NFA Closure Properties

Sipser pages pages 58-63

NFAs also have closure properties

- We have given constructions for showing that DFAs are closed under
 - 1. Complement
 - 2. Intersection
 - 3. Difference
 - 4. Union
- We will now establish that NFAs are closed under
 - 1. Reversal
 - 2. Kleene star
 - 3. Concatenation

Reversal of $\epsilon\text{-NFAs}$

- Closure under reversal is easy using ε-NFAs. If you take such an automaton for L, you need to make the following changes to transform it into an automaton for L^{Rev}:
 - 1. Reverse all arcs
 - 2. The old start state becomes the only new final state.
 - 3. Add a new start state, and an ε -arc from it to all old final states.





- **Reverse all arcs**
- The old start state becomes the only new final state.
- Add a new start state, and an ϵ -arc 3. from it to all old final states.

Concatentation

• $L \bullet R = \{x \bullet y \mid x \text{ in } L \text{ and } y \text{ in } R\}$

- To form a new ε-NFA that recognizes the concatenation of two other ε-NFAs with the same alphabet do the following
 - Union the states (you might have to rename them)
 - Add an ϵ -transition from each final state of the first to the start state of the second.

Formally

• Let

$$-L = (Q_{L}, A, T_{L}, s_{L}, F_{L})$$

$$-R = (Q_{R}, A, T_{R}, s_{R}, F_{R})$$

$$-R = (Q_{L}, Q_{R}, A, T, s_{L}, F_{R})$$

- Where T s ε | $s \in F_L = S_R \cup T_L s \varepsilon$ T s c | $s \in Q_L = T_L s c$
 - $T S C | S \in Q_L T_L S C$ $T S C | S \in Q_R = T_R S C$







Kleene - Star

 If a language L is recognized by an NFA then so is the language L*

- Add a new state.
- Make it the start state in the new NFA.
- Add an ε -arc from this state to the old start state.
- Add ε-arcs from every final state to this new state.

Example

• Add a new state.

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- Make it the start state in the new NFA, and an accepting state.
- Add an ε-arc from this state to the old start state.
- Add ε-arcs from every final state to this new state

