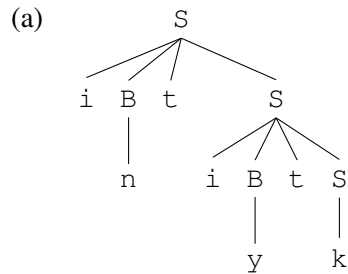


## CS558 Programming Languages – Fall 2023 – Suggested Study Question Solutions for Lecture 1b

1. This grammar is inspired by the well-known “dangling else” ambiguity problem, which arises in languages where the `else` clause of an `if-then-else` statement is optional and the `then` and `else` bodies can be arbitrary statements (including `if-then-else` statements!) For example, in the sentence of part (c), the ambiguity of the grammar means that we cannot tell which `i(f)` `the e(lse)` should be paired up with.



(b) Here are two possible derivations; the first is the left-most derivation for the tree in (a); the second is the right-most derivation.

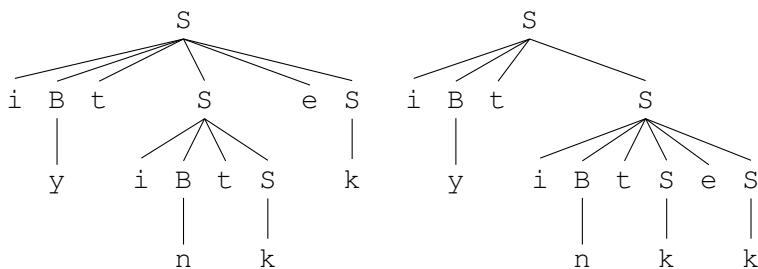
```

S → i B t S
  → i n t S
  → i n t i B t S
  → i n t i y t S
  → i n t i y t k
  
```

```

S → i B t S
  → i B t i B t S
  → i B t i B t k
  → i B t i y t k
  → i n t i y t k
  
```

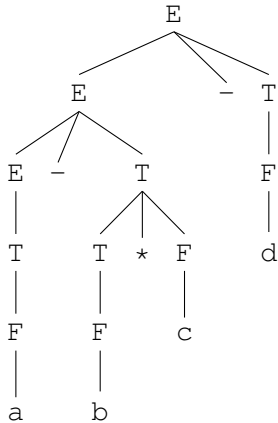
(c) Here are two trees:



(d) The grammar is ambiguous because there is at least one sentence—in particular, the one in part (c)—that has two different parse trees.

Note that it is *not* correct to claim that the grammar is ambiguous because the same sentence has more than one linear derivation sequence. Even when a sentence has a unique parse tree, as in part (a), it may have multiple corresponding linear derivation sequences, as in part (b). Such a sentence is not enough to demonstrate ambiguity: we must find a sentence that has two different *trees*.

2. (a) Here's the parse tree:



(b) We simply switch the order of the nonterminals in the right-hand side of the production for E:

$E ::= T + E \mid T - E \mid T$

The rest of the grammar remains the same. Try it on an example like  $a + b + c$ .

3. Here is a possible AST tree grammar:

*Add* :  $Exp \rightarrow Exp Exp$

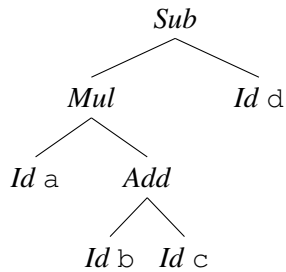
*Sub* :  $Exp \rightarrow Exp Exp$

*Mul* :  $Exp \rightarrow Exp Exp$

*Div* :  $Exp \rightarrow Exp Exp$

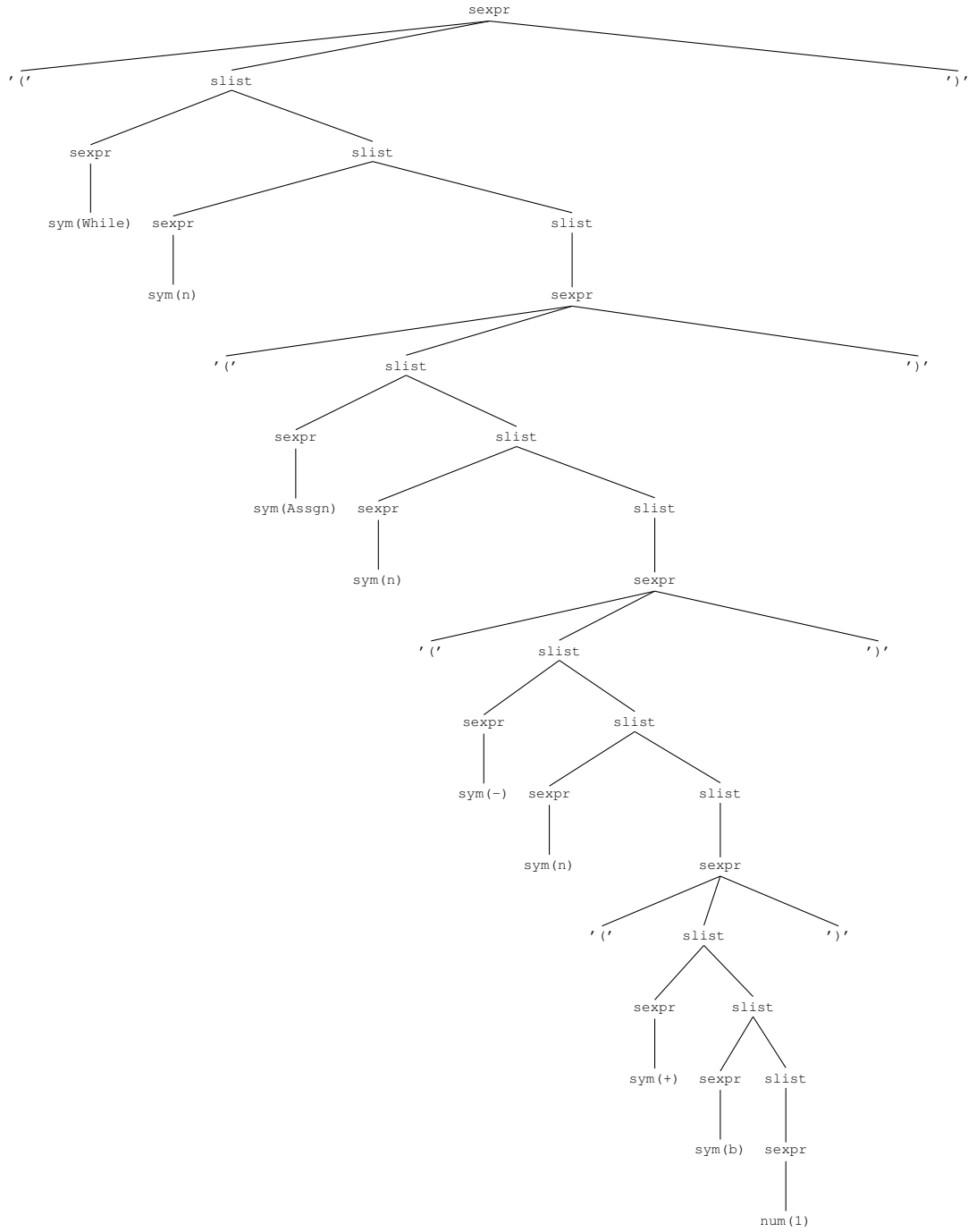
*Id* :  $Exp \rightarrow (string)$

(b)



The most important differences from a parse tree are that: (a) intermediate levels of nonterminals (such as T and F), have disappeared, because they were only in the concrete grammar to enforce the intended precedence and associativity, which are now directly reflected in the shape of the tree; and (b) there are no explicit parentheses, again because their grouping effect is now directly reflected in the shape of the tree.

4. The concrete parse tree (zoom in!):



The abstract parse tree:

