d: line 04, g: line 02, f: line 09}
07: {a : line 03, c: line 02, d: line 04, e: line 06, g: line 02, f: line 09, h: line 06 }
08: {a : line 03, c: line 02, d: line 04, g: line 02, f: line 09, h: line 06 }
09: {a: line 01}
10: {a: line 01 b: line 09, g: line 02, f: line 09 }
11: {a: line 01 g: line 02, f: line 09 }
12: {a: line 01 g: line 02, f: line 09, j: line 11}
13: {a line 01 b: line 12, g: line 02, f: line 09, j: line 11}
14: {a: line 01 g: line 02, f: line 09, j: line 11}
```

(d) The free identifiers of let $b=g(a)$ in $b+a$ are $a$ and $g$. When we ask for the free identifiers "of a function" we exclude the the parameters and the function name itself (assuming it is allowed to be recursive), so the sole free identifier of f is a. (If we asked for the "free identifiers of the body of f ," i.e. of the expression $a+b$, the answer would be $\{a, b\}$.) The entire expression has no free variables; such expressions are said to be "closed."
2. With static scoping, the uses of $x$ in set_x and print_x always resolve to the global $x$, so we see 112 2. With dynamic scoping, the calls to set_x and print_x from inside second see the local $x$ instead, so we see $1 \quad 1 \quad 21$.
3. Here's one simple solution. Function $g$ uses a variable (a) before it is initialized if-and-only-if control ever reaches the third line of the function, and that occurs if-and-only-if $f$ returns.

```
void g() {
    int a;
    f();
    int b = a; // uninitialized use, but do we ever get here?
}
```

A similar line of reasoning (not always quite so trivial!) can be used to show that essentially every interesting property of programs is undecidable in general. (Google "Rice's theorem" for a theoretical exposition.) The consequence is that compilers must necessarily approximate their analyses of program behavior.

