

CS 581: Theory of Computation  
Fall 2009  
Mid-term exam  
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This is a closed-notes, closed-book exam.

1. PDA construction

- (a) Construct a PDA accepting the language

$$A = \{w \mid w \text{ has an equal number of } a\text{'s and } b\text{'s}\}$$

- (b) Justify your construction

- (c) Illustrate a computation of your machine on the string *abbba*.

2. Not Regular

Consider the language

$$A = \{w \mid w \text{ has an equal number of } a\text{'s and } b\text{'s}\}$$

Use this language to demonstrate three techniques for showing that  $A$  is not regular.

- (a) Show  $A$  is not regular using the pumping lemma.

- (b) Show  $A$  is of infinite index.

- (c) Show  $A$  is not regular by using closure properties and the fact that  $\{a^i b^i \mid i \geq 0\}$  is not regular.

3. Shuffle Let  $A, B \subseteq \Sigma^*$  be languages. Define the *shuffle* of  $A$  and  $B$ ,  $A \odot B$  as follows:

$$A \odot B = \{x_1 y_1 \cdots x_k y_k \mid x_1 \cdots x_k \in A \text{ and } y_1 \cdots y_k \in B, x_i, y_i \in \Sigma^*\}$$

For example,  $\{000\} \odot \{111\}$  includes the strings 000111, 111000, 101010, 010101, 011100, ...

Define the *shuffle closure* of  $A$ ,  $A^\otimes$ , as follows:

$$\begin{aligned} A^{\odot 0} &= \{\epsilon\} \\ A^{\odot^{n+1}} &= A^{\odot^n} \odot A \\ A^\otimes &= \bigcup_{i \geq 0} A^{\odot^i} \end{aligned}$$

- (a) Show the regular sets are closed under shuffle ( $\odot$ ).

- (b) Show the regular sets are *not* closed under shuffle closure ( $\otimes$ ).